

# TOPS-10 LINK Reference Manual

AA-0988D-TB, AD-0988D-T1, AD-0988D-T2

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This document describes LINK-10, the linking loader for TOPS-10.

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**SOFTWARE:** LINK-10 V6.0

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## PREFACE

This manual is the reference document for LINK, the TOPS-10 linking loader. The manual is aimed at the intermediate to highly-experienced applications programmer, and contains complete documentation of LINK.

Chapter 1 provides a general introduction to LINK.

Chapter 2 describes automatic use of LINK through one of the system commands DEBUG, EXECUTE, or LOAD. This chapter is sufficient for most loading tasks.

Chapter 3 describes direct use of LINK. This discussion is useful for large or complicated loads. This chapter also discusses libraries and library searches.

Chapter 4 describes output from LINK: executable programs, most output files, and LINK messages. Included are descriptions of the internal format of save (.EXE) files.

Chapter 5 discusses overlays, including overlay structures, overlay-related output files, the overlay handler and its messages, and the FUNCT. subroutine. This chapter has an extensive example of an overlay load. Many of the elements of this example are of interest outside the context of overlays.

Appendix A gives a technical description of the output from the language translators, which is in the form of REL Blocks.

Appendix B lists all LINK messages.

Appendix C describes the job data area.

The current copies of TOPS-10 documents are also useful:

TOPS-10 Utilities Manual

TOPS-10 Operating System Commands Manual

TOPS-10 MAKLIB User's Guide

TOPS-10/TOPS-20 MACRO Assembly Reference Manual

TOPS-10/TOPS-20 COBOL-68 Language Manual

TOPS-10/TOPS-20 COBOL-74 Language Manual

TOPS-10/TOPS-20 SPEAR Manual

TOPS-10/TOPS-20 ALGOL Programmer's Guide

TOPS-10/TOPS-20 FORTRAN Language Manual

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Processor Reference Manual



## CHAPTER 1

### INTRODUCTION TO LINK

LINK is the TOPS-10's linking loader. It merges independently compiled or assembled modules into a single executable program.

This merging process requires LINK to perform the following functions:

1. Perform the relocation calculations by converting relocatable addresses to virtual addresses, and by binding segments and PSECTs to addresses.
2. Resolve global symbol references by global chain fixups, Polish fixups, and library searches.
3. Produce an executable program by providing some JOBDAT information and a DDT runtime symbol table.

The virtual address space used for loading your program is not hardware memory. During loading and execution, the system simulates this virtual space by swapping code between disk and hardware memory as required. For simplicity, we will refer to the virtual address space as memory.

#### 1.1 INPUT TO LINK

The primary input to LINK is the output from the language translators; it is a binary file containing machine language code corresponding to your program, called object modules. Other input may include your commands to LINK, and libraries containing object modules.

##### 1.1.1 Object Modules

An object module is output from a language translator; it is part of a binary file (REL file) containing machine language code corresponding to your program. This file is formatted into blocks, called REL Blocks, that LINK recognizes and can handle appropriately. The format of each REL Block Type is described in Appendix A.

Most object modules contain relocatable code. This means that the addresses in the module are relative to the zero address. LINK loads the relocatable code at an arbitrary memory address, but adds a constant to each address referenced in the program. This resolves relative addresses to absolute addresses.

## INTRODUCTION TO LINK

Using relocatable code simplifies your programming task. Your programming task is simpler because you need not worry about the loading addresses of your programs.

Besides relocating and loading your object modules, LINK resolves values for global symbols: those that are defined in one module and used in others. LINK also resolves references to entry name symbols when modules containing these symbols are loaded.

Using symbols in your programs makes your programming simpler. If you need to revise a program, it is much easier to change the value of a symbol than to change each occurrence of the value. This is especially important for global symbols. You need only change the value in the defining module; the other modules do not need retranslation.

### 1.1.2 Commands to Link

LINK is controlled during loading by the command strings you give. Commands consist of file specifications and switches. LINK command strings are discussed in Chapter 3.

### 1.1.3 Libraries

A library is a file containing object modules that may be needed to resolve references in your program. For example, the FORTRAN library FORLIB contains subroutines that may be referenced by the output from the FORTRAN compiler. When loading FORTRAN-compiled code, LINK usually searches this library to satisfy any unresolved subroutine calls. Most language translators have their own libraries.

You can construct your own libraries, and have LINK search them for necessary subroutines. Libraries and searching are discussed in Section 3.4.

## 1.2 OUTPUT FROM LINK

The primary output from LINK is the executable program, called the core image. In the core image, all addresses are resolved to absolute memory locations, and all symbols (including subroutine calls) are resolved to absolute values or addresses.

This core image may be executed immediately or saved as a sharable save (.EXE) file. The .EXE file may be created automatically by LINK. This occurs if you specify /SSAVE when you run LINK, or if the program is too complex to be left in core with LINK.

You can also execute the core image under the control of a debugging program.

During its processing, LINK generates messages, which are output to your terminal or a log file. Some of these give information about LINK's operation; some warn you about possible problems; some identify errors. LINK messages are described in Appendix B.



## INTRODUCTION TO LINK

At your option, LINK can generate three special files: the map file, the log file, and the symbol file. The map file contains information about symbols in your program modules. The log file records LINK's messages so that you can save them. The symbol file contains a symbol table for the load and has a file extension of .SYM. LINK's output files are described in Chapter 4.

### 1.3 LINK'S OVERLAY FACILITY

If your program is larger than your available memory, you can use LINK's overlay facility to make it fit in memory. To do this, you define a tree structure for the program's modules. Then, at execution time, only part of the tree is in memory at one time. This reduces the amount of memory needed for execution. See Chapter 5 for a discussion of overlays.

### 1.4 LINK AND EXTENDED ADDRESSING

The KL Model B processor is capable of using an address space consisting of 32 sections, each containing 512 pages. As of TOPS-10 Version 7.03, programs have been able to reference this expanded address space. For information on using extended addressing with a specific programming language, consult the documentation for that language.

To load a program into a particular section, use the appropriate monitor command with the /USE switch. Refer to the TOPS-10 Operating System Commands Manual for more information.

Use PSECTs to load a program into a nonzero section. See Section 6.1 for information on loading PSECTs.

When loading a program that uses extended addressing, pay particular attention to the use of 18-bit and 30-bit addresses. If a program uses 30-bit addresses and you reference a 30-bit address as an 18-bit address, LINK truncates the 30-bit address and notifies you with the following message:

```
%LNKFTH Fullword value [symbol] truncated to halfword
```

LINK issues this warning if the truncation results in the loss of a section number. Refer to Appendix B for more information about this message.

While writing an extended addressing program, keep the following restrictions in mind:

- Programs that use overlays cannot use nonzero sections.
- Programs should not store executable code into locations 0 through 17 of nonzero sections. However, you can store data that is not executed in these locations. If you store data in locations 0 through 17, use global addresses to reference the locations. If you use local addresses, ACs (accumulators) are referenced instead. In Section 1, locations 0 through 17 refer to ACs.

## INTRODUCTION TO LINK

### 1.5 USING LINK

You have two ways to use LINK:

1. You can use LINK automatically by means of the LOAD, EXECUTE, or DEBUG system commands. This is the easiest and best way to load many programs. Chapter 2 describes automatic use of LINK.
2. You can run LINK directly by typing R LINK to the monitor. This is necessary only for very large or complicated loads, such as those involving overlays. Chapter 3 discusses direct use of LINK.

## CHAPTER 2

### USING LINK AUTOMATICALLY

The system commands LOAD, EXECUTE, and DEBUG invoke LINK automatically. Each of these commands uses a simple command string; the system converts the string into more complicated LINK commands.

This discussion of the LOAD, EXECUTE, and DEBUG commands does not attempt to describe them completely. Only those switches applying directly to loading will be discussed here. For a full discussion, see the TOPS-10 Operating System Commands Manual.

These system commands invoke LINK:

- The LOAD command uses LINK to load your object modules into memory, but does not execute the program. Before loading, your source files are compiled, if necessary; this compilation will occur if there are no object modules for the specified source files, or if the object files are older than their source files.
- The EXECUTE command uses LINK to load your program, and then executes the loaded program. Before loading, your source files are compiled, if necessary.
- The DEBUG command works like the EXECUTE command, except that your program is executed under the control of a debugging program. The debugging program that is loaded depends on the type of program being loaded. See the /TEST switch for a list of languages. The system uses the file extension to determine the language in which the program is written. Therefore, it is highly recommended that you use standard file extensions when naming the files of your programs. Standard file extensions are listed in the appropriate Commands Manual for the operating system.

#### 2.1 COMMAND FORMATS

The formats for the LOAD, EXECUTE, and DEBUG commands are the same. Each can accept a list of input file specifications and switches. The format for these commands is:

```
.command/switches input-spec/switches, input-spec/switches,...
```

Where the command is one of the three system commands (LOAD, EXECUTE, or DEBUG), input-spec is the file specification of the program you want to load, and the switches are any of the valid switches for the command.

## USING LINK AUTOMATICALLY

If you separate the input file specifications with commas, each source file will be compiled into a separate object file. If you separate the input file specifications with plus signs, they will be compiled into a single object file.

Section 2.3 shows examples of using LINK automatically.

### 2.2 COMMAND SWITCHES

You can use switches with the LOAD, EXECUTE, and DEBUG commands to control LINK's loading. Table 2-1 briefly describes some of the command switches that apply to LINK. Refer to the TOPS-10 Operating System Commands Manual for complete descriptions of the switches for these commands.

Table 2-1  
Switches for System Commands

Switch	Meaning
/COMPILE	Forces compilation of source files even if a sufficiently recent REL file exists.
/DDT	Loads DDT. This supersedes the default debugger selection, which is usually based on the extension of the first file in the command string.
/DEBUG	Causes the FORTRAN compiler to generate debugging information. /NODEBUG is the default.
/MAP	Produces a map file at the end of loading. This file shows all global symbols loaded.
/NOCOMPILE	Compiles source files only if their REL files are older than the source files. /NOCOMPILE is the default.
/NODEBUG	Prevents the FORTRAN compiler from generating debugging information.
/NOSEARCH	Suspends the effect of an earlier global /SEARCH switch. This is the default action.
/SEARCH	Loads only the modules from the specified library file that satisfy global references in the program.

You can use any LINK program switches with the system commands LOAD, EXECUTE, or DEBUG by using a special switch format. This format requires that you use a percent sign (%) instead of the usual slash (/), and that the entire switch specification be enclosed in double quotation marks ("). For example, you can pass the /ERRORLEVEL switch to LINK by using the command:

```
.EXECUTE MYPROG %"ERRORLEVEL:0"
```

## USING LINK AUTOMATICALLY

Used directly with LINK, the command strings would include:

```
*MYPROG/ERRORLEVEL:0
```

If you give more than one switch in this format, succeeding switches within the quotation marks must have the usual slashes:

```
.EXECUTE MYPROG%"ERRORLEVEL:0/SEGMENT:LOW"
```

LINK program switches are described in Section 3.2.

### 2.3 EXAMPLE OF USING LINK AUTOMATICALLY

For this example, the following program, named MYPROG.FOR, is used:

```
      TYPE 10
10    FORMAT (' This is written by MYPROG')
      STOP
      END
```

The following example shows an interactive execution of the program using the EXECUTE command:

```
.EXECUTE MYPROG.FOR (RET)
FORTRAN: MYPROG
MAIN.
LINK: Loading
[LNKXCT MYPROG execution]
This is written by MYPROG.
CPU time 0.16   Elapsed time 0.28

EXIT

.
```

The following example shows how to load a program for debugging using the DEBUG command:

```
.DEBUG MYPROG.FOR (RET)
FORTRAN: MYPROG
MAIN.
LINK: Loading
[LNKDEB FORDDT execution]

STARTING FORTRAN DDT

>>START
This is written by MYPROG.
CPU time 0.16   Elapsed time 0.42

EXIT

.
```



## CHAPTER 3 USING LINK DIRECTLY

If you have a loading task that cannot be handled conveniently by the EXECUTE, LOAD, or DEBUG system commands (such as loading overlays or PSECTs), you can load your program by using LINK directly. To do this, you must already have compiled or assembled all required object modules.

To use LINK directly, type R LINK to the system. LINK will respond with an asterisk:

```
.R LINK (RET)
*
```

Continue typing command strings, ending each one with a carriage return. For example,

```
.R LINK (RET)
*/OVERLAY (RET)
*TEST/LINK:TEST (RET)
*      /NODE:TEST SPEXP/LINK:SPEXP (RET)
*
```

A command string consists of file specifications and switches. You can continue a command string to the next line by typing a hyphen immediately before pressing carriage return; LINK continues the line by responding with a number sign (#). For example,

```
.R LINK (RET)
*MYPROG,MYMAP/MAP/CONTENTS:ALL- (RET)
# /ERRORLEVEL:0/LOG/LOGLEVEL:5 (RET)
*
```

The use of continuation lines is more efficient as the command scanner must be invoked for every distinct command string.

You can include a comment on a command line by beginning the comment with a semicolon; the remaining text on the line is not processed by LINK.

When LINK sees the end of the command string (a carriage return), it processes the entire string, then prints an asterisk to begin the next line. This processing continues until one of the following occurs:

1. LINK finds a /GO switch in a command string. It then completes loading and exits to system command level (if you did not specify execution), or passes control to the loaded program for execution.

## USING LINK DIRECTLY

2. A fatal error occurs. LINK prints an error message and exits to system command level.
3. A /RUN switch is encountered.
4. Either /EXIT or ^Z is encountered.

### 3.1 COMMAND STRING FORMAT

A LINK command string can contain file specifications, LINK switches, and command scanner switches. Command scanner switches are described in Section 3.2.1. LINK switches are described in Section 3.2.2.

Some LINK switches take output file specifications as arguments; some switches are suffixed to output file specifications. Other file specifications specify input files. For example, the following command string tells LINK to use an input file called MYREL.REL to generate a saved output file called MYEXE.EXE:

```
*MYREL,MYEXE/SAVE/GO
```

LINK supplies the missing parts of the file specifications from its defaults.

#### DEFAULTS

For output files, the defaults are:

device	DSK:	
file name	name of last module with start address or, if none, then nnnLNK where nnn is your job number with leading zeros if necessary	
extension	log file	LOG
	map file	MAP
	overlay file	OVL
	plotter file	PLT
	executable file	EXE
	symbol file	SYM

directory your current path, i.e. [30,5526,LINK,TEST]

For input files, the defaults are:

device DSK:

extension REL

directory your current path, i.e. [30,5526,LINK,TEST]

You can change these defaults by using the /DEFAULT switch (see Section 3.2.2).

You can have LINK read command strings from an indirect command file. To do this, prefix an at-sign (@) to the command file specification. For example, the following commands tell LINK to read all command strings from the file LNKPRG.CCL. (.CCL is the default file extension for indirect command files, but if .CCL is not found, .CMD will be tried.):

```
.R LINK (RET)
*@LNKPRG (RET)
```



## USING LINK DIRECTLY

### 3.2 SWITCHES TO LINK

LINK's handling of files depends on your use of LINK switches. There are two sets of switches to LINK. The first set of switches (command scanner switches) are optional switches that define your request to the system command scanner. These are described in Section 3.2.1. The second set of switches are switches to LINK that you can use to control and modify the linking and loading process. These are described in Section 3.2.2.

#### 3.2.1 Command Scanner Switches

The system SCAN module scans command lines for various system programs, one of which is LINK. You can include SCAN switches in your command strings for LINK; however, none of these switches is required in order to run LINK.

The following SCAN switches are meaningful to LINK. The remaining SCAN switches, which are listed in LINK's HELP file, are ignored by LINK.

Like LINK switches, SCAN switches are preceded by a slash (/), and can be abbreviated up to their first unique characters.

#### SCAN Switches Meaningful to LINK

Switch	Meaning
/ESTIMATE:n	Reserves n 128-word disk blocks for the specified output file. This does not work globally, but only on the file for which it is specified. By default, disk space is automatically allocated, as required, for files output to disk. This allocates blocks contiguously on the disk, thus providing improved access times.
/EXIT	Exits, but leaves LINK's core image in place.
/HELP:arg	Displays the HLP:LINK.HLP file. Specifying /HELP:SWITCHES types LINK and SCAN switches with no explanations.
/MESSAGE:keyword	Displays messages in the format specified by keyword. The keywords and their meanings are:  PREFIX           Displays only the message code from SCAN or LINK which are of the forms SCNxxx, or LNKxxx.  FIRST            Displays the prefix and a short message.  CONTINUATION    Displays the prefix and a longer message.
/NOOPTION	Ignores any LINK switches found in the file DSK:SWITCH.INI[,]

## USING LINK DIRECTLY

Switch	Meaning
/OPTION:name	Reads default LINK switches from the file SWITCH.INI, on the line that begins with LINK:name. If you use the /OPTION switch, it must appear in the first command string to LINK.
/PROTECTION:n	Assigns the 3-digit octal number n as the protection for the specified output file.
/RUN:file	Runs the specified program after loading is finished. This switch is ignored if you have specified program execution.
/RUNOFFSET:n	Begins execution of the program given by the /RUN switch at the address n locations after the normal start address. If you omit n, the default is 1. If you omit /RUNOFFSET, the default is 0.
/TMPFIL:file:"str"	Creates a TMPCOR file with the specified name. The name must be a 3-character filename. Usually, the string given in quotation marks is a command string to be executed by the program given in the /RUN switch when it is started one location after its normal start address (by /RUNOFFSET:1).

### 3.2.2 Link Switches

This section lists the switches that may be used to instruct LINK to take special action while loading your programs. The switches are described in this section in alphabetical order, and for each switch the following information is shown, if appropriate:

FORMAT  
FUNCTION  
EXAMPLES  
OPTIONAL NOTATIONS  
RELATED SWITCHES

Switches can be abbreviated to save typing. However, in most cases, the switch must include enough characters to make it unique from other switches. For example, the switch /NOSYMBOL cannot be abbreviated to /NOSY, because this result in a conflict with the switch /NOSYSLIB. However, /NOSYM is a unique set of characters, and thus is a legal abbreviation for /NOSYMBOL.

Certain switches that can be abbreviated to a single letter are:

/D for /DEBUG  
/E for /EXECUTE  
/G for /GO  
/H for /HELP  
/L for /LOCALS  
/M for /MAP  
/N for /NOLOCAL  
/S for /SEARCH  
/T for /TEST  
/U for /UNDEFINE  
/V for /VERSION

## USING LINK DIRECTLY

Many switches accept a value that may be specified in decimal (which is the default) or octal. If the value can be specified in octal, this is noted in the OPTIONAL NOTATIONS section of the switch description. To specify an octal value, type a number sign (#) before the octal number. For example, /ARSIZE:39 can be specified in octal as /ARSIZE:#47.

Some switches accept a value that specifies an amount of memory area. This value is interpreted as the number of pages, by default. However, this value may be specified as nK, where n is the number 1024-word blocks of memory (1K=2P). Where this is valid, it is noted in the OPTIONAL NOTATIONS section of the switch description. Some switches can be used either locally or globally (in particular, /LOCALS, /NOLOCAL, /NOSTART, /SEGMENT, /INCLUDE, /ONLY, /SEARCH, /START and /NOSEARCH). This means that if the switch is suffixed to a file specification, it applies only to that file; if it is not suffixed to a file specification, it applies to the files that follow on that command line. For example, in the following command strings /SEARCH is used both locally and globally:

1. \*FILE1,FILE2/SEARCH,FILE3
2. \*FILE4,/SEARCH FILE5,FILE6

In the first line, /SEARCH is suffixed to the file specification FILE2; only that file is loaded in search mode. In the second line, /SEARCH is not suffixed to a file specification; all the remaining files named in the command string are to be searched.

In general, a switch used globally is disabled at the end of its command string, unless it is overridden by another switch. The second switch, if used locally, will override the first only for the local file. If the second switch is used globally, it will persist for the following files. For example, in the following command string, a globally-used switch (/SEARCH) is overridden by a locally used switch:

```
*/SEARCH FILE1,FILE2/NOSEARCH,FILE3
```

In this command string, FILE1 and FILE3 will be loaded in search mode, but FILE2 will be loaded normally.

### NOTE

The effects of a global switch on the same line as a /GO switch persist beyond the /GO switch and apply to any modules loaded during library searches. However, for certain languages default switches may override user-specified global switches.

The following pages contain the switches and their descriptions, listed in alphabetical order.

## USING LINK DIRECTLY

### /ARSIZE

FORMAT        /ARSIZE:n

Where n is a positive decimal integer.

FUNCTION      Sets the size of the overlay handler's table of multiply-defined global symbols. Use this switch if you have received LNKARL, LNKDMA, and LNKABT messages in a previous attempt to load your program. These messages will give instructions for the argument to the /ARSIZE switch.

EXAMPLES      \*/ARSIZE:39 RET  
              \*

Allocates 39 words for the multiply-defined global symbol table in each link of an overlay structure.

OPTIONAL  
NOTATIONS     You can specify the table size in octal.

## USING LINK DIRECTLY

### /BACKSPACE

FORMAT        /BACKSPACE:n

              Where n is a positive decimal integer.

FUNCTION      Backspaces over n files on the current tape device.    (The  
              switch is ignored for non-tape devices.)

EXAMPLES     \*MTA0:/BACKSPACE:3 RET  
              \*

              Backspaces magtape MTA0: by three files.

OPTIONAL  
NOTATIONS    If you omit n, it defaults to 1.

RELATED  
SWITCHES     /MTAPE, /REWIND, /SKIP, /UNLOAD

## USING LINK DIRECTLY

### /COMMON

FORMAT        /COMMON:name:n

Where name is up to six SIXBIT-compatible ASCII characters.

n = a positive decimal integer.

FUNCTION      Allocates n words of labeled COMMON storage for FORTRAN and FORTRAN-compatible programs. The COMMON label is a name, which becomes defined as a global symbol.

For unlabeled COMMON storage, use .COMM. as the name, or simply omit the name.

You cannot expand a given COMMON area during loading. If your program modules define a given COMMON area to have different sizes, the module giving the largest definition must be loaded first. If the /COMMON switch gives the largest definition, it must precede the loading of the modules.

EXAMPLES     \*/COMMON:A:1000 (RET)  
\*

Creates a labeled COMMON area of 1000 words.

\*/COMMON:.COMM.:1000 (RET)  
\*

Creates an unlabeled COMMON area of 1000 words.

\*/COMMON::1000 (RET)  
\*

Creates an unlabeled COMMON area of 1000 words.

OPTIONAL  
NOTATIONS

You can specify the number of words in octal.

RELATED  
SWITCHES

/PSCOMMON

## USING LINK DIRECTLY

### /CONTENTS

FORMAT /CONTENTS: (keyword, ..., keyword)

FUNCTION Each keyword gives a symbol type to be included in the map file if the file is generated. To generate the map file, use the /MAP switch.

The keywords ALL, NONE, and DEFAULT reset all symbol types. Otherwise, using the /CONTENTS switch resets only those symbol types specified by keywords. In the following list of keywords, the defaults are in boldface:

ABSOLUTE	Include absolute symbols.
ALL	Include all symbols.
COMMON	Include COMMON symbols.
DEFAULT	Reset to LINK's defaults.
ENTRY	Include entry-name symbols.
GLOBAL	Include global symbols.
LOCALS	Include local symbols. The local symbols cannot be included in the map file unless the /LOCALS switch is also given.
NOABSOLUTE	Exclude absolute symbols.
NOCOMMON	Exclude COMMON symbols.
NOENTRY	Exclude entry-name symbols.
NOGLOBAL	Exclude global symbols.
NOLOCAL	Exclude local symbols.
NONE	Exclude all symbols.
NORELOCATABLE	Exclude relocatable symbols.
NOUNDEFINED	Exclude undefined symbols.
NOZERO	Exclude symbols in zero-length programs. (a zero-length program contains no code or data; it contains only symbol definitions, e.g., JOBDAT.)
RELOCATABLE	Include relocatable symbols.
UNDEFINED	Include undefined symbols.
ZERO	Include symbols in zero-length programs.

Only those symbols that satisfy all conditions in the keyword list will appear in the .MAP file. For example, if both the NOGLOBAL and RELOCATABLE settings are in force, all global symbols are excluded regardless of their relocatability.

EXAMPLES \*/CONTENTS: (NOCOMMON, NOENTRY) (RET)  
\*

Excludes COMMON and entry-name symbols.

\*/CONTENTS: ALL (RET)  
\*

Includes all symbols.

OPTIONAL NOTATIONS You can omit parentheses if you give only one keyword.

RELATED SWITCHES /MAP

## USING LINK DIRECTLY

### /CORE

FORMAT        /CORE:nP

Where n is a positive decimal integer.

FUNCTION       Gives the initial low segment core size for LINK. The core size given with /CORE should not be larger than that given with the /MAXCOR switch; if it is not, the core will be allocated initially, but the first time LINK expands core size, the allocation will be reduced to the size given by /MAXCOR.

EXAMPLES       \*/CORE:17P RET  
                 \*

Allocates 17P (1P=512 words) of core for the initial low segment.

OPTIONAL NOTATIONS    You can specify the core size in octal.  
                      You can also specify the argument in K instead of P.

RELATED SWITCHES     /FRECOR, /MAXCOR, /RUNCOR



## USING LINK DIRECTLY

### /COUNTERS

FORMAT /COUNTERS

FUNCTION Displays information about relocation counters on the terminal. A relocation counter is an address counter that LINK uses while loading relocatable code.

/COUNTERS returns the name, initial value, current value, and limit value of each counter. /COUNTERS first prints the following header:

```
Reloc. ctr.    initial value    current value    limit value
```

where:

Reloc. ctr. gives the name of relocation counter.

initial value is the origin of the relocation counter.

current value is the address of the next free location after the relocation counter has been loaded.

limit value is an upper bound that you set using /LIMIT or that LINK sets by default for the relocation counter. This upper bound defines a point the relocation counter should not load beyond. If /LIMIT is used and the counter loads beyond this bound, LINK returns messages. See /LIMIT for more information.

/COUNTERS may be used to determine the size of overlays when loading programs that might be too for the allocated memory space. Refer to Section 5.4 for more information.

You can also use /COUNTERS to determine the size of PSECTs when loading extended addressing programs or programs that use PSECTs to conserve memory space. Refer to Chapter 6.

EXAMPLES The following examples illustrate the various /COUNTERS displays.

The following display results from loading a module that does not contain code.

```
* /COUNTERS (RET)
[LNKRLC No relocation counters set]
*
```

The following display results from loading only absolute code.

```
* ABCODE /COUNTERS (RET)
[LNKRLC No relocation counters set
 Absolute code loaded]
*
```

## USING LINK DIRECTLY

The following display results from loading only PSECT code.

\*PSCODE/COUNTERS (RET)

[LNKRLC Reloc. ctr.	initial value	current value	limit value
PSCODE	20	25	1000000]

\*

The following display results from loading code that contains both absolute and PSECT code.

\*MIXED/COUNTERS (RET)

[LNKRLC Reloc. ctr.	initial value	current value	limit value
PSECTA	1400000	1400001	4000000
PSECTB	2500000	2500001	4000000
PSECTC	3600000	3600001	4000000
Absolute code loaded]			

\*

The following display results from loading two-segment formatted code.

\*TWOPT/COUNTERS (RET)

[LNKRLC Reloc. ctr.	initial value	current value	limit value
.LOW.	0	1642	1000000
.HIGH.	400000	400753	1000000]

\*

RELATED  
SWITCHES

/NEWPAGE, /SET, /LIMIT

## USING LINK DIRECTLY

### /CPU

FORMAT        /CPU:keyword

Keyword:      KA10  
               KI10  
               KL10  
               KS10

FUNCTION      This switch is used to override LINK's handling of the processor information found in the .REL files being loaded. (See the description of the type 6 block in Appendix A). Ordinarily LINK prints a warning if all .REL files being loaded together do not have identical CPU types. This switch can be used either to make LINK flag certain modules built for a specific CPU type (by specifying all but that CPU type as keywords to /CPU) or to suppress LINK's warning message (by specifying all the CPU types associated with the .REL files being loaded).

EXAMPLE       \*/CPU:KI10<RET>  
               \*

Will cause LINK to issue the %LNKCCD message if any modules with the KL10 CPU type are encountered.

OPTIONAL  
NOTATIONS     /CPU: (keyword,keyword)

USING LINK DIRECTLY

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## USING LINK DIRECTLY

### /DDEBUG

FORMAT /DDEBUG:keyword

FUNCTION Specifies a default debugging program to be loaded if the /DEBUG or /TEST switch appears without an argument.

The permitted keywords and the debugging programs they specify are listed below. Only those printed in **boldface** are supported by DIGITAL.

<b>ALGDDT</b>	Specifies <b>ALGDDT</b> as the default.
<b>ALGOL</b>	Specifies <b>ALGDDT</b> as the default.
<b>COBDDT</b>	Specifies <b>COBDDT</b> as the default.
<b>COBOL</b>	Specifies <b>COBDDT</b> as the default.
<b>DDT</b>	Specifies <b>DDT</b> as the default.
<b>FAIL</b>	Specifies <b>SDDT</b> as the default.
<b>FORDDT</b>	Specifies <b>FORDDT</b> as the default.
<b>FORTRAN</b>	Specifies <b>FORDDT</b> as the default.
<b>MACRO</b>	Specifies <b>DDT</b> as the default.
<b>PASCAL</b>	Specifies <b>PASDDT</b> as the default.
<b>PASDDT</b>	Specifies <b>PASDDT</b> as the default.
<b>SAIL</b>	Specifies the <b>SAIL</b> debugger as the default.
<b>SDDT</b>	Specifies the <b>SAIL</b> debugger as the default.
<b>SIMDDT</b>	Specifies <b>SIMDDT</b> as the default.
<b>SIMULA</b>	Specifies <b>SIMDDT</b> as the default.

EXAMPLES \*/DDEBUG:FORTRAN (RET)  
\*

Specifies **FORDDT** as the default debugging program for the /DEBUG or /TEST switch.

RELATED SWITCHES /DEBUG, /TEST

## USING LINK DIRECTLY

### /DEBUG

FORMAT /DEBUG:keyword

FUNCTION Requests loading of a debugging program and sets the start address for execution as the normal start address of the debugging program. The /DEBUG switch also sets the /EXECUTE switch because it is assumed that the program is to be executed. The /GO switch is still required to end loading and begin execution.

The /DEBUG switch turns on the /LOCALS switch for the remainder of the load. You can override this by using the /NOLOCAL switch, but the override lasts only during processing of the current command string.

Local symbols for the debugging program itself are never loaded.

If debugging overlaid programs, you must specify /DEBUG when loading the root node. (Refer to Section 5.4 for more information.)

The permitted keywords and the programs they load are listed below. Only those printed in boldface are supported by DIGITAL.

ALGDDT	Loads ALGDDT.
ALGOL	Loads ALGDDT.
COBDDT	Loads COBDDT.
COBOL	Loads COBDDT.
DDT	Loads DDT.
FAIL	Loads SDDT.
FORDDT	Loads FORDDT.
FORTRAN	Loads FORDDT.
MACRO	Loads DDT.
PASCAL	Loads PASDDT.
PASDDT	Loads PASDDT.
SAIL	Loads the SAIL debugger.
SDDT	Loads the SAIL debugger.
SIMDDT	Loads SIMDDT.
SIMULA	Loads SIMDDT.

If you give no keyword with /DEBUG, the default is either DDT or the debugging program specified by the /DDEBUG switch.

EXAMPLES \*/DEBUG:DDT RET  
\*

Loads DDT, sets the /EXECUTE switch, and specifies that the execution will be controlled by DDT.

OPTIONAL NOTATIONS Abbreviate /DEBUG to /D.

RELATED SWITCHES /DDEBUG, /TEST

## USING LINK DIRECTLY

### /DEFAULT

FORMAT        /DEFAULT:keyword filespec  
              filespec/DEFAULT:keyword

FUNCTION      Changes default specifications for input or output files.  
              The defaults specified remain in effect until changed by  
              another /DEFAULT switch.

The keywords allowed are:

              INPUT                Specifies the defaults for input file  
   specifications.

              OUTPUT               Specifies the defaults for output file  
   specifications.

For input files, the initial defaults are:

              device               DSK:  
              extension            REL  
              directory            User's connected directory

For output files, the initial defaults are:

              device               DSK:  
              filename            Name of main program  
              directory            User's connected directory

EXAMPLES     \*/DEFAULT:INPUT .BIN (RET)  
              \*

Resets input file default extension to BIN.

\*/DEFAULT:OUTPUT MTA0: (RET)  
\*

Resets output file default device to MTA0:.

OPTIONAL     If you omit the keyword, INPUT is assumed.  
NOTATIONS

## USING LINK DIRECTLY

### /DEFINE

FORMAT            /DEFINE:(symbol:value,...,symbol:value)

FUNCTION          Assigns each symbol the decimal value following it. This causes them to be global symbols. You can use the /UNDEFINED switch to get a list of any undefined symbols, and then define them with /DEFINE.

                  Defining an already defined symbol with /DEFINE generates an error message.

EXAMPLES          \*/UNDEFINED RET  
                  [LNKUGS 2 UNDEFINED GLOBAL SYMBOLS]  
                  A            400123  
                  IGOR        402017

                  \*/DEFINE:(A:591,IGOR:1) RET  
                  \*

                  Gives the decimal values 591 and 1 to A and IGOR, respectively.

OPTIONAL          You can give the value in octal by typing # before the value.

NOTATIONS        You can omit the parentheses if you define only one symbol. Specifying /DEFINE:FOO:BAR will define FOO to have the value of BAR if BAR is already defined.

RELATED          /UNDEFINED, /VALUE  
SWITCHES



## USING LINK DIRECTLY

### /ENTRY

FORMAT /ENTRY

FUNCTION Requests terminal typeout (in octal) of all entry name symbols loaded so far. Each entry name symbol will have been defined by an ENTRY statement (MACRO, FORTRAN, or BLISS), a FUNCTION statement (FORTRAN), a SUBROUTINE statement (FORTRAN, or COBOL), or a PROCEDURE declaration (ALGOL).

If you are using the overlay facility, /ENTRY requests only the entry name symbols for the current link.

EXAMPLES \*/ENTRY (RET)  
[LNKLSS LIBRARY SEARCH SYMBOLS (ENTRY POINTS)]  
SQRT. 3456  
\*

RELATED /NOENTRY  
SWITCHES

## USING LINK DIRECTLY

### /ERRORLEVEL

FORMAT        /ERRORLEVEL:n

FUNCTION      Suppresses terminal typeout of LINK messages with message level n and less, where n is a decimal number between 0 and 30 inclusive. You cannot suppress level 31 messages. LINK's default is /ERRORLEVEL:10.

See Appendix B for the level of each LINK message.

EXAMPLES     \*/ERRORLEVEL:10 (RET)  
              \*

Suppresses all messages of level 10 and less.

\*/ERRORLEVEL:0 (RET)  
\*

Permits typeout of all messages.

RELATED  
SWITCHES     /VERBOSITY, /MESSAGE

## USING LINK DIRECTLY

### /EXCLUDE

FORMAT        /EXCLUDE:(subroutine,...,subroutine)

FUNCTION      Prevents loading of the specified modules from the current file even if they are required to resolve global symbol references. You can use the /EXCLUDE switch for any of the following purposes:

- If a library has several modules with the same search symbols, you can select the module you want by excluding the others.
- You can prevent modules from giving multiple definitions of a symbol by selectively excluding one or more of them.
- In defining an overlay structure, you can delay loading of a module until a later link by excluding it.

EXAMPLES      \*/SEARCH LIBFIL.REL/EXCLUDE:(MOD1,MOD2) RET  
\*

Searches LIBFIL as a library but prevents loading of MOD1 and MOD2 even if they are referenced.

OPTIONAL NOTATIONS    You can omit the parentheses if you specify only one module.

RELATED SWITCHES      /INCLUDE

## USING LINK DIRECTLY

### /EXECUTE

FORMAT        /EXECUTE

FUNCTION      Tells LINK that the loaded program is to be executed beginning at its normal start address. Loading continues until a /GO switch is found.

              The /EXECUTE and /DEBUG switches are mutually exclusive.

EXAMPLES      \*/EXECUTE (RET)  
              \*

OPTIONAL  
NOTATIONS     You can abbreviate /EXECUTE to /E.

RELATED  
SWITCHES      /DEBUG, /GO, /RUN, /TEST

## USING LINK DIRECTLY

### /FRECOR

FORMAT        /FRECOR:nK

Where n is a positive decimal integer.

FUNCTION      Requires LINK to maintain a minimum amount of free memory after any expansions. LINK's default free memory is 4K. If you use the /FRECOR:nK switch, LINK computes n times 1024 words and maintains the resulting number of words of free memory, if possible.

If the modules to be loaded are quite large, a larger amount of free memory avoids some moving of areas.

The following areas may be expanded during loading:

1. ALGOL symbol information (AS).
2. Bound global symbols (BG).
3. Dynamic area (DY).
4. Fixup area (FX).
5. Global symbol tables (GS).
6. User's high segment code (HC).
7. User's low segment code (LC).
8. Local symbol tables (LS).
9. Relocation tables (RT).
10. Argument typechecking (AT)

Each of these areas has a lower bound, an actual upper bound, and a maximum upper bound. LINK normally maintains space between the actual and maximum upper bounds for each area. The total of these areas is at least the space given by the /FRECOR switch, if possible.

If the low segment code area exceeds the limit set by the user (with /MAXCOR) or by the system (CORMAX), LINK recovers free core by concatenating these free areas. When all this recovered space is used, one or more of the areas overflows to disk, and free core is no longer maintained.

EXAMPLES      \*/FRECOR:7K (RET)  
              \*

Maintains 7K of free core, if possible.

OPTIONAL      You can specify the free core in octal.  
NOTATIONS

RELATED        /CORE, /MAXCOR, /RUNCOR  
SWITCHES

## USING LINK DIRECTLY

/GO

FORMAT /GO

FUNCTION Ends loading after the current module. LINK then performs any required library searches, generates any required output files, and does one of the following:

- Begins execution at the normal start address of the loaded program (if you used /EXECUTE).
- Begins execution at the start address of the debugging program (if you used /DEBUG, or both /TEST and /EXECUTE).
- Exits to the monitor (if you used no execution switch).

EXAMPLES \*MYPROG/EXECUTE/GO (RET)

[LNKXCT MYPROG EXECUTION]

Begins execution of the loaded program at its normal start address.

\*MYPROG/DEBUG/GO (RET)

[LNKDEB DDT EXECUTION]

Begins execution of the loaded program at the normal start address of DDT.

OPTIONAL NOTATIONS Abbreviate /GO to /G.

RELATED SWITCHES /DEBUG, /EXECUTE, /RUN, /TEST

## USING LINK DIRECTLY

### /HASHSIZE

FORMAT        /HASHSIZE:n

Where n is a positive decimal integer.

FUNCTION      Gives a minimum for the initial size of the global symbol table. LINK selects a prime number larger than n for the initial size.

If you know that you will need a large global symbol table, you can save time and space by allocating space for it with /HASHSIZE. You should give a hash size at least 10 percent larger than the number of global symbols in the table.

If LINK gives the message [LNKRGS Rehashing Global Symbol Table] during a load, you should use the /HASHSIZE switch for future loads of the same program. The minimum hash size for loading a program appears in the header lines of the map file.

The default hash size is a LINK assembly parameter (initially 251 decimal).

EXAMPLES      \*/HASHSIZE:1000 (RET)  
              \*

Sets the hash size to the prime number 1021.

## USING LINK DIRECTLY

### /INCLUDE

**FORMAT**            /INCLUDE:(module,...,module)

**FUNCTION**        Specifies modules to be loaded regardless of any global requests for them.

In library search mode, an /INCLUDE switch requests loading of the specified modules. If the switch is associated with a file, the request is cleared after that file is searched. If not, the request persists until the modules are found.

When LINK is not in library search mode, the /INCLUDE switch associated with a file requests that only the specified modules be loaded, and the request is cleared after that file is processed. An /INCLUDE switch not associated with a file requests loading of the specified modules, and the request persists until the modules are found.

You can use /INCLUDE in an overlay load to force a module to be loaded in an ancestor link common to successor links that reference that module. This makes the module available to all links that are successors to its link.

**EXAMPLES**        \*/SEARCH LIB1/INCLUDE:(MOD1,MOD2) (RET)  
                  \*

Searches LIB1 and loads MOD1 and MOD2 even if they are not referenced.

**OPTIONAL NOTATIONS**    You can omit the parentheses if you specify only one module.

**RELATED SWITCHES**    /EXCLUDE, /NOINCLUDE, /MISSING



## USING LINK DIRECTLY

### /LIMIT

FORMAT        /LIMIT:psect:address

FUNCTION      Allows you to specify an upper bound for a specific PSECT. In the format description, psect should be the PSECT name, which has been defined with either the /SET switch or in one of the modules already loaded. Address should be the upper bound address of the specified PSECT, expressed in either numeric or symbolic form. This address should be one greater than the highest location which may be loaded in the PSECT.

If the PSECT grows beyond the address specified in the /LIMIT switch, LINK will send a warning to your terminal, but will continue to process input files and to load code. The warning message will take the following form:

%LNKPEL PSECT <psect> exceeded limit of <address>

No chained references will be resolved, and LINK will suppress program execution, producing the following fatal error:

?LNKCF5 Chained fixups have been suppressed

This action prevents unintended PSECT overlays. PSECT overlays can cause loops and other unpredictable behavior, because LINK uses address relocation chains in the user image that is being built.

#### EXAMPLE

```
*TEST1 (RET)
*/COUNTERS (RET)
[LNKRLC RELOC. CTR.        INITIAL VALUE    CURRENT VALUE    LIMIT VALUE
  .LOW.                    0                140              1000000
          Q                1000            4000             1000000
          R                4500            10500            1000000]
*/LIMIT:Q:4000 (RET)
*TEST2 (RET)
%LNKPEL PSECT Q EXCEEDED LIMIT OF 4000
DETECTED IN MODULE .MAIN FROM FILE DSK:TEST2.REL
*/COUNTERS (RET)
[LNKRLC RELOC. CTR.        INITIAL VALUE    CURRENT VALUE    LIMIT VALUE
  .LOW.                    0                140              1000000
          Q                1000            5000             4000
          R                4500            10500            1000000]
*TEST/SAVE/GO (RET)
%LNKPOV PSECTS R AND Q OVERLAP FROM ADDRESS 4500 TO 5000
?LNKCF5 CHAINED FIXUPS HAVE BEEN SUPPRESSED
```

In this example, a program named TEST1, which contains two PSECTS, is loaded. The PSECTS are named Q and R. After TEST1 is loaded, the /COUNTERS switch shows that the upper bound of PSECT Q is 4000.

## USING LINK DIRECTLY

The /LIMIT switch is used to limit PSECT Q to 4000.

A second program, TEST2, also requires storage for PSECT Q. Therefore, when TEST2 is loaded, LINK produces a warning to the effect that the limit that was set has been exceeded. The /COUNTERS switch shows that PSECT Q now requires an upper bound of 5000.

When the programs are started (with /GO), LINK produces the POV warning message and the CFS fatal error message.

### RELATED SWITCH

/COUNTER

### OPTIONAL NOTATIONS

A defined global symbol can be used to specify the limit value, e.g., /LIMIT:FOO:BAR.

## USING LINK DIRECTLY

### /LINK

FORMAT /LINK:name

Where name is up to 6 RADIX-50 compatible characters.

FUNCTION Directs LINK to give the specified name to the current core image and outputs the core image to the overlay file. /LINK is used to close an overlay link. LINK first performs any required library searches and assigns a number to the link.

For a discussion of overlay structures, see Chapter 5.

The current core image has all modules loaded since the beginning of the load or since the last /LINK switch.

EXAMPLES \*SPEXP/LINK:ALPHA (RET)  
\*

Loads module SPEXP and outputs the core image to the overlay file as a link called ALPHA.

OPTIONAL NOTATIONS If you omit the link name, LINK uses only its assigned number.

RELATED SWITCH /NODE

## USING LINK DIRECTLY

### /LOCALS

FORMAT        /LOCALS

FUNCTION       Includes local symbols from a module in the symbol table. LINK does not need these tables, but you may want them for debugging. To have the symbol table included in a program use the /SYMSEG switch.

The /LOCALS and /NOLOCAL switches may be used either locally or globally. If the switch is suffixed to a file specification, it applies only to that file; if it is not suffixed to a file specification, it applies to all following files in the same command line.

EXAMPLES       \*/SYMSEG (RET)  
                \*/LOCALS A,B/NOLOCAL,C,/NOLOCAL D (RET)

\*

Loads A with local symbols, B without local symbols, C with local symbols, and D without local symbols.

OPTIONAL NOTATIONS    You can abbreviate /LOCALS to /L.

RELATED SWITCHES     /NOLOCAL, /SYMSEG

## USING LINK DIRECTLY

### /LOG

**FORMAT**           logfilespec/LOG

**FUNCTION**       Specifies a file specification for the log file (see Section 4.2.2). Any LINK messages output before the /LOG switch is encountered are not entered in the log file.

**EXAMPLES**       \*LOGFIL/LOG RET  
                  \*  
                  Specifies the file DSK:LOGFIL.LOG in the user's directory.

                  \*TTY:/LOG RET  
                  \*  
                  Directs log messages to the user's terminal.

**OPTIONAL NOTATIONS**   You can omit all or part of the logfilespec.  
                  The defaults are:

device	DSK:
filename	name of main program
extension	LOG
directory	your logged-in directory

                  You can change the defaults using the /DEFAULT switch.

**RELATED SWITCHES**   /LOGLEVEL

## USING LINK DIRECTLY

### /LOGLEVEL

FORMAT        /LOGLEVEL:n

FUNCTION      Suppresses logging of LINK messages with level n and less, where n is a decimal number between 0 and 30 inclusive. You cannot suppress level 31 messages.

See Appendix B for the level of each LINK message.

The default is /LOGLEVEL:10.

EXAMPLES      \*/LOGLEVEL:0 RET  
\*

Logs all messages.

RELATED  
SWITCHES      /ERRORLEVEL, /LOG

## USING LINK DIRECTLY

### /MAP

FORMAT mapfilespec/MAP:keyword

FUNCTION Specifies a file specification for the map output file (see Section 4.2). The contents of the file are determined by the /CONTENTS switch or its defaults.

Permitted keywords and their meanings are:

END Produces a map file at the end of the load. This is the default if you omit the keyword.

ERROR Produces a map file if a fatal error occurs. Any modules loaded after this switch will not appear in the log. To ensure that a .MAP file is generated, specify this switch before the loading of .REL files.

NOW Produces a map file immediately. Library searches will not have been performed unless forced.

EXAMPLES \*MAPFIL/MAP:END RET  
\*

Generates a map in the file DSK:MAPFIL.MAP in your disk area at the end of loading.

OPTIONAL NOTATIONS You can omit all or part of the mapfilespec. The defaults are:

device	DSK:
filename	name of main program
extension	MAP
directory	user's connected directory

You can change the defaults using the /DEFAULT switch.

You can abbreviate /MAP to /M.

RELATED SWITCHES /CONTENTS

## USING LINK DIRECTLY

### /MAXCOR

**FORMAT**            /MAXCOR:nP  
Where n is a positive decimal integer.

**FUNCTION**        Requires LINK to overflow low segment code larger than n pages to disk. The disk overflow may contain symbol areas, low segment code, and high segment code.

                  If you specify MAXCOR smaller than the memory already used, disk overflow occurs at the next expansion and memory is reduced to your specified MAXCOR.

                  If you specify MAXCOR smaller than LINK's minimum requirement, you will get an error message. You must then use another /MAXCOR switch to enlarge the core.

**EXAMPLES**        \*/MAXCOR:30P   
                  \*

                  Specifies a maximum core use of 30 pages (1P=512 words).

**OPTIONAL NOTATIONS**    You can specify the core in octal.  
                  You can specify the switch argument as K instead of P.

**RELATED SWITCHES**    /CORE, /FRECOR, /RUNCOR



## USING LINK DIRECTLY

### /MAXNODE

FORMAT /MAXNODE:n

Where n is a positive decimal integer.

FUNCTION Specifies the number of links to be defined when the overlaid program requires more than 256 links. LINK will allocate extra space in the OVL file for certain fixed-length tables based on the number of links specified with this switch.

Note that this switch must be placed after the /OVERLAY switch and it must precede the first /NODE switch in the set of commands to LINK.

EXAMPLES \*TEST/OVERLAY/MAXNODE:500   
\*

Reserves space for 500 defined links. See Chapter 5 for a discussion on overlays.

RELATED SWITCHES /OVERLAY

## USING LINK DIRECTLY

### /MISSING

FORMAT /MISSING

FUNCTION Requests terminal typeout of modules requested with the /INCLUDE switch that have not yet been loaded.

EXAMPLES \*MYPROG (RET)  
\*/SEARCH/INCLUDE:(MOD1,MOD2) LIB1 (RET)  
\*/MISSING (RET)  
[LNKIMM 1 INCLUDED MODULE MISSING]  
\*LIB2/INCLUDE:(MOD2) (RET)  
\*/MISSING (RET)  
[LNKIMM NO INCLUDED MODULES MISSING]  
\*

This example shows the use of /MISSING to see if all the required modules have been loaded. The module MOD2 was not yet loaded, and it was in LIB2.

In response to the first use of the switch, LINK indicated that one necessary module was missing. After the missing module was included (module named LIB2), the switch is used again. LINK responded to the second use of the switch by indicating that all necessary modules were present.

RELATED SWITCHES /INCLUDE, /UNDEFINED

## USING LINK DIRECTLY

### /MTAPE

FORMAT /MTAPE:keyword

FUNCTION Specifies tape operations to be performed on the current device. (A tape device remains current only until end-of-line or until another device is specified, whichever is earlier.) The switch is ignored if the current device is not a tape.

The operation is performed immediately if /MTAPE is given with an input file or with an already initialized output file. Otherwise, the operation is performed when the output file is initialized.

The valid keywords and the operations they specify are:

MTBLK	Writes 3 inches of blank tape.
MTBSF	Backspaces one file.
MTBSR	Backspaces one record.
MTDEC	Initializes DIGITAL-compatible 9-channel tape.
MTEOF	Writes an end-of-file mark.
MTEOT	Spaces to logical end-of-tape.
MTIND	Initializes industry-compatible 9-channel tape.
MTREW	Rewinds tape to the load point (BOT).
MTSKF	Skips one file.
MTSKR	Skips one record.
MTUNL	Rewinds and unloads tape.
MTWAT	Waits for tape I/O to finish.

EXAMPLES \*MTA0:/MTAPE:MTEOT (RET)  
\*MTA0:/MAP:NOW (RET)  
\*

Spaces to logical end-of-tape on MTA0: and writes a map file.

RELATED SWITCHES /BACKSPACE, /REWIND, /SKIP, /UNLOAD

## USING LINK DIRECTLY

### /NEWPAGE

FORMAT            /NEWPAGE:keyword

FUNCTION         Sets the relocation counter to the first word of the next page. If the counter is already at a new page, this switch is ignored.

The permitted keywords and their relocation counters are:

      LOW         Resets the low-segment counter to new page.  
                  If you omit the keyword, this is the default.

      HIGH        Resets the high-segment counter to new page.

EXAMPLES        \*/NEWPAGE:HIGH (RET)  
                  \*SUBR1 (RET)  
                  \*/NEWPAGE:LOW (RET)  
                  \*SUBR2 (RET)  
                  \*

Sets the high-segment counter to a new page, loads SUBR1, sets the low-segment counter to a new page, and loads SUBR2. Note that SUBR1 and SUBR2 are not necessarily loaded into the high and low segments respectively; the /NEWPAGE switch sets a counter, but does not force the next loaded module into the specified segment.

RELATED  
SWITCHES        /SET, /COUNTER

## USING LINK DIRECTLY

### /NODE

FORMAT /NODE:argument

FUNCTION Opens an overlay link. /NODE places LINK's relocation counter at the end of a previously defined link in an overlay structure, which becomes the immediate ancestor to the next link defined. (For a discussion of overlay structures, see Chapter 5.)

The /NODE switch must precede any modules to be placed in the new link.

Three kinds of arguments are permitted:

- A name given with a previous /LINK switch. LINK will place the relocation counter at the end of the specified link.
- A negative number (-n). LINK backs up n links along the current path.
- A positive number n or 0. LINK begins further loading at the end of link number n. You can use 0 to begin loading at the root link.

### NOTE

It is recommended that you use a link name (or 0 for the root link) rather than a nonzero number. This is because a change in commands defining an overlay may change some of the link numbers.

EXAMPLES For examples defining overlay structures, see Chapter 5.

RELATED SWITCHES /LINK, /OVERLAY, /PLOT

## USING LINK DIRECTLY

### /NOENTRY

FORMAT        /NOENTRY:(symbol,symbol,...)

FUNCTION      Deletes entry name symbols from LINK's overhead tables when loading overlays, thereby saving space at run time. If you know that execution of the current load will not reference certain entry points, you can use /NOENTRY to delete them.

/NOENTRY differs from /NOREQUEST in that /NOREQUEST deletes requests for symbols, while /NOENTRY deletes symbols that might be requested.

EXAMPLES     \*/ENTRY (RET)  
              [LNKLSS LIBRARY SEARCH SYMBOLS (ENTRY POINTS)]  
              SQRT.    3456  
              \*/NOENTRY:(SQRT.) (RET)  
              \*/ENTRY (RET)  
              \*

Deletes SQRT. so that it cannot be used to fulfill a symbol request.

OPTIONAL NOTATIONS    You can omit the parentheses if only one symbol is given.

RELATED SWITCHES     /ENTRY, /EXCLUDE, /NOEXCLUDE, /INCLUDE, /NOINCLUDE,  
                      /MISSING, /REQUEST, /NOREQUEST

## USING LINK DIRECTLY

### /NOINCLUDE

FORMAT        /NOINCLUDE

FUNCTION      Clears requests for modules that were specified in a previous /INCLUDE.

EXAMPLE      \*LIB1/INCLUDE:(MOD1,MOD3) (RET)  
              \*/NOINCLUDE (RET)  
              \*

Loads MOD1 and MOD3 from LIB1. However, if the modules are not found immediately, stop searching.

RELATED  
SWITCHES     /INCLUDE, /EXCLUDE, /MISSING

## USING LINK DIRECTLY

### /NOINITIAL

FORMAT        /NOINITIAL

FUNCTION      Prevents loading of LINK's initial global symbol table (JOB DAT). The /NOINITIAL switch cannot operate after the first file specification because JOB DAT will be already loaded. The initial global symbol table contains the JBxxx symbols described in Appendix C.

The /NOINITIAL switch is commonly used for:

- Loading LINK itself (to get the latest copy of JOB DAT).
- Loading a private copy of JOB DAT (to alter if necessary).
- Building an .EXE file that will eventually run in executive mode (for example, a monitor or bootstrap loader).

EXAMPLES     \*/NOINITIAL RET  
              \*



## USING LINK DIRECTLY

### /NOLOCAL

FORMAT /NOLOCAL

FUNCTION Suspends the effect of a preceding /LOCALS switch so that local symbol tables will not be loaded with their modules.

The /LOCALS and /NOLOCAL switches may be used either locally or globally. If the switch is suffixed to a file specification, it applies only to that file; if it is not suffixed to a file specification, it applies to all following files in the same command string.

This switch is useful if you need to conserve memory space, because local symbols are loaded into the low segment by default.

EXAMPLES \*/LOCALS A,B/NOLOCAL,C,/NOLOCAL D RET  
\*

Loads A with local symbols, B without local symbols, C with local symbols, and D without local symbols.

OPTIONAL NOTATIONS Abbreviate /NOLOCAL to /N.

RELATED SWITCHES /LOCALS

## USING LINK DIRECTLY

### /NOREQUEST

FORMAT        /NOREQUEST:(symbol,symbol,...)

FUNCTION      Deletes references to links from LINK's overhead tables when loading overlay programs. If you know that the execution of the current load will not require certain links, you can use /NOREQUEST to delete references to them.

/NOREQUEST differs from /NOENTRY in that /NOENTRY deletes symbols that might be requested, while /NOREQUEST deletes the requests for them.

EXAMPLES     \*/REQUEST (RET)  
              [LNKRER REQUEST EXTERNAL REFERENCES]  
              ROUTN.  
              SQRT.  
              \*/NOREQUEST:(ROUTN.,SQRT.) (RET)  
              \*/REQUEST (RET)  
              \*

Deletes references to ROUTN. and SQRT.

OPTIONAL NOTATIONS    You can omit the parentheses if only one symbol is given.

RELATED SWITCH        /NOENTRY

## USING LINK DIRECTLY

### /NOSEARCH

FORMAT /NOSEARCH

FUNCTION Suspends the effect of a previous /SEARCH switch. Files named between a /SEARCH and the next /NOSEARCH are searched as libraries, so that modules are loaded only to resolve global references.

The /SEARCH and /NOSEARCH switches may be used either locally or globally. If the switch is suffixed to a file specification, it applies only to that file; if it is not suffixed to a file specification, it applies to all following files in the same command string.

EXAMPLES \*FILE1 (RET)  
\*/SEARCH A,B/NOSEARCH,C,/NOSEARCH D (RET)  
\*

Searches A, loads B, searches C, and loads D.

RELATED SWITCHES /SEARCH

## USING LINK DIRECTLY

### /NOSTART

FORMAT        /NOSTART

FUNCTION      Directs LINK to disregard any start addresses found after the /NOSTART switch. Normally LINK keeps the most recent start address found, overwriting any previously found. The /NOSTART switch prevents this replacement.

EXAMPLES     \*MAIN1,/NOSTART MAIN2,MAIN3 (RET)  
              \*

Directs LINK to save the start address from MAIN1 instead of replacing it with other start addresses from MAIN2 and MAIN3.

RELATED  
SWITCHES     /START

## USING LINK DIRECTLY

### /NOSYMBOL

FORMAT        /NOSYMBOL

FUNCTION      Prevents construction of user symbol tables. Symbols are then not available for the map file, but the header for the file can still be generated by the /MAP switch.

The /NOSYMBOL switch prevents writing an ALGOL SYM file if it would otherwise have been written.

If you do not need the map file or symbols, you can speed loading by using the /NOSYMBOL switch.

EXAMPLES     \*/NOSYMBOL (RET)  
              \*

## USING LINK DIRECTLY

### /NOSYSLIB

FORMAT /NOSYSLIB:(keyword,...,keyword)

FUNCTION Prevents automatic search of the system libraries named as keywords. LINK usually searches system libraries at the end of loading to satisfy unresolved global references. The /NOSYSLIB switch prevents this search.

The /NOSYSLIB switch can also be used to terminate searching of libraries that were specified in a previous /SYSLIB switch. When you specify searching of a library with /SYSLIB, that library will continue to be searched for every module you load. You can use /NOSYSLIB to specify libraries that should not be searched. Refer to /SYSLIB for more information.

The permitted keywords and the libraries they specify are listed below. Only those printed in **boldface** specify libraries supported by DIGITAL.

<b>ANY</b>	<b>Prevents all library searches.</b>
<b>ALGOL</b>	<b>Prevents search of ALGLIB.</b>
BCPL	Prevents search of BCPLIB.
<b>COBOL</b>	<b>Prevents search of LIBOL or C74LIB.</b>
F40	Prevents search of LIB40.
<b>FORTRAN</b>	<b>Prevents search of FORLIB.</b>
NELIAC	Prevents search of LIBNEL.
PASCAL	Prevents search of PASLIB.
SAIL	Prevents search of SAILIB.
SIMULA	Prevents search of SIMLIB.

EXAMPLES \*/NOSYSLIB:(ALGOL,COBOL) RET  
\*

Prevents search of the system libraries ALGLIB and LIBOL.

OPTIONAL NOTATIONS If you omit keyword it defaults to ANY.  
You can omit parentheses if only one keyword is given.

RELATED SWITCH /SYSLIB

## USING LINK DIRECTLY

### /NOUSERLIB

FORMAT filespec/NOUSERLIB

FUNCTION Discontinues automatic searching of the specified file at each /LINK or /GO switch. If you need a file searched for some links but not others, you can use the /USERLIB and /NOUSERLIB switches to enable and disable automatic search of the file.

EXAMPLES \*/OVERLAY (RET)  
\*MYFORL/USERLIB:FORTRAN (RET)  
\*MOD1/LINK:MOD1 (RET)  
\*/NODE:MOD1 MOD2/LINK:MOD2 (RET)  
\*MYFORL/NOUSERLIB (RET)  
\*

Loads the overlay handler; requests search of MYFORL as a FORTRAN library; loads MOD1 and MOD2 as links; discontinues search of MYFORL.

OPTIONAL NOTATIONS If you omit the filespec, LINK discontinues search of all user libraries.

RELATED SWITCHES /USERLIB

## USING LINK DIRECTLY

### /ONLY

FORMAT /ONLY:keyword

FUNCTION Directs LINK to load only the specified segment of two-segment modules. The permitted keywords are:

HIGH Loads only high segments.  
LOW Loads only low segments.  
BOTH Loads both segments.

The /ONLY switch is ignored for one-segment modules and for PSECTed modules.

EXAMPLES \*/ONLY:HIGH MOD1,MOD2 (RET)  
\*MOD3/ONLY:BOTH (RET)  
\*

Loads high segment for MOD1 and MOD2; loads both segments for MOD3.



## USING LINK DIRECTLY

### /OTSEGMENT

FORMAT /OTSEGMENT:keyword

FUNCTION Specifies the time and manner of loading the object-time system.

The permitted keywords are:

DEFAULT Suspends the effect of a previous /OTSEGMENT:SHARABLE or /OTSEGMENT:NONSHARABLE switch.

HIGH As for /OTSEGMENT:SHARABLE.

LOW AS FOR /OTSEGMENT:NONSHARABLE.

NONSHARABLE Loads the object-time system into user core image at load time. The user program may have code in both segments. The object-time system may have code in both segments.

SHARABLE Binds the object-time system at execution time. The user program is in the low segment and the object-time system is in the high segment.

LINK's default action is to bind the object-time system at execution time. This normal action occurs if none of the following are true.

- You specify /OTSEGMENT:NONSHARABLE.
- You have loaded any code into the high segment.
- You have specified /SEGMENT:HIGH for some modules.
- You have specified /SYMSEG:HIGH.
- Your low segment is too big for sharable object-time systems to fit.

If any of these is true, a non-sharable object-time system is loaded as part of your program.

EXAMPLES \*MYPROG/SYSLIB/OTSEGMENT:NONSHAR (RET)  
\*

Loads a non-sharable copy of the object-time system as part of your program.

RELATED SWITCHES /SEGMENT

## USING LINK DIRECTLY

### /OVERLAY

FORMAT	filespec/OVERLAY:(keyword,...,keyword)
FUNCTION	Initiates construction of an overlay structure. For a discussion of overlay structures, see Chapter 5.  The permitted keywords and their meanings are listed below. The default settings are printed in <b>boldface</b> .
<b>ABSOLUTE</b>	Specifies that links are absolute. This is the default situation when overlays are loaded. The inverse situation is to use /OVERLAY:RELOCATABLE. Relocatable overlays are described in Chapter 5.
<b>LOGFILE</b>	Outputs runtime overlay messages to your terminal.
<b>NOLOGFILE</b>	Suppresses output of runtime overlay messages.
<b>NOWARNING</b>	Suppresses overlay warning messages.
<b>PATH</b>	Specifies that each link path will be loaded with its link.
<b>RELOCATABLE</b>	Specifies that links are relocatable.
<b>TREE</b>	Specifies that the overlay will have a tree structure.
<b>WARNING</b>	Outputs overlay warning messages to user terminal.
<b>WRITABLE</b>	Specifies that the links are writable. Refer to Chapter 5 for more information.
EXAMPLES	See Chapter 5.
OPTIONAL NOTATIONS	You can omit the parentheses if only one keyword is given.

## USING LINK DIRECTLY

### /PATCHSIZE

FORMAT        /PATCHSIZE:n

Where n is a positive decimal integer.

FUNCTION      Allocates n words of storage to precede the symbol table. The allocated storage is in the same segment (high or low) as the symbol table. The default is /PATCHSIZE:64.

The storage allocated is available for patching or for defining new symbols with DDT, and is identified by the global symbol "PAT.."

EXAMPLES     \*/SYMSEG:HIGH/PATCHSIZE:200 (RET)  
              \*

Loads the symbol table in the high segment after allocating 200 words between the last loaded module and the symbol table.

OPTIONAL NOTATIONS    You can specify the patchsize in octal.

RELATED SWITCHES     /SYMSEG

## USING LINK DIRECTLY

### /PLOT

FORMAT filespec/PLOT

FUNCTION Directs LINK to output a tree diagram of your overlay structure. You can have the diagram formatted for a plotter (by default) or for a line printer (by giving the device as LPT:).

Each box in the diagram shows a link number, its name (if you gave one with the /LINK switch), and its relationship to other links (as defined by your commands).

The /PLOT switch cannot precede the /OVERLAY switch.

EXAMPLES See Chapter 5.

OPTIONAL LINK has default settings for the size of the overlay NOTATIONS diagram and the increment for drawing lines. You can override these by giving the /PLOT switch in the form:

filespec/PLOT:(LEAVES:value,INCHES:value,STEPS:value)

Where the values for each parameter define:

INCHES Width of diagram in inches. The defaults are INCHES:29 for plotter and INCHES:12 for line printer.

LEAVES Number of links without successors that can appear in one row. The defaults are LEAVES:16 for plotter and LEAVES:8 for line printer.

STEPS Increments per inch for drawing lines. The defaults are STEPS:100 for plotter and STEPS:20 for line printer.

For line printer diagrams, you cannot give INCHES or LEAVES different from the defaults. The STEPS parameter should be between 10 and 25.

For plotter diagrams, you should give INCHES and LEAVES in a ratio of about 2 to 1. For example, INCHES:40 and LEAVES:20.

If LINK cannot design the diagram on one page, it will automatically design subtrees for diagrams on more pages.

RELATED /LINK, /NODE, /OVERLAY SWITCHES

## USING LINK DIRECTLY

### /PLTTYP

**FORMAT** /PLTTYP:keyword

**FUNCTION** Allows a user to specify the type of plot file to be generated by the /PLOT switch.

**KEYWORDS**

DEFAULT Generate output for a printer only if the device is a printer or terminal.

PLOTTER Generate output for a plotter.

PRINTER Generate output for a printer.

**EXAMPLES**

.

.R LINK

\*TEST/OVERLAY  
\*DSK:TEST/PLOT/PLTTYP:PRINTER  
\*OVL0,OVL1/LINK:TEST  
\*/NODE:TEST OVL2 /LINK:LEFT  
\*/NODE:LEFT OVL5 /LINK:LEFT1  
\*/NODE:LEFT OVL6 /LINK:LEFT2  
\*/NODE:TEST OVL3,OVL4 /LINK:RIGHT  
\*TEST /SAVE /GO

EXIT

.

Causes all output from the /PLOT switch to be in line printer format.

**RELATED SWITCHES** /PLOT

## USING LINK DIRECTLY

### /PSCOMMON

FORMAT        /PSCOMMON:psect:common

FUNCTION      Specifies where LINK is to load COMMON blocks. This switch causes the FORTRAN common specified by the argument common to be loaded into the PSECT specified in the argument psect. Use the /PSCOMMON switch before loading the specified common and before declaring the common's size with the /COMMON switch.

/PSCOMMON only affects common blocks defined with the /COMMON switch. If the common block is created by a REL block, /PSCOMMON is ignored, and the PSECT specified by the REL file is used.

EXAMPLES      In the following example, /SET defines the SECTA PSECT's origin, /PSCOMMON specifies that COMABC is loaded into SECTA, and /COMMON defines the common size.

```
*/SET:SECTA:30000000 (RET)
*/PSCOMMON:SECTA:COMABC (RET)
*/COMMON:COMABC:10000 (RET)
*PROG (RET)
*
```

RELATED  
SWITCHES      /COMMON

# USING LINK DIRECTLY

## /REDIRECT

FORMAT /REDIRECT:lowpsect:highpsect

FUNCTION Loads two-segment formatted REL files as part of a program using PSECTs. The argument lowpsect is the name of the PSECT to receive the low-segment code and highpsect is the name of the PSECT to receive the high-segment code.

You must redirect both the high and the low segments.

EXAMPLES The following example loads a two-segment program (TWOVRT), and displays the low- and high-segment values using /COUNTERS.

```
*TWOVRT (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr. initial value current value limit value
  .LOW. 0 1642 1000000
  .HIGH. 400000 400753 1000000]
*
```

Next, PSECT origins are set for PSHI and PSLO, .LOW. is redirected into PSLO, .HIGH. is redirected into PSHI, and /COUNTERS is used to display PSHI and PSLO values.

```
*/SET:PSHI:400010 (RET)
*/SET:PSLO:3500 (RET)
*/REDIRECT:PSLO:PSHI (RET)
*TWOVRT (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr. initial value current value limit value
  PSHI 400010 400753 1000000
  PSLO 3500 5202 1000000]
*
```

## USING LINK DIRECTLY

### /REQUEST

FORMAT        /REQUEST

FUNCTION      Requests terminal typeout of all external references to other links.

If you use /REQUEST to get the names of external references, you can then either delete the references with the /NOREQUEST switch, or load the referenced modules.

EXAMPLES     \*/REQUEST<RET>  
              [LNKRER REQUEST EXTERNAL REFERENCES]  
              ROUTN.  
              SQRT.  
              \*/NOREQUEST:ROUTN.<RET>  
              \*/SEARCH LIB1<RET>  
              \*

Obtains the external references ROUTN. and SQRT.; deletes the request for ROUTN.; searches the file LIB1 for a module containing the entry point SQRT.

RELATED  
SWITCHES     /NOREQUEST



## USING LINK DIRECTLY

### /REQUIRE

FORMAT        /REQUIRE:(symbol,...,symbol)

FUNCTION      Generates global requests for the specified symbols. LINK uses these symbols as library search symbols (entry points).

/REQUIRE differs from /INCLUDE in that /INCLUDE requests a module by name, while /REQUIRE requests an entry name symbol. Thus you can use /REQUIRE to specify a function (for example, SQRT.) even if you do not know the module name.

You can use /REQUIRE to load a module into a link common to all links that reference the module.

Note that the global requests generated by the /REQUIRE switch do not use the standard calling sequence, and are therefore not visible to the /REQUEST switch.

EXAMPLES      \*/UNDEFINED (RET)  
              [LNKUGS NO UNDEFINED GLOBAL SYMBOLS]  
              \*/REQUIRE:(ROUTN.,SQRT.) (RET)  
              \*/UNDEFINED (RET)  
              [LNKUGS 2 UNDEFINED GLOBAL SYMBOLS]  
                  ROUTN.  
                  SQRT.  
              \*

OPTIONAL NOTATIONS      You can omit the parentheses if only one symbol is given.

RELATED SWITCHES        /SEARCH, /NOSEARCH

## USING LINK DIRECTLY

### /REWIND

FORMAT        /REWIND

FUNCTION      Rewinds the current input or output device if the device  
is a tape. If not, the switch is ignored.

EXAMPLES     \*MTA0:/REWIND (RET)  
              \*

Rewinds tape on MTA0:.

## USING LINK DIRECTLY

### /RUNAME

FORMAT        /RUNAME:name

FUNCTION      Assigns a job name for execution of your program. This name is stored inside the monitor and is used in the SYSTAT display.

If you give no /RUNAME switch, the default name is the name of the module with the start address. If there is no such module, the name nnnLNK is used, where nnn is your 3-digit job number.

EXAMPLES     \*/RUNAME:LNKDEV (REF)  
              \*

Assigns the name LNKDEV for job execution.

## USING LINK DIRECTLY

### /RUNCOR

**FORMAT**            /RUNCOR:nP  
                      Where n is a positive decimal integer.

**FUNCTION**         Allocates n pages of memory for the program's low segment  
                      at execution time.

**EXAMPLES**         \*MYPROG/RUNCOR:22P/EXECUTE/GO (RET)  
                      [LNKXCT MYPROG Execution]  
                      Specifies that MYPROG will execute with a low-segment  
                      allocation of 22 pages.

**OPTIONAL**         You can specify the argument in octal.  
**NOTATIONS**        You can specify the switch argument as K instead of P.

## USING LINK DIRECTLY

### /SAVE

FORMAT filespec/SAVE

FUNCTION Directs LINK to create an .EXE file. The file extension defaults to .EXE. For example, if you enter FOO.BAR/SAVE, LINK creates a file FOO.EXE.

Note that if you want to run the saved file with the system command, the file extension must be .EXE.

EXAMPLES \*MYPROG   
\*DSKZ:GOODIE.EXE/SAVE/GO   
\*

Directs LINK to save the linked version of MYPROG as GOODIE.EXE on DSKZ:.

RELATED /SSAVE  
SWITCH

## USING LINK DIRECTLY

### /SEARCH

FORMAT /SEARCH

FUNCTION Directs LINK to load selectively from all following files up to the next /NOSEARCH or /GO. These files are searched as libraries, and only modules whose entry point name resolves a global request are loaded.

Using /NOSEARCH discontinues the library search mode, but for each link the system libraries are still searched (unless you used the /NOSYSLIB switch), and user libraries are still searched (if you used the /USERLIB switch).

The /SEARCH and /NOSEARCH switches may be used either locally or globally.

Note that search requests in .TEXT blocks may be processed in the reverse order of the entered /SEARCH switches. Keep this in mind when specifying the order in which the modules are to be searched. See Block Types Greater Than 3777 in Appendix A for more information.

EXAMPLES \*/SEARCH A,B/NOSEARCH,C,/NOSEARCH D

Searches A, loads B, searches C, and loads D.

RELATED SWITCHES /NOSEARCH

## USING LINK DIRECTLY

### /SEGMENT

FORMAT /SEGMENT:keyword

FUNCTION Specifies which segment is to be used for loading following modules. FORTRAN object code is an exception; both segments are loaded into the low segment unless one or more of the following is true:

- You used the /OTSEGMENT:NONSHARABLE switch.
- You used the /SEGMENT:HIGH switch to load code into the high segment.
- You used the /SEGMENT:DEFAULT switch to load code into both segments.
- Some code is already loaded into the high segment.

The keywords for the /SEGMENT switch are:

DEFAULT Suspends effect of /SEGMENT:LOW or /SEGMENT:HIGH.

HIGH Load into high segment, even if impure code.

LOW Loads into low segment.

NONE Same as DEFAULT.

If the switch is suffixed to a file specification, it applies only to that file; if it is not suffixed to a file specification, it applies to all following files in the same command string.

EXAMPLES \*/SEGMENT:LOW MOD1,MOD2,/SEGMENT:HIGH MOD3 (RET)  
\*

Loads MOD1 and MOD2 into the low segment; loads MOD3 into the high segment even if its code is impure.

RELATED SWITCHES /OTSEGMENT

## USING LINK DIRECTLY

### /SET

FORMAT        /SET:name:address

Where name is .HIGH., .LOW., or a PSECT name, and address is a 30-bit octal address or a defined symbol.

FUNCTION      Sets the loading position of a PSECT, or sets the .HIGH. or .LOW. relocation counter.

For setting the loading position of a PSECT, name is the name of the PSECT, and address is a virtual memory address. The /SET switch must precede the modules that will make up the specified PSECT. The /SET switch is not needed if the REL files already contain origin information.

### NOTE

If you load PSECTs so that the resulting core image contains gaps, you must generate an EXE file and execute that file (rather than executing the loaded core image). It is good practice to generate an .EXE file for all PSECTed programs.

If you do not ask for an .EXE file and you need one, LINK will generate one for you.

EXAMPLES      \*/SET:A:200000 (RET)  
\*

Specifies that the PSECT named A is to be loaded with its origin at address 200000.

\*/SET:.HIGH.:400000 (RET)  
\*

Sets the high segment relocation counter .HIGH. to the address 400000. Note that saying /SET:.HIGH. causes a high segment to appear and a vestigial JOB DAT area to be built.

RELATED        /COUNTER, /LIMIT  
SWITCHES



## USING LINK DIRECTLY

### /SEVERITY

FORMAT        /SEVERITY:n

FUNCTION      Specifies that messages of severity level greater than n will terminate the load, where n is a decimal number between 0 and 30 inclusive. Level 31 messages always terminate the load.

The defaults are /SEVERITY:24 for timesharing jobs, and /SEVERITY:16 for batch jobs.

EXAMPLES     \*/SEVERITY:30 (RET)  
              \*

Specifies that only level 31 messages are fatal.

## USING LINK DIRECTLY

### /SKIP

FORMAT /SKIP:n

Where n is a positive decimal integer.

FUNCTION Skips forward over n files on the current tape device. (A tape device remains current only until end-of-line or until another device is specified, whichever occurs first.) If the device is not a tape, the switch is ignored.

EXAMPLES \*MTA0:/SKIP:4 RET  
\*

Skips forward over 4 files on MTA0:.

RELATED /BACKSPACE, /MTAPE, /REWIND, /UNLOAD  
SWITCHES

## USING LINK DIRECTLY

### /SPACE

FORMAT        /SPACE:n

Where n is a positive decimal integer.

FUNCTION       Specifies that n words of memory will follow the current link at execution time. This memory allocation will not increase the size of the overlay file, but it will increase the size of the program at run time.

The /SPACE switch is used to allocate space for use by the object time system. The OTS uses this space for I/O buffers, and as scratch space in FORTRAN and heap space in ALGOL.

You should place the /SPACE switch before the first /LINK switch, to ensure allocation for the root link. It is possible to allocate space after one or more overlays are linked. This might be useful if an overlay has unusual storage requirements: buffers for a file which is open only while that overlay is resident, or a large local matrix. To allocate space between overlays, use /SPACE when loading the overlay that will be using this file or matrix. LINK allows one /SPACE switch for the root node, and one for each overlay.

The default amount of memory allocated, if you do not specify /SPACE, is 2000 for the root link and 0 (zero) for other links.

If the space allocated for a relocatable link is too small, the overlay handler can relocate it. If the space allocated for an absolute link is too small, a fatal error occurs.

EXAMPLES

```
* /OVERLAY (RET)
* TEST /SPACE:90 /LINK:MAIN (RET)
*        /NODE:MAIN SUB1 /LINK:SUB1 (RET)
*        /NODE:MAIN SUB2 /LINK:SUB2 (RET)
*
```

Allocates 90 words of memory to follow the root link for the program. See Chapter 5 for a discussion on overlay.

OPTIONAL  
NOTATIONS

You can specify the number of words in octal.

## USING LINK DIRECTLY

### /SSAVE

FORMAT filespec/SSAVE

FUNCTION Specifies the same actions as the /SAVE switch, except that at execution time the program's high segment will be sharable. The extension of the file created by LINK is always EXE; other file extensions are ignored. If, for example you enter FOO.BAR/SSAVE, LINK creates a file FOO.EXE.

EXAMPLES \*DSK:SHRPRG/SSAVE RET  
\*

Requests a sharable save file DSK:SHRPRG.EXE.

RELATED SWITCH | /SAVE

## USING LINK DIRECTLY

### /START

FORMAT        /START:symbol  
              /START:address  
              /START

Where symbol is a defined global symbol and address is a 30-bit octal address.

FUNCTION      Specifies the start address for the loaded program, and prevents replacement by any start addresses found later. You can use the /START switch with no argument to disable a previously given /NOSTART switch.

EXAMPLES      \*MAIN1/START:ENTRY1,MAIN2,MAIN3 (RET)  
              \*

Defines the start address as ENTRY1 in MAIN1, and prevents replacement of this start address by any others found in MAIN2 or MAIN3.

RELATED  
SWITCHES      /NOSTART

## USING LINK DIRECTLY

### /SUPPRESS

FORMAT /SUPPRESS:symbol

Where symbol is a previously defined global symbol.

FUNCTION Used to suppress a previously defined global symbol. If the symbol is unknown, this switch has no effect. Use this switch if a global symbol is defined in two modules and you wish to suppress one of the definitions.

LINK suppresses a defined global symbol by setting its definition to undefined in the global symbol table. LINK does not remove the symbol definition from the symbol table. As a result, the symbol table built for debugging contains both the old and new values of the symbol.

Since LINK sets the symbol to undefined in the symbol table, it expects that a subsequent module will be loaded that contains a global definition for the symbol. If the symbol is not defined later, LINK issues the Undefined Global Symbol (LNKUGS) error.

EXAMPLES In the following example, the ENTPTR symbol is used in both the TEST and TEST2 programs. First, LINK is run, TEST is loaded, and the value of ENTPTR is shown using the /VALUE switch.

```
.R LINK (RET)
*TEST (RET)
*/VALUE:ENTPTR (RET)
[LNKVAL Symbol ENTPTR 140 defined]
```

Next, ENTPTR's value is suppressed using /SUPPRESS and its current value is shown. Note that the value is now undefined.

```
*/SUPPRESS:ENTPTR (RET)
*/VALUE:ENTPTR (RET)
[LNKVAL Symbol ENTPTR 0 undefined]
```

Finally TEST2 is loaded and the value is shown again.

```
*TEST2 (RET)
*/VALUE:ENTPTR (RET)
[LNKVAL Symbol ENTPTR 200 defined]
```

In the next example, TEST and TEST2 are loaded, but ENTPTR is not suppressed after TEST is loaded. In this example, LINK issues the Multiply-defined global symbol warning.

```
.R LINK (RET)
*TEST (RET)
*TEST2
%LNKMDS Multiply-defined global symbol ENTPTR
Detected in module .MAIN from file TEST2.REL
Defined value = 140, this value = 200
```

## USING LINK DIRECTLY

### /SYFILE

**FORMAT** filespec/SYFILE:keyword

**FUNCTION** Requests LINK to output a symbol file to the given filespec, and sets the /SYMSEG:DEFAULT switch. If you previously specified /NOSYM, the /SYFILE switch has no effect.

The symbol file contains global symbols sorted for DDT. If you used the /LOCALS switch, it also contains local symbols, module names, and module lengths.

The permitted keywords and their meanings are:

**ALGOL** Requests symbols in ALGOL's format. The first word of the table is "XWD 1044, count." The remaining words are copied out of Type 1044 REL blocks. If an ALGOL main program has been loaded, then /SYFILE:ALGOL becomes the default.

**RADIX50** Requests symbols in Radix-50 format. The first word of the table is negative. Each symbol requires two words in the table: the first is the symbol name in Radix-50 format; the second is the symbol value.

**TRIPLET** Requests symbols in triplet format. The first word of the table is zero. Each symbol requires three words in the table: the first word contains flags; the second is the symbol name in SIXBIT; the third is the symbol value.

**EXAMPLES** \*SYMBOL/SYFILE (RET)  
\*

Creates a symbol file called SYMBOL with the symbols in Radix-50 format.

**OPTIONAL NOTATIONS** If you omit the keyword, RADIX50 is assumed.

## USING LINK DIRECTLY

### /SYMSEG

FORMAT /SYMSEG:keyword

FUNCTION Places the symbol table so that it will not be overwritten during execution or debugging.

Keywords and their meanings are:

DEFAULT Places the symbol table in the low segment, except for overlaid programs. For overlays, symbols are not loaded.

HIGH Places the symbol table in the high segment.

LOW Places the symbol table in the low segment.

NONE Prevents loading of the symbol table.

PSECT:name Places the symbol table at the end of the PSECT (after allocating any space required by the /PATCHSIZE switch).

EXAMPLES \*/SYMSEG:LOW (RET)  
\*

Places the symbol table in the program low segment.

RELATED SWITCHES /LOCALS, /NOLOCALS



## USING LINK DIRECTLY

### /SYSLIB

FORMAT /SYSLIB:keyword

FUNCTION Forces searching of one or more system libraries, immediately after you end the command line. LINK will also automatically search a system library if code from the corresponding compiler has been loaded. By default, LINK searches the system libraries that are appropriate for the language compiler, after all the modules of the program are loaded. /SYSLIB forces the search to take place immediately.

After you specify a library with /SYSLIB, the library you specified will be searched every time you load a module, until you use /NOSYSLIB to end searching of that library.

The permitted keywords and the libraries they specify are listed below. Those printed in **boldface** specify libraries supported by DIGITAL.

<b>ANY</b>	<b>Forces search of all system libraries.</b>
<b>ALGOL</b>	<b>Forces search of ALGLIB.</b>
BCP	Forces search of BCPLIB.
<b>COBOL</b>	<b>Forces search of LIBOL or C74LIB.</b>
F40	Forces search of LIB40.
<b>FORTRAN</b>	<b>Forces search of FORLIB.</b>
NELIAC	Forces search of LIBNEL.
PASCAL	Forces search of PASLIB.
SAIL	Forces search of SAILIB.
SIMULA	Prevents search of SIMLIB.

EXAMPLES \*TEST1/SYSLIB:ALGOL (RET)  
\*TEST2/NOSYSLIB:ALGOL  
\*

Where TEST1 is a FORTRAN module, LINK will search both FORLIB and ALGLIB for TEST1. Where TEST2 is a FORTRAN module, LINK will search only FORLIB when TEST2 is loaded.

OPTIONAL NOTATIONS You can omit the keyword. LINK will search all libraries for which corresponding code has been loaded.

RELATED SWITCHES /NOSYSLIB

## USING LINK DIRECTLY

### /TEST

FORMAT /TEST:keyword

FUNCTION Loads the debugging program indicated by keyword. Unlike the /DEBUG switch, /TEST causes execution to begin in the loaded program (not in the debugging module). This switch is useful if you expect the program to run successfully, but want the debugger available in case the program has errors.

The /TEST switch turns on the /LOCALS switch for the remainder of the load. You can override this by using the /NOLOCAL switch, but the override lasts only during processing of the current command string.

Local symbols for the debugging module itself are never loaded.

The permitted keywords and the programs they load are listed below. Only those printed in **boldface** are supported by DIGITAL.

<b>ALGDDT</b>	Loads <b>ALGDDT</b> .
<b>ALGOL</b>	Loads <b>ALGDDT</b> .
<b>COBDDT</b>	Loads <b>COBDDT</b> .
<b>COBOL</b>	Loads <b>COBDDT</b> .
<b>DDT</b>	Loads <b>DDT</b> .
<b>FAIL</b>	Loads <b>SDDT</b> .
<b>FORDDT</b>	Loads <b>FORDDT</b> .
<b>FORTRAN</b>	Loads <b>FORDDT</b> .
<b>MACRO</b>	Loads <b>DDT</b> .
<b>PASCAL</b>	Loads <b>PASDDT</b> .
<b>PASDDT</b>	Loads <b>PASDDT</b> .
<b>SAIL</b>	Loads the <b>SAIL</b> debugger.
<b>SDDT</b>	Loads the <b>SAIL</b> debugger.
<b>SIMDDT</b>	Loads <b>SIMDDT</b> .
<b>SIMULA</b>	Loads <b>SIMDDT</b> .

EXAMPLES \*MYPROG/TEST:FORTRAN (RET)  
\*

Loads MYPROG and FORDDT.

OPTIONAL NOTATIONS If you give no keyword with /TEST, the default is either DDT or the debugging program specified by the /DDEBUG switch.

RELATED SWITCHES /DDEBUG, /DEBUG

## USING LINK DIRECTLY

### /UNDEFINED

FORMAT        /UNDEFINED

FUNCTION       Requests terminal typeout (in octal) of undefined global symbols. You can use /UNDEFINED to get a list of undefined symbols, and then define them with the /DEFINE switch.

EXAMPLES       \*/UNDEFINED RET  
                [LNKUGS 2 UNDEFINED GLOBAL SYMBOLS]  
                  A        400123  
                  IGOR    402017  
                \*/DEFINE:(A:591,IGOR:1) RET  
                \*

Gives the decimal values 591 and 1 to A and IGOR, respectively.

OPTIONAL NOTATIONS    You can abbreviate /UNDEFINE to /U.

RELATED SWITCHES     /DEFINE, /VALUE

## USING LINK DIRECTLY

### /UNLOAD

FORMAT        device/UNLOAD

FUNCTION      Rewinds and unloads the specified tape device. (This switch is ignored if the current device is not a tape device.) The /UNLOAD is not performed until the current file processing is completed.

EXAMPLES     \*MTA0:/UNLOAD   
              \*

              Rewinds and unloads MTA0.

RELATED  
SWITCHES     /BACKSPACE, /MTAPE, /REWIND, /SKIP

## USING LINK DIRECTLY

### /UPTO

FORMAT        /UPTO:addr

Where addr is a 30-bit octal address that specifies the upper limit to which the symbol table can grow. The address can be replaced by a symbol.

FUNCTION      Sets an upper limit to which the symbol table can expand.

EXAMPLE       \*/UPTO:550000 (RET)  
              \*

Included in a FORTRAN load, this switch would override the default upper bound for the symbol table. This might be used if FOROTS begins above 550000.

RELATED  
SWITCH        /SYMSEG

## USING LINK DIRECTLY

### /USERLIB

FORMAT filespec/USERLIB:(keyword,...,keyword)

FUNCTION Directs LINK to search the user library given by filespec before searching system libraries. The keyword indicates that the given library is to be searched only if code from the corresponding compiler was loaded.

Keywords and their meanings are given below. Only those printed in **boldface** indicate compilers and libraries supported by DIGITAL.

ALGOL	Search as an ALGOL library.
ANY	Always search this library.
BCPL	Search as a BCPL library.
COBOL	Search as a COBOL library.
<b>FORTRAN</b>	Search as a <b>FORTRAN</b> library.
NELIAC	Search as a NELIAC library.
PASCAL	Search as a PASCAL library.
SAIL	Search as a SAIL library.
SIMULA	Search as a SIMULA library.

EXAMPLES \*MYFORL/USERLIB:FORTRAN (RET)  
\*

Directs LINK to search the user library MYFORL (before searching FORLIB) if any FORTRAN-compiled code is loaded.

OPTIONAL NOTATIONS You can omit the parentheses if only one keyword is given.

RELATED SWITCHES /NOUSERLIB, /SYSLIB

## CHAPTER 4

### OUTPUT FROM LINK

The primary output from LINK is the executable program formed from your input modules and switches. During its processing, LINK gives errors, warnings, and informational messages. At your option, LINK can generate any of several files.

#### 4.1 THE EXECUTABLE PROGRAM

The executable program that LINK generates (called the core image) consists mostly of data and machine instructions from your object modules. In the core image, all relocatable addresses have been resolved to absolute addresses, and the values of all global references have been resolved.

You have several options for loading the program, depending on the purpose of the load. Those options are:

- Execute the program. To do this, include the /EXECUTE switch any place before the /GO switch. LINK will pass control to your program for execution.
- Execute the program under the control of DDT. To do this, use the /DEBUG switch before the first input file specification.
- Execute the program and debug it after execution. To do this, use the /TEST and /EXECUTE switches before the first input file specification. After execution, type DDT to the system to enter the debugging program.
- Save the core image as an EXE file. To do this, use the /SAVE switch. See Section 4.2.

#### 4.2 OUTPUT FILES

At your option, LINK can produce any of the following output files:

- Saved (executable) file.
- Log file.
- map file.

## OUTPUT FROM LINK

- Symbol file.
- Plotter file (see Section 5.1).
- Overlay file (see Section 5.1).

### 4.2.1 Executable Files

The executable file, sometimes called the saved or .EXE file, is a copy of the completed core image generated by LINK. You can create an executable file by supplying the /SSAVE switch before the /GO switch when you are loading the program with direct commands to LINK. The executable file will retain the same file name as the source program, with a file extension .EXE.

Alternatively, you can type the file specification, followed by /SSAVE (or /SAVE), and the executable file will be written to the file you specified. If you load the program with the system LOAD command, you may then save the executable file by typing the system SAVE command.

You can run the executable file later, without running LINK, by using the system command RUN, or the two system commands GET and START. The following section describes the internal format of the executable file.

See Chapter 3 for descriptions of /SAVE and /SSAVE switches.

**4.2.1.1 Format of Sharable Save Files** - A sharable save file is divided into two main areas: the directory area, which contains information about the structure of the file, and the data area, which contains the data of the file.

The following diagram illustrates the general format of a sharable save file:

```
Directory Area:  =====
                  !  Directory Section  !
                  !-----!
                  ! Entry Vector Section !
                  !-----!
                  ! Terminating Section !
                  =====
Data Area:      !  Data Section  !
                !               !
                !               !
                !               !
                !               !
                !               !
                !               !
                !               !
                !               !
                !               !
                !               !
                =====
```

The directory area of the sharable save file has three distinct sections: the directory section, the entry vector section, and the terminating section. The size of the directory area depends on the access characteristics of the pages in the data area of the save file.

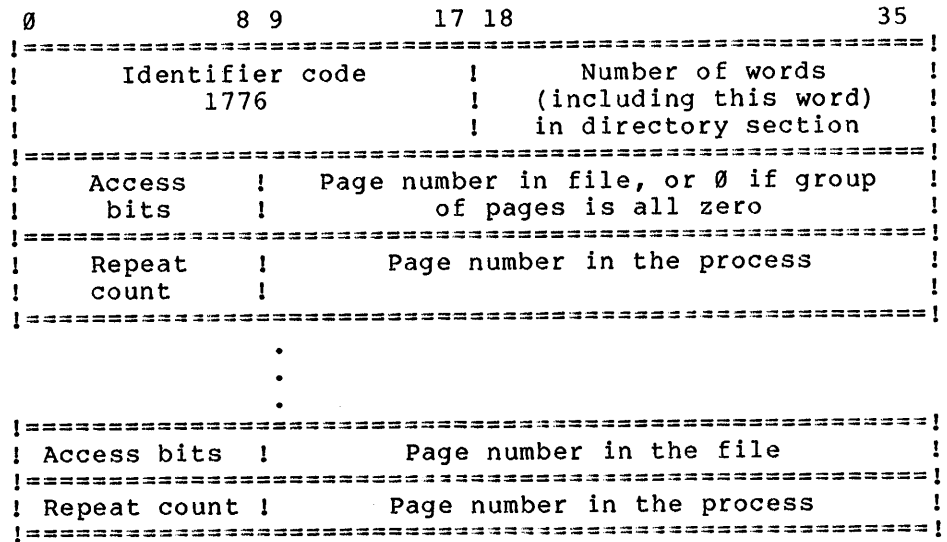


OUTPUT FROM LINK

Each of the sections in the directory area begins with a header word containing its identifier code in the left half and its length in the right half. Each section is described in the following paragraphs.

The directory section is the first of the three sections and describes groups of contiguous pages that have identical access. The length of this section varies according to the number of groups that can be generated from the data portion of the save file. The more data pages that can be combined into a single group, the fewer groups required, and the smaller the directory section.

The format of the directory section is as follows:



PSECT attributes are used to set the access bits. Refer to the description of Block Type 24 in Appendix A.

The directory section has one header word containing 1776 count, where count is the number of words in the directory section, including this header word.

The header word is followed by word pairs. Each pair of words is formatted as follows:

Word 0, .SVFPPF, specifies the access flags and the page number in the file. The flag bits are:

Bit	Symbol	Meaning
0	SV%HIS	Page is in high segment.
1	SV%SHR	Page is sharable.
2	SV%WRT	Page is writable.
3	SV%CON	Page is concealed.
4	SV%SYM	Page is part of symbol table.

Word 1, .SVPPC, contains the repeat count in Bits 0-8, and the page number in the process in the remaining bits.

OUTPUT FROM LINK

The repeat count is the number (minus 1) of consecutive pages in the group described by the word pair. Pages are considered to be in a group when the following three conditions are met:

1. The pages are contiguous.
2. The pages have the same access.
3. The pages are allocated but not loaded.

A group of all zero pages is indicated by a file page number of 0.

The word pairs are repeated for each group of pages in the address space.

The entry vector section follows the directory section. It points to the first word of the entry vector, and gives the length of the vector.

```

0                               17 18                               35
!=====!
! Identifier code           ! Number of words           !
!       1775                ! (including this word)    !
!                               ! in entry vector section  !
!=====!
!                               254000                               !
!=====!
!                               Starting address                       !
!=====!

```

This format is the default. However, if you make special provisions in your program, the format becomes the following. (Refer to the description of Block Type 7 in Appendix A for further information.)

```

0                               17 18                               35
!=====!
! Identifier code           ! Number of words           !
!       1775                ! (including this word)    !
!                               ! in entry vector section  !
!=====!
!                               Number of words in entry vector      !
!=====!
!                               Address of entry vector                !
!=====!

```

The data for this section is the address of the entry vector.

The terminating section, called the end section, always immediately precedes the data section. The format of the terminating section is the following:

```

!=====!
! Identifier code           !                               !
!       1777                !                               !
!=====!

```

The data area follows the terminating section, beginning at the next page boundary.

## OUTPUT FROM LINK

### 4.2.2 LOG Files

A LOG file is generated if you use the /LOG switch. LINK then writes most of its messages into the specified file. You can control the kinds of messages entered in the LOG file by using the /LOGLEVEL switch. For an example of a LOG file, see Section 5.1.

### 4.2.3 Map files

The map file is generated if you use the /MAP switch. LINK constructs a symbol map in this file. The kinds of symbols included depends on your use of the /CONTENTS, /LOCALS, /NOLOCALS, /NOINITIAL, and /NOSYMBOLS switches. For an example of a map file, see Section 5.1. For a list of /MAP options, refer to Section 3.2.2.

### 4.2.4 Symbol Files

The symbol file (or SYM file) is generated if you use the /SYFILE switch. This file contains all global symbols, module names, and module lengths, and, if you used the /LOCALS switch, all local symbols.

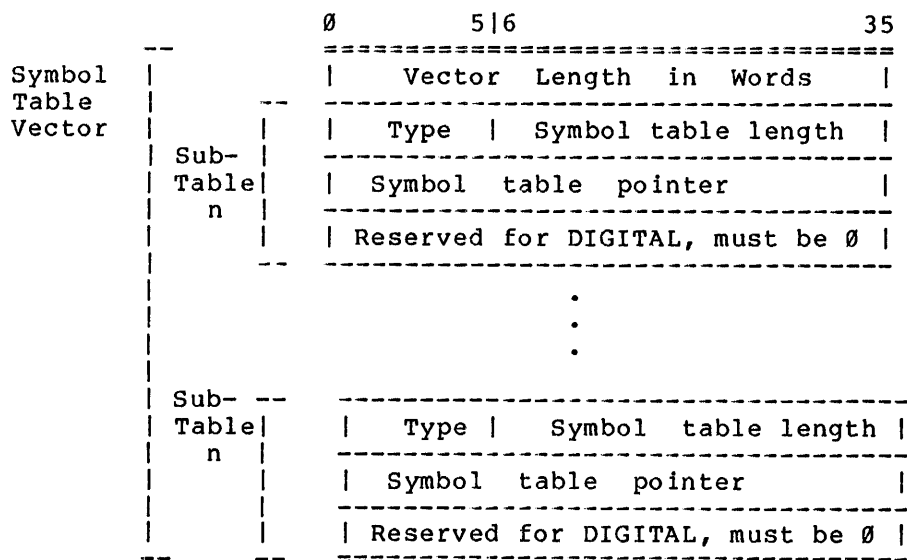
## 4.3 SYMBOL TABLE VECTOR

A symbol table vector is a pointer to the symbol tables of a program. There is one symbol table vector, and an undefined and defined symbol table per program.

When an extended symbol table is used, the contents of .JBUSY in JOBDAT are zeroed, and the address of the symbol table vector is loaded into .JBSYM. The symbol table vector contains two pointers: one to defined RADIX-50 symbols, the other to undefined RADIX-50 symbols.

The symbol table vector contains subtables that point to each symbol table and give their length and type. Each subtable is three words long, although only the first two words are currently used. The format of a symbol table vector is illustrated below.

OUTPUT FROM LINK



Symbol Table

Word	Symbol	Meaning															
0	.SYSTL	Defines the length in words of the symbol table vector including this word.															
0	.SYTYP	First subtable word, containing the following two fields:  SY.TYP is a 6-bit field that contains the symbol table type.  The types are:  <table border="0"> <thead> <tr> <th>Code</th> <th>Name</th> <th>Type</th> </tr> </thead> <tbody> <tr> <td>1</td> <td>.SYRD5</td> <td>Radix-50 defined symbols</td> </tr> <tr> <td>2</td> <td>.SYR5U</td> <td>Radix-50 undefined symbols</td> </tr> <tr> <td>3-37</td> <td></td> <td>Reserved for DIGITAL</td> </tr> <tr> <td>40-77</td> <td></td> <td>Reserved for customers</td> </tr> </tbody> </table>	Code	Name	Type	1	.SYRD5	Radix-50 defined symbols	2	.SYR5U	Radix-50 undefined symbols	3-37		Reserved for DIGITAL	40-77		Reserved for customers
Code	Name	Type															
1	.SYRD5	Radix-50 defined symbols															
2	.SYR5U	Radix-50 undefined symbols															
3-37		Reserved for DIGITAL															
40-77		Reserved for customers															
	SY.LEN	is a 30-bit field that contains the length in words of the particular symbol table.															
1	.SYADR	Lowest word in the table.  If bit 0 is 1, this word contains a section-local address. If bit 0 is 0, this word contains a global address. A section-local address is an 18-bit address. A global address is a 30-bit address.															

The Reserved word must be zero.

## OUTPUT FROM LINK

### 4.4 MESSAGES

During its processing, LINK issues messages about what it is doing, and about errors or possible errors it finds. LINK also responds to query switches such as /COUNTER, /ENTRY, /MISSING, /REQUEST, and /UNDEFINED.

Each LINK message has an assigned level and an assigned severity. (See Appendix B for the level and severity of each message.)

The level of a message determines whether it will be output to your terminal, the log file, or both. You can control this output by using the /ERRORLEVEL switch for the terminal and the /LOGLEVEL switch for the log file. LINK's defaults are /ERRORLEVEL:10 and /LOGLEVEL:10.

Responses to query switches and messages that require you to do something immediately are never output to the LOG file. For example, if you use the /UNDEFINE switch, LINK responds with the LNKUGS message; this message is output to the terminal but not to the log file.

The severity of a message determines whether LINK considers the message fatal (that is, whether the job is terminated). You can set the fatal severity with the /SEVERITY switch. The default severities are 24 for interactive jobs and 16 for batch jobs.

For both terminal messages and log file entries, LINK can issue short, medium, or long messages, depending on your use of the /VERBOSITY switch. For /VERBOSITY:SHORT, LINK gives only a 6-letter code; for /VERBOSITY:MEDIUM, LINK gives the code and a medium-length message; for /VERBOSITY:LONG, LINK gives the code, a medium-length message, and a long message.

Appendix B gives each 6-letter message code, its medium-length and long messages, and its level and severity.



## CHAPTER 5

### OVERLAYS

If your loaded program is too large to execute in one piece, you may be able to define an overlay structure for it. This permits the system to execute the program with only some parts at a time in your virtual address space. The overlay handler removes and reads in parts of the program, according to the overlay structure.

#### NOTE

You only need an overlay structure if your program is too large for your virtual address space. If the program can fit in your virtual space, you should not define an overlay structure for it; the monitor's page swapping facility is faster than overlay execution.

#### 5.1 OVERLAY STRUCTURES

An overlay program has a tree structure. (The tree is usually pictured upside down.) The tree is made up of links, each containing one or more program modules. These links are connected by paths. Using LINK switches, you define each link and each path.

At the top of the (upside down) tree is the root link, which must contain the main program. First-level links are below the root link; each first-level link is connected to the root link by one path.

Second-level links are below the first-level links, and each is connected by a path to exactly one first-level link. A link at level  $n$  is connected by a path to exactly one link at level  $n-1$ .

Notice that a link can have more than one downward path (to successor links), but only one upward path (to predecessor links).

Figure 5-1 shows a diagram of an overlay structure with 5 links. The root link is TEST; the first-level links are LEFT and RIGHT; the second-level links are LEFT1 and LEFT2.

## OVERLAYS

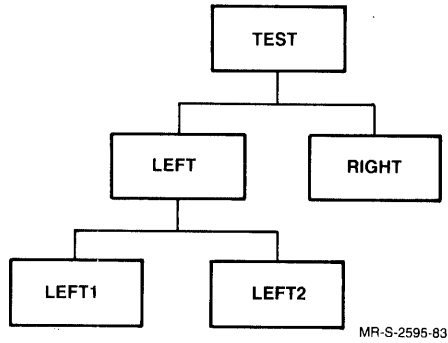


Figure 5-1 Example of an Overlay Structure

Defining an overlay structure allows your program to execute in a smaller space. This is because the code in a given link is allowed to make reference to memory only in links along a direct upward or downward path.

In the structure in Figure 5-1, the link LEFT can reference memory in itself, in the root link (TEST), or in its successor links LEFT1 and LEFT2. More generally, a link can reference memory in any link that is vertically connected to it.

Referencing memory in any other link is not allowed. For example, a path from LEFT1 to LEFT2 is not a direct upward or downward path.

Because of this restriction on memory references, only one complete vertical path (at most) is required in the virtual address space at any one time. The remaining links can be stored on disk while they are not needed.

### 5.1.1 Defining Overlay Structures

LINK has a family of overlay-related switches. These switches are described in detail in Section 3.2.2. The following example shows command strings for defining the overlay diagrammed in Figure 5-1. (Some of the command lines in this example are indented for clarity.)

```
*TEST/LOG/LOGLEVEL:2                ;Define TEST.LOG
*/ERRORLEVEL:5                       ;Important messages
*TEST/OVERLAY                         ;Define TEST.OVL
*TEST/MAP                             ;Define TEST.MAP
*LPT:TEST/PLOT                        ;Request diagram
*OVL0,OVL1/LINK:TEST                 ;Root link
*      /NODE:TEST OVL2/LINK:LEFT      ;Left branch
*      /NODE:LEFT OVL5/LINK:LEFT1     ;Left-left branch
*      /NODE:LEFT OVL6/LINK:LEFT2     ;Left-right branch
*      /NODE:TEST OVL3,OVL4/LINK:RIGHT ;Right branch
*TEXT/SSAVE                           ;Define TEST.EXE
*/EXECUTE/GO
```



## OVERLAYS

The first command string above defines the log file for the overlay. TEST/LOG specifies that the file is named TEST.LOG. The /LOGLEVEL:2 switch directs that messages of level 2 and above be entered in the log file.

In the second command string, the /ERRORLEVEL:5 switch directs that messages of level 5 and above be typed out on the terminal. The third command string, TEST/OVERLAY, tells LINK that an overlay structure is to be defined, and that the file for the overlay is to be TESTOVL.

The fourth command string, TEST/MAP, defines the file TEST.MAP, which will contain symbol maps for each link.

The next command string, LPT:TEST/PLOT, directs that a tree diagram of the overlay links be printed on the line printer.

The next command string, OVL0,OVL1/LINK:TEST, loads the files OVL0.REL and OVL1.REL into the root link. The /LINK:TEST switch tells LINK that no more modules are to be in the root link, and that the link name is TEST.

Each of the next four lines defines one link with a string of the form:

```
/NODE:linkname filename/LINK:linkname
```

The /NODE:linkname switch specifies the previously defined link to which the present link is an immediate successor. The filenames/LINK:linkname part of the line names the files containing modules to be included in the current link and specifies the name of the link.

The first of these four lines begins with /NODE:TEST, which tells LINK that the link being defined is to be an immediate successor to TEST, the root link. Then (on the same line), the string OVL2/LINK:LEFT loads the file OVL2.REL, ends the link, and names it LEFT.

The next line, /NODE:LEFT OVL5/LINK:LEFT1, defines a link named LEFT1 containing the file OVL5.REL, and this link is an immediate successor to the link LEFT.

The next line, /NODE:LEFT OVL6/LINK:LEFT2, defines another immediate successor to LEFT, this time containing the file OVL6.REL and called LEFT2.

The last link is defined in the next line, /NODE:TEST OVL3,OVL4/LINK:RIGHT. This string defines the link RIGHT, which is an immediate successor to TEST and contains the files OVL3.REL and OVL4.REL.

The next-to-last line, TEST/SSAVE, directs LINK to create the saved file TEST.EXE. The last line, /EXECUTE/GO, specifies that the loaded program is to be executed, and that all commands to LINK are completed.

The process also produced an executable file TEST.EXE, which can be run using the RUN system command. However, to run the program, the file TEST.OVL must be present, because it provides the code for the links.

## OVERLAYS

### 5.1.2 An Overlay Example

The following pages show terminal listings of the files associated with the example above. These pages are:

1. Terminal copy of the FORTRAN source files used in the overlay.
2. Terminal copy of the compilation of the source files.
3. Terminal copy of the interactive use of LINK to define and execute the overlay.
4. The file TEST.LOG generated by LINK, which shows the log messages issued during the load.
5. The file TEST.MAP generated by LINK, which shows symbol maps for the overlay.
6. The tree diagram requested by the LPT:/PLOT switch.

## OVERLAYS

```
.type ov10.for
  TYPE 1
  FORMAT('1','Execution begins in main Program OVL0')
  TYPE 11
  FORMAT(1X,'OVL0 calls OVL2A')
  CALL OVL2A
  TYPE 2
  FORMAT(/1X,'Return to OVL0')
  TYPE 21
  FORMAT(1X,'OVL0 calls OVL4')
  CALL OVL4
  TYPE 2
  TYPE 3
  FORMAT(/1X,'Execution ends in main Program OVL0')
  STOP
  END

.type ov11.for
  SUBROUTINE OVL1
  TYPE 1
  FORMAT(/1X,' OVL1 calls OVL3')
  CALL OVL3
  TYPE 2
  FORMAT(/1X,' Return to OVL1')
  RETURN
  END

.type ov12.for
  SUBROUTINE OVL2A
  TYPE 1
  FORMAT(/1X,' OVL2A calls OVL5')
  CALL OVL5
  TYPE 2
  FORMAT(/1X,' Return to OVL2A')
  TYPE 3
  FORMAT(1X,' OVL2A calls OVL6')
  CALL OVL6
  TYPE 2
  RETURN
  END
  SUBROUTINE OVL2B
  TYPE 1
  FORMAT(/1X,' OVL2B doesn't call anything')
  RETURN
  END

.type ov13.for
  SUBROUTINE OVL3
  TYPE 1
  FORMAT(/1X,' OVL3 doesn't call anything')
  RETURN
  END

.type ov14.for
  SUBROUTINE OVL4
  TYPE 1
  FORMAT(/1X,' OVL4 calls OVL1')
  CALL OVL1
  TYPE 2
  FORMAT(/1X,' Return to OVL4')
  RETURN
  END

.type ov15.for
  SUBROUTINE OVL5
  TYPE 1
  FORMAT(/1X,' OVL5 doesn't call anything')
  RETURN
  END

.type ov16.for
  SUBROUTINE OVL6
  TYPE 1
  FORMAT(/1X,' OVL6 calls OVL2B')
  CALL OVL2B
  TYPE 2
  FORMAT(/1X,' Return to OVL6')
  RETURN
  END
```

## OVERLAYS

```
.COMPILE OVLO,OVL1,OVL2,OVL3,OVL4,OVL5,OVL6
FORTRAN: OVLO
OVL0
FORTRAN: OVL1
OVL1
FORTRAN: OVL2
OVL2A
OVL2B
FORTRAN: OVL3
OVL3
FORTRAN: OVL4
OVL4
FORTRAN: OVL5
OVL5
FORTRAN: OVL6
OVL6

.R LINK
*TEST/LOG/LOGLEVEL:5
*/ERRORLEVEL:5/NOINITIAL
*TEST/OVERLAY
*TEST/MAP
*DSK:TEST/PLOT/PLTTY:PRINTER
*OVLO,OVL1/LINK:TEST
[LNKLMN Loading module OVLO from file DSK:OVLO.REL[10,3551,LINK5A]]
[LNKLMN Loading module OVL1 from file DSK:OVL1.REL[10,3551,LINK5A]]
[LNKLMN Loading module OVRLAY from file SYS:OVRLAY.REL[1,5]]
[LNKLMN Loading module JOBDAT from file SYS:JOBDAT.REL[1,4]]
[LNKLMN Loading module FORINI from file SYS:FORLIB.REL[1,5]]
[LNKLMN Loading module FORDST from file SYS:FORLIB.REL[1,5]]
[LNKLMN Loading module FORPSE from file SYS:FORLIB.REL[1,5]]
[LNKELN End of link number 0 name TEST]
*/NODE:TEST OVL2/LINK:LEFT
[LNKLMN Loading module OVL2A from file DSK:OVL2.REL[10,3551,LINK5A]]
[LNKLMN Loading module OVL2B from file DSK:OVL2.REL[10,3551,LINK5A]]
[LNKELN End of link number 1 name LEFT]
*/NODE:LEFT OVL5/LINK:LEFT1
[LNKLMN Loading module OVL5 from file DSK:OVL5.REL[10,3551,LINK5A]]
[LNKELN End of link number 2 name LEFT1]
*/NODE:LEFT OVL6 /LINK:LEFT2
[LNKLMN Loading module OVL6 from file DSK:OVL6.REL[10,3551,LINK5A]]
[LNKELN End of link number 3 name LEFT2]
*/NODE:TEST OVL3,OVL4/LINK:RIGHT
[LNKLMN Loading module OVL3 from file DSK:OVL3.REL[10,3551,LINK5A]]
[LNKLMN Loading module OVL4 from file DSK:OVL4.REL[10,3551,LINK5A]]
[LNKELN End of link number 4 name RIGHT]
*TEST/SSAVE
*/EXECUTE/GO
[LNKXCT OVLO execution]

Execution begins in main Program OVLO
OVLO CALLS OVL2A
OVL2A CALLS OVL5

                                OVL5 DOESN'T CALL ANYTHING

RETURN TO OVL2A
OVL2A CALLS OVL6

                                OVL6 CALLS OVL2B
                                OVL2B DOESN'T CALL ANYTHING
                                RETURN TO OVL6

RETURN TO OVL2A

RETURN TO OVLO
OVLO CALLS OVL4

                                OVL4 CALLS OVL1
                                OVL1 CALLS OVL3
                                OVL3 DOESN'T CALL ANYTHING

                                RETURN TO OVL1
                                RETURN TO OVL4

RETURN TO OVLO

Execution ends in main Program OVLO

CPU time 0.32   Elapsed time 0.95

EXIT
```

# OVERLAYS

.TYPE TEST.LOG

```

12:55:15 6 1 LMN Loading module OVL0 from file DSK:OVL0.REL[10,3551,LINK5A]
12:55:15 6 1 LMN Loading module OVL1 from file DSK:OVL1.REL[10,3551,LINK5A]
12:55:16 6 1 LMN Loading module OVL2 from file DSK:OVL2.REL[10,3551,LINK5A]
12:55:16 6 1 LMN Loading module OVL3 from file DSK:OVL3.REL[10,3551,LINK5A]
12:55:16 6 1 LMN Loading module OVL4 from file DSK:OVL4.REL[10,3551,LINK5A]
12:55:16 6 1 LMN Loading module OVL5 from file DSK:OVL5.REL[10,3551,LINK5A]
12:55:16 6 1 LMN Loading module OVL6 from file DSK:OVL6.REL[10,3551,LINK5A]
12:55:19 7 1 ELN End of link number 0 name TEST

12:55:19 6 1 LMN Loading module OVL2A from file DSK:OVL2A.REL[10,3551,LINK5A]
12:55:19 6 1 LMN Loading module OVL2B from file DSK:OVL2B.REL[10,3551,LINK5A]
12:55:20 7 1 ELN End of link number 1 name LEFT

12:55:20 6 1 LMN Loading module OVL5 from file DSK:OVL5.REL[10,3551,LINK5A]
12:55:20 7 1 ELN End of link number 2 name LEFT1

12:55:20 6 1 LMN Loading module OVL6 from file DSK:OVL6.REL[10,3551,LINK5A]
12:55:20 7 1 ELN End of link number 3 name LEFT2

12:55:21 6 1 LMN Loading module OVL3 from file DSK:OVL3.REL[10,3551,LINK5A]
12:55:21 6 1 LMN Loading module OVL4 from file DSK:OVL4.REL[10,3551,LINK5A]
12:55:21 7 1 ELN End of link number 4 name RIGHT
  
```

.TYPE TEST.MAP

LINK symbol map of TEST /KL/KS Page 1  
Produced by LINK version 5A(2030) on 17-Dec-82 at 12:55:23

```

Overlay no.      0      name TEST
Overlay is absolute
Low segment starts at      0 ends at 10024 length 10025 = 9P
Control Block address is 7763, length 32 (octal), 26. (decimal)
491 words free in Low segment
95 Global symbols loaded, therefore min. hash size is 106
Start address is 235, located in program OVL0
  
```

\*\*\*\*\*

```

OVL0 from DSK:OVL0.REL[10,3551,LINK5A] created by FORTRAN /KL/KS on 4-Nov-82 at 15:00:00
Low segment starts at 140 ends at 234 length 75 (octal), 61. (decimal)
High segment starts at 235 ends at 345 length 111 (octal), 73. (decimal)

MAIN. 235 Global Relocatable OVL0 235 Entry Relocatable
  
```

\*\*\*\*\*

```

OVL1 from DSK:OVL1.REL[10,3551,LINK5A] created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00
Low segment starts at 346 ends at 377 length 32 (octal), 26. (decimal)
High segment starts at 400 ends at 442 length 43 (octal), 35. (decimal)

OVL1 401 Entry Relocatable FORDT% 500010 Global Absolute
  
```

\*\*\*\*\*

```

OVL2 from SYS:OVL2.REL[1,5] created by MACRO on 2-Aug-82 9:18:00
Low segment starts at 5223 ends at 6226 length 1004 (octal), 516. (decimal)
High segment starts at 443 ends at 5107 length 4445 (octal), 2341. (decimal)

CLROV. 2712 Entry Relocatable GETOV. 2610 Entry Relocatable
INIOV. 2570 Entry Relocatable LOGOV. 2731 Entry Relocatable
REMOV. 2624 Entry Relocatable RUNOV. 2642 Entry Relocatable
SAVOV. 2665 Entry Relocatable %OVL2 50100204 Global Absolute Suppressed
,OVL2A 5224 Entry Relocatable ,OVL2B 5302 Global Relocatable
,OVL2B 3541 Entry Relocatable ,OVL2A 5301 Global Relocatable
  
```

\*\*\*\*\*

```

JOBDAT from SYS:JOB DAT.REL[1,4] created by MACRO on 26-FEB-81 at 19:25:00
Zero length module
  
```

\*\*\*\*\*

# OVERLAYS

FORINI from SYS:FORLIB.REL[1,5] created by MACRO on 15-Sep-82 at 19:45:00  
 Low segment starts at 6437 ends at 7476 length 1040 (octal), 544. (decimal)  
 High segment starts at 6227 ends at 6436 length 210 (octal), 136. (decimal)

ABORT,	6432	Entry	Relocatable	ALCHN,	6406	Entry	Relocatable
ALCOR,	6402	Entry	Relocatable	CERPT,	6437	Global	Relocatable
CLOSE,	6344	Entry	Relocatable	DBMS,	6416	Entry	Relocatable
DEC,	6362	Entry	Relocatable	DECHN,	6410	Entry	Relocatable

LINK symbol map of TEST /KL/KS Page 2

FORINI	DECOR,	6404	Entry	Relocatable	ENC,	6360	Entry	Relocatable
	EXIT,	6400	Entry	Relocatable	EXIT1,	6346	Entry	Relocatable
	FIN,	6372	Entry	Relocatable	FIND,	6376	Entry	Relocatable
	FDRER,	6340	Entry	Relocatable	FOROP,	6422	Entry	Relocatable
	FUNCT,	6414	Entry	Relocatable	IFI,	6424	Entry	Relocatable
	IFO,	6426	Entry	Relocatable	IN,	6350	Entry	Relocatable
	INIT,	6336	Entry	Relocatable	INQ,	6420	Entry	Relocatable
	IDLST,	6370	Entry	Relocatable	MTHER,	6430	Entry	Relocatable
	MTOP,	6374	Entry	Relocatable	NLI,	6364	Entry	Relocatable
	NLO,	6366	Entry	Relocatable	OPEN,	6342	Entry	Relocatable
	OUT,	6352	Entry	Relocatable	RESET,	6227	Entry	Relocatable
	RTB,	6354	Entry	Relocatable	TRACE,	6412	Entry	Relocatable
	WTB,	6356	Entry	Relocatable				

\*\*\*\*\*

FORDST from SYS:FORLIB.REL[1,5] created by MACRO on 15-Sep-82 at 19:45:00  
 High segment starts at 7477 ends at 7477 length 1 (octal), 1. (decimal)

DBSTP*	7477	Entry	Relocatable
--------	------	-------	-------------

\*\*\*\*\*

FORPSE from SYS:FORLIB.REL[1,5] created by MACRO on 15-Sep-82 at 19:45:00  
 Low segment starts at 7702 ends at 7762 length 61 (octal), 49. (decimal)  
 High segment starts at 7500 ends at 7701 length 202 (octal), 130. (decimal)

PAUS,	7501	Entry	Relocatable	STOP,	7504	Entry	Relocatable
-------	------	-------	-------------	-------	------	-------	-------------

\*\*\*\*\*

Index to LINK symbol map of off%e TEST /KL/KS Page 3

Name	Page	Name	Page	Name	Page	Name	Page
FORDST	2	FORPSE	2	OVL0	1	OVLAY	1
FORINI	1	JOB DAT	1	OVL1	1		

LINK symbol map of TEST /KL/KS #1 Page 4

Overlay no. 1 name LEFT  
 Overlay is absolute  
 Low segment starts at 14025 ends at 14262 length 236 = 1P  
 Control Block address is 14221, length 30 (octal), 24. (decimal)  
 Path is 0  
 333 words free in Low segment  
 6 Global symbols loaded, therefore min. hash size is 7

\*\*\*\*\*

OVL2A from DSK:OVL2.REL[10,3551,LINK5A] created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00  
 Low segment starts at 14025 ends at 14076 length 52 (octal), 42. (decimal)  
 High segment starts at 14077 ends at 14173 length 75 (octal), 61. (decimal)

OVL2A	14100	Entry	Relocatable
-------	-------	-------	-------------

\*\*\*\*\*

OVL2B from DSK:OVL2.REL[10,3551,LINK5A] created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00  
 Low segment starts at 14174 ends at 14206 length 13 (octal), 11. (decimal)  
 High segment starts at 14207 ends at 14220 length 3 12 (octal), 10. (decimal)

OVL2B	14210	Entry	Relocatable
-------	-------	-------	-------------

\*\*\*\*\*

# OVERLAYS

LINK symbol map of TEST /KL/KS #2 Page 5

```

Overlay no.      2      name  LEFT1
Overlay is absolute
Low segment starts at 14263 ends at 14346 length      64 = 1P
Control Block address is 14137, length      20 (octal), 16. (decimal)
Path is 0,1
281 words free in Low segment
3 Global symbols loaded, therefore min. hash size is 4
    
```

\*\*\*\*\*

```

OVL5  from DSK:OVL5.REL[10,3551,LINK5A]      created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00
Low segment starts at 14263 ends at 14276 length      14 (octal), 12. (decimal)
High segment starts at 14277 ends at 14316 length      20 (octal), 16. (decimal)
    
```

```

OVL5      14300      Entry  Relocatable
    
```

\*\*\*\*\*

LINK symbol map of TEST /KL/KS #3 Page 6

```

Overlay no.      3      name  LEFT2
Overlay is absolute
Low segment starts at 14263 ends at 14447 length      165 = 1P
Control Block address is 14420, length      20 (octal), 16. (decimal)
Path is 0,1
216 words free in Low segment
4 Global symbols loaded, therefore min. hash size is 5
    
```

\*\*\*\*\*

```

OVL6  from DSK:OVL6.REL[10,3551,LINK5A]      created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00
Low segment starts at 14263 ends at 14331 length      47 (octal), 39. (decimal)
High segment starts at 14332 ends at 14417 length      66 (octal), 54. (decimal)
    
```

```

OVL6      14333      Entry  Relocatable
    
```

\*\*\*\*\*

LINK symbol map of TEST /KL/KS #4 Page 7

```

Overlay no.      4      name  RIGHT
Overlay is absolute
Low segment starts at 14025 ends at 14207 length      163 = 1P
Control Block address is 14156, length      22 (octal), 16. (decimal)
Path is 0
376 words free in Low segment
5 Global symbols loaded, therefore min. hash size is 6
    
```

\*\*\*\*\*

```

OVL3  from DSK:OVL3.REL[10,3551,LINK5A]      created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00
Low segment starts at 14025 ends at 14040 length      14 (octal), 12. (decimal)
High segment starts at 14041 ends at 14060 length      20 (octal), 16. (decimal)
    
```

```

OVL3      14042      Entry  Relocatable
    
```

\*\*\*\*\*

```

OVL4  from DSK:OVL4.REL[10,3551,LINK5A]      created by FORTRAN /KL/KS on 4-Nov-82 at 14:47:00
Low segment starts at 14061 ends at 14112 length      32 (octal), 26. (decimal)
High segment starts at 14113 ends at 14155 length      43 (octal), 35. (decimal)
OVL4      14114      Entry  Relocatable
    
```

\*\*\*\*\*

Index to overlay number of TEXT /KL/KS Page 8

Overlay Page	Overlay Page	Overlay Page	Overlay Page
#0	3	#2	5
#1	4	#3	6
		#4	7

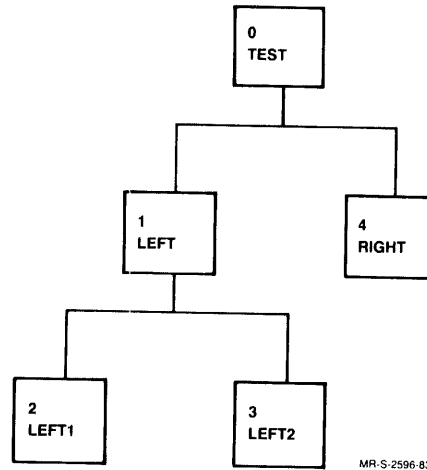
Index to overlay names of TEST

Name	Page	Name	Page	Name	Page	Name	Page
LEFT	4	LEFT2	6	RIGHT	7	TEST	3
LEFT1	5						

[End of LINK map of TEST]

## OVERLAYS

The listing file TEST.OVL will look similar to the following:



### 5.2 WRITABLE OVERLAYS

Ordinarily each overlay link built by LINK is copied by the overlay handler from the OVL file to the address space at runtime. The contents of any locations that have been modified will be lost each time the overlay link is copied from the OVL file. This can be prevented by the use of writable overlays.

If a link is specified as writable, the overlay handler copies that link to a temporary file on disk before overwriting it. Later, when the copied link is needed, the overlay handler retrieves the link from the temporary file rather than the OVL file. In this way, any modified values are preserved. Because writable overlays involve more file I/O, they are slower than the default (nonwritable) overlays and should only be used when the program structure and storage requirements demand dynamic storage in overlay links.

To specify that an overlay is writable, use the FORTRAN SAVE statement in the program, and specify /OVERLAY:WRITABLE when loading the program with LINK.

#### 5.2.1 Writable Overlay Syntax

To build a writable overlay, specify the keyword WRITABLE with the /OVERLAY switch in the LINK command line:

```
filespec/OVERLAY:WRITABLE
```

#### 5.2.2 Writable Overlay Error Messages

The overlay handler must write and update a temporary file. In addition to the error messages associated with all overlays, there are two additional error messages for writable overlays:

```
? OVLCWF Cannot write file [filename]: [reason]
```

```
? OVLCUF Cannot update file [filename]: [reason]
```

If either of these messages appears, you should check for disk quota violations or other conditions that could prevent the overlay handler from writing a temporary file.



## OVERLAYS

### 5.3 RELOCATABLE OVERLAYS

LINK ordinarily allocates 2000 extra words at the end of the root link and no extra space at the end of each subsequent link. This is adequate for programs with static storage requirements. If a link requires extra storage at run-time, you can use the /SPACE switch to make the necessary allowances for the program's requirements. The /SPACE switch allows you to specify the number of words to be allocated after the current link is loaded.

However, there are programs whose dynamic run-time storage requirements are unpredictable. For example, a program's run-time storage requirements may vary according to the program's input. For this class of programs, relocatable overlays can be useful.

For relocatable overlays LINK places extra relocation information in the OVL file, permitting overlay links to be relocated at runtime. The overlay handler, using the FUNCT. subroutine, can determine where the link will fit in the address space and resolve relocatable addresses within the link. This extra processing causes relocatable overlays to run slower than nonrelocatable overlays. Relocatable overlays should only be used when you cannot determine the dynamic storage requirements of a program.

#### 5.3.1 Relocatable Overlay Syntax

To build a relocatable overlay, specify the RELOCATABLE keyword to the /OVERLAY switch in the LINK command line:

```
filespec/OVERLAY:RELOCATABLE
```

#### 5.3.2 Relocatable Overlay Messages

If /OVERLAY:(LOGFILE,RELOCATABLE) is specified during the loading of a program, informational messages of the following form are sent to the user's terminal:

```
%OVLRLR Relocating link [linkname] at [address]
```

## OVERLAYS

### 5.4 RESTRICTIONS ON OVERLAYS

The following restrictions apply to all overlaid programs:

- Overlaid programs cannot be run execute-only.
- PSECTed programs cannot be overlaid.
- Overlaid programs with large buffer requirements must use the /SPACE switch. If an %OVLMAN (Memory not available) error is encountered, the program should be reloaded using the /SPACE switch with each link.
- If the program uses more than 256 links, use the /MAXNODE switch to specify the number of links necessary for the program. LINK will allocate extra space in the the OVL file for tables that require it, based on the number of links you specify.

#### 5.4.1 Restrictions on Absolute Overlays

The following restrictions apply to absolute overlaid programs:

1. Any intermediate results stored in non-root links are lost as soon as the links are overlaid. Do not expect to retain a value stored in a non-root link unless /OVERLAY:WRITABLE has been specified.
2. Certain forms of global, inter-overlay references are not recommended because you cannot be sure that the necessary modules will be in memory at the right time. Some of these references are:
  - Additive fixups, in the form FOO##+BAR where FOO is in another overlay.
  - Left-hand fixups, in the form XWD FOO##,BAR, where FOO is in another overlay.
  - Fullword fixups, in the form EXP FOO##, where FOO is in another overlay.
  - Similarly, MOVEI 1,FOO##, where FOO is in a different overlay, should not be used, because the necessary module may not be in memory.

In fact, the only predictable inter-overlay global reference is one that brings the necessary module into memory, such as PUSHJ P,FOO##.

## OVERLAYS

### 5.4.2 Restrictions on Relocatable Overlays

The following restriction applies to relocatable overlays:

- Complex expressions involving relocatable symbols are not relocated properly in a relocatable overlay. No standard DEC compiler produces such expressions. MACRO programmers should avoid using them in subroutines that are to be loaded as part of an overlaid program. Any expression that causes MACRO to generate a Polish fixup block will not be properly relocated at runtime. The following are examples of such complex expressions: such a complex expression:

```
MOVEI 1,A## + B## + C##  
A,,0
```

### 5.4.3 Restrictions on FORTRAN Overlays

The following restriction applies to FORTRAN programs that are written with associate variables and using the overlay facility.

- If the associate variable is declared in a subroutine, that subroutine must be loaded in the root link of the overlay structure. Accessing a file opened with an associate variable changes the value of the specified variable. If this variable is in a nonresident overlay link when the access is made, program execution will produce unpredictable results. Moreover, the value of the variable will be reset to zero each time its overlay link is removed from memory. Only variables declared in routines that are loaded into the root link will always be resident. However, variables declared in COMMON and in the root link will always be resident, and may be safely used as associate variables.
- If you place COMMON in a writable overlay, be sure that all references to the variables in that COMMON are in the same overlay or its successors.
- A FORTRAN ASSIGN statement may be used in a relocatable overlay. If the ASSIGN is made in a subroutine, the value of the assigned variable may be preserved from one call of that subroutine to the next. However, the overlay containing that subroutine could then be replaced in memory by a different overlay. If the overlay containing the subroutine is relocated differently when brought back into memory, any subsequent GOTO may fail.

## OVERLAYS

### 5.5 SIZE OF OVERLAY PROGRAMS

Although most programs have a consistent size, the size of an overlay program depends on which overlays are in memory. This can be ascertained by using the /COUNTER switch when linking the program. To do this, place /COUNTER after the /LINK switch for the overlay of which you want to know the size, but before the next /NODE switch. This will give you the size of the program when the overlay is actually loaded into memory. The display will include all routines loaded from the runtime libraries. This allows you to determine which overlay is the largest, and whether the program can be loaded without restructuring.

### 5.6 DEBUGGING OVERLAYED PROGRAMS

COBDDT and ALGDDT can be used to debug overlay programs, but FORDDT cannot. To use DDT with an overlaid program, the program should be loaded using /SYMSEG:LOW, with local symbols for the desired modules.

To set breakpoints in an overlay, put a subroutine in the root node, and call the subroutine from the overlay. Such a subroutine need consist only of a SUBROUTINE statement, a RETURN, and an END. The breakpoint can be set at this subroutine before the program starts running.

When a FORTRAN program starts running, it calls RESET. in FOROTS, which removes the symbol table. The symbol table will return after the first overlay is called. If you need the symbols for debugging the root link, insert a CALL INIOVL at the beginning of the main program (refer to Section 5.7.1 for more information). This call will reinstall the symbol table. LINK builds a separate symbol table for each overlay, so that all the symbols known to DDT are for modules that are currently in memory. Note that it is not possible to single-step through RESET. (\$X and \$\$X will not work). Set a breakpoint after RESET. if you are debugging a root link, and use \$G.

### 5.7 THE OVERLAY HANDLER

LINK's overlay handler is the program that supervises execution of overlay structures defined by LINK switches.

When you load an overlay structure, the overlay handler is loaded into the root link of the structure. From there it can supervise overlaying operations, because the root link is always in your virtual address space during execution. During execution, when a link not in memory is called, the overlay handler brings in the link, possibly overlaying one or more links already in memory.

## OVERLAYS

The overlay handler consists of self-modifying code and data, and two 128-word buffers. One of these buffers, IDXBFR, contains a 128-word section of the link number index table. This allows 256 links to be directly referenced at any one time. The second buffer, INBFR, contains the preambles and relocation tables, if required, of the individual links.

There are two ways of overlaying links during execution:

1. A call to a link not in memory implicitly calls the overlay handler to overlay one or more links with the required links. This action of the overlay handler is transparent to the user.
2. An explicit call to one of several entry points in the overlay handler can cause one or more links to be overlaid. These entry points and calls to them are discussed in the sections below.

### 5.7.1 Calls to the Overlay Handler

Overlays can be used transparently, or they can be explicitly called from the program. Such calls are made to one of the entry points in the overlay handler.

The overlay handler has five entry points that are available for calls from user programs. To call the overlay handler from a MACRO program, you must use the standard calling sequence, which is:

```
MOVEI    16,arglst
PUSHJ    17,entry-name
```

Where arglst is the address of the first argument in the argument list, and entry-name is the entry-point name.

The argument list must be of the form:

```
arglst:  -n,,0          ;n is number of arguments
          Z code,addr1  ;For first argument
          .
          .
          Z code,addrn  ;For nth argument
```

Where addr... is the address of the argument.

The legal values of "code" are 2 (for a link number), 17 (for an ASCII string), and 15 (for a character string descriptor).

For each word of the argument list, the code indicates the type of argument. The code occupies the AC field, bits 9 through 12. The address gives the location of the argument; it can be indirect and indexed.

To call the overlay handler from a FORTRAN program, the call must be of the form:

```
CALL subroutine (arglst)
```

Where subroutine is the name of the desired subroutine, and arglst is a list of arguments separated by commas.

## OVERLAYS

### 5.7.2 Overlay Handler Subroutines

Each of the seven callable subroutines in the overlay handler has an entry name symbol for use with MACRO, and a subroutine name for use with FORTRAN, as follows:

MACRO Entry Name Symbol	FORTRAN Subroutine	Subroutine Function
CLROV.	CLROVL	Specifies a non-writable overlay.
GETOV.	GETOVL	Brings specified links into memory.
INIOV.	INIOVL	Specifies the file from which the overlay program will be read, if the load time specification is to be overridden.
LOGOV.	LOGOVL	Specifies or closes the file in which runtime messages from the overlay handler will be written.
REMOV.	REMOVL	Removes specified links from memory.
RUNOV.	RUNOVL	Moves into memory a specified link and begins execution at its start address.
SAVOV.	SAVOVL	Specifies a writable overlay.

#### Declaring a Non-Writable Link (CLROV.)

You can declare an overlay link to be non-writable, using the CLROV. entry point. This does not immediately affect the program, but waits until the link is about to be overlaid or read in. If the link is already non-writable, this entry point has no effect.

#### Example

```
                MOVEI      16,arglst
                PUSHJ     17,CLROV.

arglst:  -n,,0           ;n is number of arguments
          Z 17,addr1     ;for first ASCIZ linkname
          .
          .
          Z 17,addrn     ;for nth ASCIZ linkname

                                OR

arglst:  -n,,0           ;n is number of arguments
          Z 2,addr1      ;for first link number
          .
          .
          Z 2,addrn      ;for nth link number
```

Where addr... is the address of the argument.

## OVERLAYS

### Getting a Specific Path (GETOV.)

The subroutine to bring a specific path into core can be used to make sure that a particular path is used when otherwise the overlay handler might have a choice of paths. It is illegal to specify a path that overlays the calling link.

To call the subroutine from a FORTRAN program, use:

```
CALL GETOVL (linkname,...,linkname)
```

where each linkname is the ASCII name of a link in the desired path.

To call the subroutine from a MACRO program, use the standard FORTRAN calling sequence:

```
MOVEI    16,arglst
PUSHJ    17,GETOV.
```

The argument list has one word for each link required to be in the path.

#### Example

```
arglst:  -n,,0                ;n is number of arguments
          Z 17,addr1
          .
          .
          Z 17,addrn
```

OR

```
arglst:  -n,,0                ;n is number of arguments
          Z 2,addr1
          .
          .
          Z 2,addrn
```

Where addr... is the address of the argument.

### Initializing an Overlay (INIOV.)

The overlay initializing subroutine specifies a file from which the overlay program will be read. This subroutine is used to override the file specified at load time. The file specified to INIOV. can have any valid specification, but it must be in the correct format for an overlay (OVL) file.

To call the subroutine from a FORTRAN program, use:

```
CALL INIOVL ('filespec')
```

where 'filespec' is a literal constant that can give a device, a filename, a file type, and a project-programmer number (PPN).

To call the subroutine from a MACRO program, use the standard FORTRAN calling sequence:

```
MOVEI    16,arglst
PUSHJ    17,INIOV.
```

## OVERLAYS

The argument list is of the form:

```
      -1,,0  
arglst: Z 17,address of ASCIZ filespec
```

where filespec is an ASCIZ string (ASCII ending with nulls) that can give a device, a filename, a file type, and a PPN

### NOTE

If you call INIOV. with no arguments, it initiates the overlay handler and reads in the symbols for the root link, using the overlay file specified at load time. This can be useful for debugging the root link before any successor links have been read in, because symbols are not normally available until the first link comes into memory.

### Specifying an Overlay Log File (LOGOV.)

You can specify an output file for runtime messages from the overlay handler. These messages are listed in Section 5.5. The log file entry includes the elapsed run time since the first call to the overlay handler.

To call this subroutine from a FORTRAN program, use:

```
CALL LOGOVL ('filespec')
```

where 'filespec' is a literal constant that can give a device, a filename, a file type, and a PPN.

To close the file, use

```
CALL LOGOVL (0)
```

To call the subroutine from a MACRO program, use the standard FORTRAN calling sequence:

```
MOVEI    16,arglst  
PUSHJ    17,LOGOV.
```

The argument list is of the form:

```
      -1,,0  
arglst: Z 17,address of ASCIZ filespec
```

Where filespec is an ASCIZ string that can give a device, a filename, a file type, and a PPN.

To close the log file, the argument list is:

```
      -1,,0  
arglst: Z 17,address of word containing zero
```



## OVERLAYS

### Removing Specific Links from Memory (REMOV.)

The subroutine to remove specific links from memory, once they are no longer required, can be used to reduce core image size for faster execution. Specifying removal of the calling link causes an error.

To call the subroutine from a FORTRAN program, use:

```
CALL REMOVEL (linkname,...,linkname)
```

Where each linkname is the ASCIZ name of a link to be removed from memory.

To call the subroutine from a MACRO program, use the standard FORTRAN calling sequence:

```
MOVEI    16,arglst
PUSHJ    17,REMOV.
```

The argument list has one word for each link to be removed.

#### Example

```
arglst:  -n,,0                ;n is number of arguments
          Z 17,addrl
          .
          .
          Z 17,addrn
```

OR

```
arglst:  -n,,0                ;n is number of arguments
          Z 2,addl
          .
          .
          Z 2,addrn
```

Where addr... is the address of the argument.

### Running a Specific Link (RUNOV.)

The subroutine for running a specific link allows you to transfer program execution to the start address of a particular link. (An error occurs if the link has no start address.) If the link is not already in memory, it and its path are brought in.

You can use this subroutine to overlay the calling link, because the next instruction executed is the start address of the named link; therefore, there is no automatic return to the calling link.

#### NOTE

The FORTRAN compiler does not generate start addresses for subroutines. FORTRAN main programs cannot be loaded into non-root links. Therefore, to use RUNOVL to transfer control to a FORTRAN subroutine in a non-root link, you must use the /START switch at load time to define a start address for the link.

## OVERLAYS

To call the subroutine RUNOVL from a FORTRAN program, use:

```
CALL RUNOVL (linkname)
```

Where linkname is the ASCIZ name of the link to be run.

To call the subroutine from a MACRO program, use the standard FORTRAN calling sequence:

```
MOVEI    16, arglst
PUSHJ    17, RUNOV.
```

The argument list is of the form:

```
      -1,,0
arglst: Z 17, address of ASCIZ linkname
```

OR

```
      -1,,0
arglst: Z 2, address of link number
```

### Declaring A Writable Link (SAVOV.)

You can dynamically declare an overlay link to be writable by calling SAVOV. This does not affect the current state of the code immediately, but waits until the link is about to be overlaid. If the link already writable, this symbol has no effect.

#### Example

```
MOVEI    16, arglst
PUSHJ    17, SAVOV.

      -n,,0                                ;n is number of arguments
arglst:  Z 17, addr1                        ;for first ASCIZ linkname
      .
      .
      Z 17, addrn                          ;for nth ASCIZ linkname
```

OR

```
arglist: -n,,0                              ;n is number of arguments
      Z 2, addr1                            ;for first link number
      .
      .
      Z 2, addrn                            ;for nth link number
```

Where addr... is the address of the argument.

If called with no arguments, SAVOV. only initializes the temporary file.

## OVERLAYS

### 5.7.3 Overlay Handler Messages

This section lists all of the overlay handler's messages. (The messages from LINK, which have the LNK prefix, are given in Appendix B.)

For each overlay handler message, the last three letters of the six-letter code, the severity, and the text of the message are given in boldface. Then, in lightface type, comes an explanation of the message.

When a message is issued, the three letters are suffixed to the letters OVL, forming a 6-letter code of the form OVLxxx. The explanation of the message will be printed only if you use the /OVERLAY:LOG switch.

The severity of a message determines whether the job will be terminated when the message is issued. Level 31 messages terminate program execution. Level 8 messages are warnings: they do not terminate execution, but the error may affect the execution of the program. Level 1 messages are informational and are printed on the terminal only if you specified /OVERLAY:LOGFILE.

Code	Sev	Message and Explanation
ARC	31	<b>Attempt to remove caller from link [name or number]</b>  The named link attempted to remove the link that called it. This error occurs when the call to the REMOV. subroutine requests removal of the calling link.
ARL	8	<b>Ambiguous request in link number [number] for [symbol], using link number [number]</b>  More than one successor link satisfies a call from a predecessor link, and none of these successors is in memory. Since all their paths are of equal length, the overlay handler has selected an arbitrary link.
CCF	31	<b>Cannot close file [file], status [octal]</b>  For some reason, the overlay handler cannot close one of its working files. This is a file I/O error.
CDL	31	<b>Cannot delete link [name or number], FUNCT. return status [number]</b>  This is an internal LINK error, and is not expected to occur. If it does, please notify your Software Specialist, or send a Software Performance Report (SPR) to DIGITAL.  Return status is one of the following:  1 Core already deallocated 3 Illegal argument passed to FUNCT. module
CFE	31	<b>Cannot find file [file] [reason]</b>  The overlay handler has attempted, unsuccessfully, to open an EXE, OVL, or .TMP file.

## OVERLAYS

Code	Sev	Message and Explanation
CGM	31	<p>Cannot get memory from OTS, FUNCT. return status [octal]</p> <p>The system does not have enough free memory to load the overlay link. The status is returned from the object-time system, and depends on the particular FUNCT. function that the overlay handler used. See Section 5.7.4 for the FUNCT. function codes and status messages.</p>
CRF	31	<p>Cannot read file [file] [reason]</p> <p>An error occurred when reading the overlay file. The file was closed after the last successful read operation.</p>
CSM	31	<p>Cannot shrink memory, FUNCT. return status [octal]</p> <p>A request to the object-time system to reduce memory, if possible, failed. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
CUF	31	<p>Cannot update file [file] [reason]</p> <p>An error occurred when updating the TMP file into which non-resident writable overlay links are written.</p>
CWF	31	<p>Cannot write file [file] [reason]</p> <p>An error occurred when creating the TMP file used to store non-resident writable overlay links.</p>
DLN	1	<p>Deleting link [name or number] after [hh:mm:ss]</p> <p>The named link has been removed from memory as a result of a call to the REMOV. subroutine. The time is elapsed time since the first call to the overlay handler. This message is output only to the overlay log file, if any.</p>
IAT	31	<p>Illegal argument type on call to [subroutine]</p> <p>A user call to the named overlay handler subroutine gave an illegal type of argument.</p>
IEF	31	<p>Input error for file [file], status [octal]</p> <p>An error occurred while reading the OVL or TMP file.</p>
ILN	31	<p>Illegal link number [number]</p> <p>A user call to one of the overlay handler subroutines gave an illegal link number as an argument.</p>
IMP	31	<p>Impossible error condition at PC=[address]</p> <p>This is an internal error caused by monitor call error returns that should not occur. This message is issued instead of the HALT message. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
IPE	31	<p>Input positioning error for file [file], status [octal]</p> <p>An error occurred while reading the OVL or TMP file.</p>

## OVERLAYS

Code	Sev	Message and Explanation
IVN	8	<b>Inconsistent version numbers</b>  The OVL and EXE files found were not created at the same time, and may not be compatible.
LNM	31	<b>Link number [decimal] not in memory</b>  A call to the REMOV. subroutine has removed the named link from memory. It must be restored by a call to GETOV. or RUNOV.
MAN	31	<b>Memory not available for absolute [link], FUNCT. return status [octal]</b>  There is not enough room for the overlay handler to load the specified link into the part of memory the link was built for. Two options are available: a) Use the /SPACE switch at load time to reserve more space for the link, or b) Build a relocatable overlay using the RELOCATABLE option to the /OVERLAY switch at load time.
MEF	31	<b>Memory expansion failed, FUNCT. return status [octal]</b>  The overlay handler was unable to get free space from the memory manager. Restructure your overlay so that the minimum number of links are in memory at any time.
NMS	8	<b>Not enough memory to load symbols, FUNCT. return status [octal]</b>  There was not enough free space available to load symbols into memory.
NRS	31	<b>No relocation table for symbols</b>  A relocation table was not included for the symbol table. It is possible that LINK failed to load the relocation table because there wasn't enough room in memory.
NSA	31	<b>No start address for link [name or number]</b>  A user call to the RUNOV. subroutine requests execution to continue at the start address of the named link, but that link has no start address.
NSD	31	<b>No such device for [file]</b>  An invalid device was specified.
OEF	31	<b>Output error for file [file], status [octal]</b>  An error occurred when writing the overlay file. The file was closed after the last successful write operation.
OPE	31	<b>Output positioning error for file [file], status [octal]</b>  An error occurred while writing the TMP file used to hold non-resident writable overlay links.

## OVERLAYS

Code	Sev	Message and Explanation
RLN	1	<b>Relocating link [name or number] at [address]</b>  The named relocatable link has been loaded at the given address. This message is output only to the overlay log file.
RLN	1	<b>Reading in link [name or number] after [time]</b>  The named link has been loaded. The time given is elapsed time since the first call to the overlay handler. This message is output only to the overlay log file.
STS	8	<b>OTS reserved space too small</b>  The object-time system does not have space for its minimum number of buffers. Reload, using the /SPACE switch for the root link with an argument greater than 2000 (octal).
ULN	31	<b>Unknown link name [name]</b>  A call to one of the overlay handler subroutines gave an invalid link name as an argument. Correct the call.
USC	8	<b>Undefined subroutine [name] called from [address]</b>  A required subroutine was not loaded. The instruction at the given program counter address calls for an undefined subroutine. Correct the call or load the required subroutine.
WLN	1	<b>Writing [link] after [time]</b>  The overlay handler is writing out a writable overlay link.

### 5.7.4 The FUNCT. Subroutine

Each DIGITAL-supplied object-time system has a subroutine that the overlay handler uses for memory management, I/O, and message handling. This subroutine has a single entry point, FUNCT., and is called by the sequence:

```
MOVEI    16,arglst
PUSHJ    17,FUNCT.
```

The format of the argument list is:

```
-<n+3>,,0
arglst:  Z 2,address of integer function code
          Z 2,address for error code on return
          Z 2,address for status code on return
          Z code,address of first argument
          .
          .
          .
          Z code,address of nth argument
```

## OVERLAYS

Where function code is one of the function codes described below; error code is a 3-letter ASCII mnemonic output by the object-time system (after ?, %, or []); and status (on return) contains one of the following values:

- 1 Function not implemented
- 0 Successful return
- n Number of the error message

Most object-time systems allocate separate space for their own use and for the use of the overlay handler. This minimizes the possibility that the overlay handler will request space that the object-time system is already using.

The permitted function code arguments, their names, and their meanings are:

Code	Name	Function
0	ILL	Illegal function; returns -1 status.
1	GAD	Get a specific segment of memory.
2	COR	Get a given amount of memory from anywhere in the space allocated to the overlay handler.
3	RAD	Return a specific segment of memory.
4	GCH	Get an I/O channel.
5	RCH	Return an I/O channel.
6	GOT	Get memory from the space allocated to the object-time system.
7	ROT	Return memory to the object-time system.
10	RNT	Get the initial runtime, in milliseconds, from the object-time system.
11	IFS	Get the initial runtime file specification of the program being run.
12	CBC	Cut back core (if possible) to reduce job size.
13	F.RRS	Read retain status (DBMS)
14	F.WRS	Write retain status (DBMS)
15	F.GPG	Get pages
16	F.RPG	Return pages
17	F.GPSI	Get TOPS-20 PSI channel
20	F.RPSI	Return TOPS-20 PSI channel

All FUNCT. codes are reserved to DEC.

The following subsections describe each function of the FUNCT. subroutine (except the reserved functions).

## OVERLAYS

### ILL (0) Function

This function is illegal. The argument list is ignored, and the status returned is -1.

### GAD (1) Function

The GAD function gets memory from a specific address in the space allocated to the overlay handler. The argument list points to:

arg 1 Address of requested memory  
arg 2 Size of requested allocation (in words)

A call to GAD with arg 2 equal to -1 requests all available memory.

On return, the status is one of the following:

0 Successful allocation  
1 Not enough memory available  
2 Memory not available at specified address  
3 Illegal arguments (address + size > 256K)

### COR (2) Function

The COR function gets memory from any available space allocated to the overlay handler. The arguments are:

arg 1 Undefined (address of allocated memory on return)  
arg 2 Size of requested allocation

On return, the status is:

0 Core allocated  
1 Not enough memory available  
3 Illegal argument (size > 256K)

### RAD (3) Function

The RAD function returns the memory starting at the specified address to the overlay handler. The arguments are:

arg 1 Address of memory to be returned  
arg 2 Size of memory to be returned (in words)

On return, the status is one of the following:

0 Successful return of memory  
1 Memory cannot be returned  
3 Illegal argument (address or size > 256K)



## OVERLAYS

### GOT (6) Function

The GOT function gets memory from the space allocated to the object-time system. Its arguments are:

arg 1 Undefined (address of allocated memory on return)  
arg 2 Size of memory requested

On return, the status is one of the following:

0 Successful allocation  
1 Not enough memory available  
3 Illegal argument (size > 256K)

### ROT (7) Function

The ROT function returns memory to the object-time system. Its arguments are:

arg 1 Address of memory to be returned  
arg 2 Size of memory to be returned (in words)

On return, the status is one of the following:

0 Successful return of memory  
1 Memory cannot be returned  
3 Illegal argument (address or size > 256K)

### RNT (10) Function

The RNT function returns the initial runtime, in milliseconds, from the object-time system. (At the beginning of the program, the object-time system will have executed a RUNTIM UUO; the result is the time returned by RNT.) Its arguments are:

arg 1 Undefined (contains initial runtime on return)  
arg 2 Ignored

On return, the runtime is in arg 1, and the status is 0. The status is 0.

### IFS (11) Function

The IFS function returns the initial runtime file specification from the object-time system. (This initial file specification is the one used to begin the program; that is, it was given with a compile-class command.) Its arguments are:

arg 1 Undefined (SIXBIT device on return)  
arg 2 Undefined (SIXBIT filename on return)  
arg 3 Undefined (project-programmer number on return)

On return, the status is one of the following:

0 Successful return  
1 Error

## OVERLAYS

### CBC (12) Function

The CBC function cuts back memory if possible, which reduces the size of the job. It uses no arguments, and the returned status is 0.

### RRS (13) Function ( Reserved for DBMS )

Returns ARG1 = 0. On return, the status is always 0.

### WRS (14) Function ( Reserved for DBMS )

Returns ARG1 = 0. On return, the status is always 0.

### GPG (15) Function

The GPG function is used to fetch a page. The arguments are:

arg2: size to be allocated, in words

On return,

arg1 = address of allocated memory, on page boundary

and the status is one of the following:

0 if allocated OK  
1 if not enough memory  
3 if argument error

### RPG (16) Function

The RPG function is used to return pages. The arguments are:

arg1: address (a word)  
arg2: size (in words)

On return, the status is:

0 if deallocated OK  
1 if wasn't allocated  
3 if argument error

### GPSI (17)

The GPSI function can be used to get a PSI channel for programs running in a TOPS-20 environment. This entry point provides only controlled access to the PSI tables. It will arrange that the tables exist and that SIR and EIR have been done but does not do AIC or any other JSYS necessary to set up the channel (ATI or MTOPR, for example).

The arguments are:

arg1: channel number,  
or -1 to allocate any user-assignable channel  
arg2: level number  
arg3: address of interrupt routine

## OVERLAYS

On return, arg1 contains the channel number allocated (if -1 was originally specified). On return, the status is:

- 0 if OK
- 1 if channel was already assigned
- 2 if no free channels
- 3 if argument error

### NOTE

This function is used by TOPS-20 programs. It is a reserved function in the TOPS-10 environment.

### RPSI (20) Function

This entry point provides only controlled access to the PSI tables. It does not do DIC or any other JSYS necessary to release a channel. It just clears the level and interrupt address fields in CHNTAB.

This function accepts the following argument:

arg1: channel number

On return the status is one of the following:

- 0 if OK
- 1 if channel wasn't in use
- 3 if argument error

### NOTE

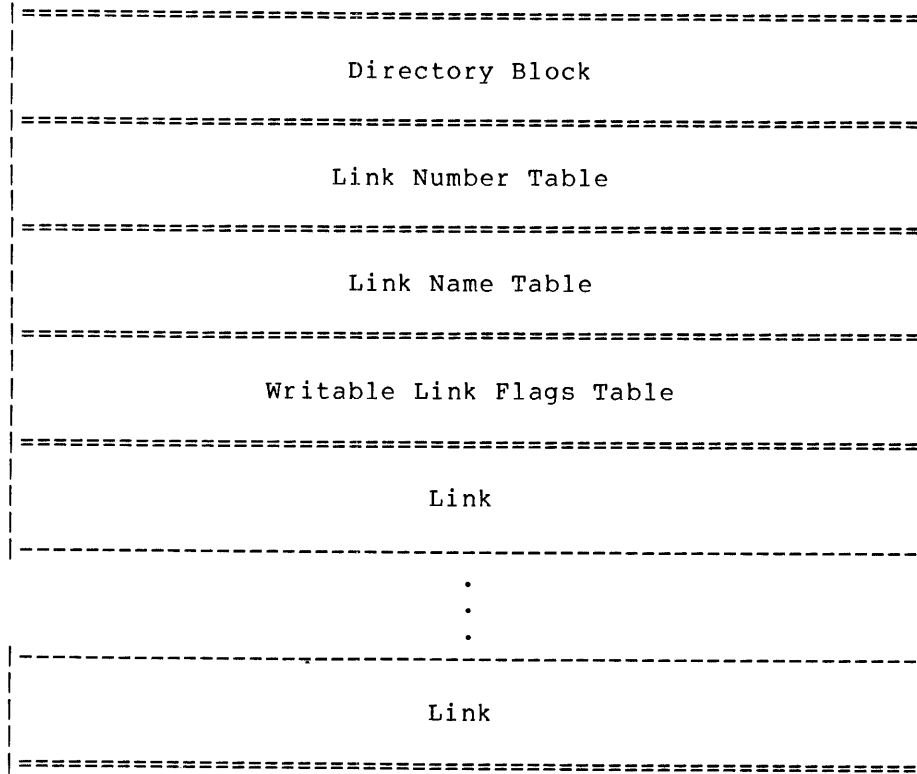
This function is used by TOPS-20 programs. It is a reserved function in the TOPS-10 environment.

## OVERLAYS

### 5.8 THE OVERLAY (OVL) FILE

This section contains diagrams of the contents of the overlay file output by LINK as a result of the /OVERLAY switch. The following diagram shows the overall scheme of the file:

Scheme of the Overlay (OVL) File



## OVERLAYS

### 5.8.1 The Directory Block

The following diagram shows the contents of the Directory Block:

#### Directory Block

.DIHDR:	0 (Reserved)	Length of Directory Block
.DIRGN:	0 (Reserved)	
.DIVER:	Version Number of Corresponding EXE file	
.DILPT:	-(Size of Link No. Table)	Link Number Table Block No.
.DINPT:	-(Size of Link Name Table)	Link Name Table Block No.
.DIWPT:	-(Size of Writable Flg Tbl)	Writable Flg Tbl Block No
.DIFLG:	Flags	
	0 (Reserved)	

In the fourth word above, the size of the Link Number Table (in words) is half the number of links (rounded upward); the Link Number Table Block No. is the number of the 128-word disk block containing the Link Number Table. (There are four disk blocks per disk page.)

In the fifth word above, the size of the Link Name Table (in words) is twice the number of links; the Link Name Table Block No. is the number of the 128-word disk block containing the Link Name Table.

The table defined by the .DIWPT word above consists of a string of two-bit bytes. The first bit, OW.WRT, indicates whether the corresponding overlay link is writable. This bit is set under the control of a REL block of type 1045 (writable links). The second bit, OW.PAG, indicates whether the corresponding overlay link is currently paged into the runtime overlay temporary file. This is strictly a run-time flag and should be zero in the overlay file. This flag is defined in the overlay file to allow the overlay handler to set up its flag table with a single read operation.

The .DIFLG word in the directory block contains a single bit flag (bit 0). If this bit is set the overlay file contains at least one writable overlay. This information is also contained in the Writable Link Table. However, by having the information available in the directory block the overlay handler can determine if any links are writable without scanning the Writable Link Table. All other bits in the .DIFLG word are reserved and must be zero.

#### NOTE

If a user requests both writable and relocatable overlays, only halfwords known to be relocatable at load time will be correctly relocated when the link is refetched.

## OVERLAYS

### 5.8.2 The Link Number Table

The following diagram shows the contents of the Link Number Table:

**Link Number Table**

Pointer to Link 0		Pointer to Link 1
Pointer to Link 2		Pointer to Link 3
⋮		
Pointer to Link n-1		Pointer to Link n

Each pointer is a disk block number. Any unused words in the last disk block of the Link Number Table are zeros.

### 5.8.3 The Link Name Table

The following diagram shows the contents of the Link Name Table:

**Link Name Table**

Link Number	
SIXBIT Link Name	
⋮	
Link Number	
SIXBIT Link Name	

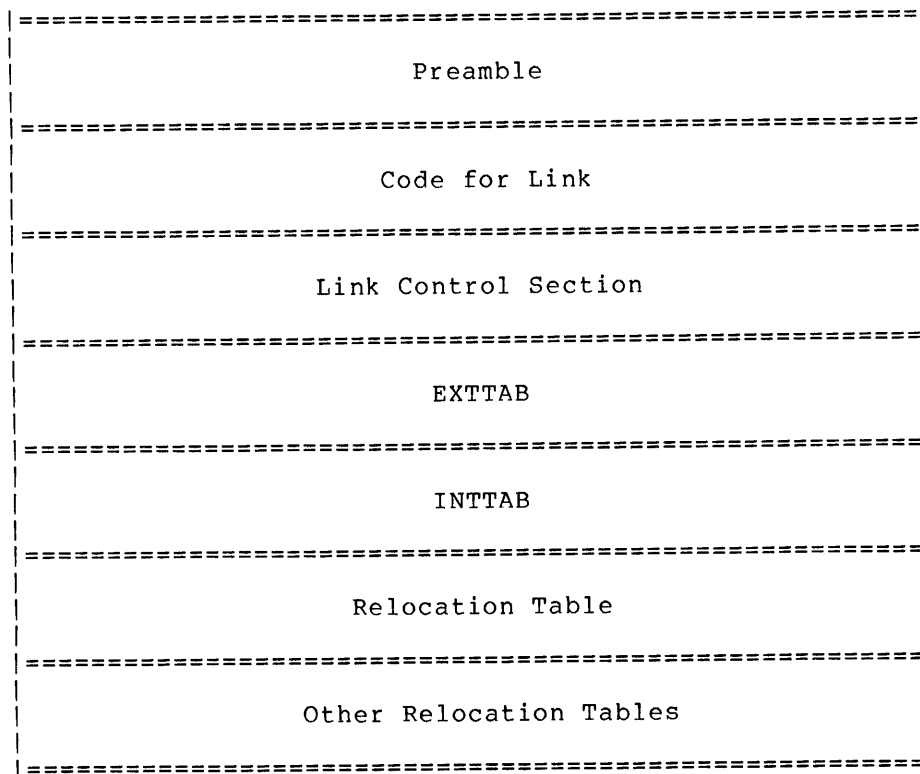
Any unused words in the last disk block of the Link Name Table are zeros.

## OVERLAYS

### 5.8.4 The Overlay Link

The following diagram shows the overall scheme of each overlay link in the overlay file:

Scheme of an Overlay Link



## OVERLAYS

### The Preamble

The following diagram shows the contents of the preamble for an overlay link:

Preamble	
0 (Reserved)	Length of Preamble
0 (Reserved)	0 (Reserved)
0 (Reserved)	Link Number
SIXBIT Link Name	
Pointer to List of Bound Links Starting with Root Link	
Pointer to List of Bound Links Ending with Root Link	
Equivalence Pointer	
Address of Control Section	
Flags	
Absolute Address at Which Link Loaded	
Length of Link (Code through INTTAB)	
Disk Block Number of Start of Link Code	
0 (Reserved)	
Disk Block Number of Relocation Table	
Disk Block Number of Other Relocation Tables	
0 (Reserved)	
Disk Block Number of Radix-50 Symbols	
Block Number of Relocation Tables for Radix-50 Symbols	
Next Free Memory Location for Next Link	



## OVERLAYS

### Code for the Link

The code for each link consists of a core image that was constructed from the REL files placed in the link. This core image contains the code and data for the link.

### The Control Section

The following diagram shows the contents of the Control Section:

#### Control Section

0 (Reserved)	Length of Header
0 (Reserved)	0 (Reserved)
0 (Reserved)	Link Number
SIXBIT Link Name	
Ptr to Ancestor in Core -(Length of Symbol Table)	Ptr to Successor in Core Address of Symbol Table
0 (Reserved)	Start Address for Link
Memory Needed to Load Link -(Length of EXTTAB)	First Address in Link Pointer to EXTTAB
-(Length of INTTAB)	Pointer to INTTAB
Address of Symbols on Disk	
Relocation Address	
Copy of Block Number for Code	
-(Length of Radix-50 SymTab)	Blk No. of Radix-50 SymTab

## OVERLAYS

### The EXTTAB Table

The following diagram shows the contents of the EXTTAB table:

**EXTTAB**

=====	
JSP 1,.OVRLA	
Flags	Address of Callee's INTTAB
Callee's Link Number	Ptr to Callee's Control Sec
Backward Pointer	Forward Pointer
=====	
:	
:	
:	
=====	
JSP 1,.OVRLA	
Flags	Address of Callee's INTTAB
Callee's Link Number	Ptr to Callee's Control Sec
Backward Pointer	Forward Pointer
=====	

The flags in the left half of the second word have the following meanings:

Bit	Meaning (if bit is on)
0	Module is in core.
1	Module is in more than one link.
2	Relocatable link is already relocated.

## OVERLAYS

### The INTTAB Table

The following diagram shows the contents of the INTTAB table:

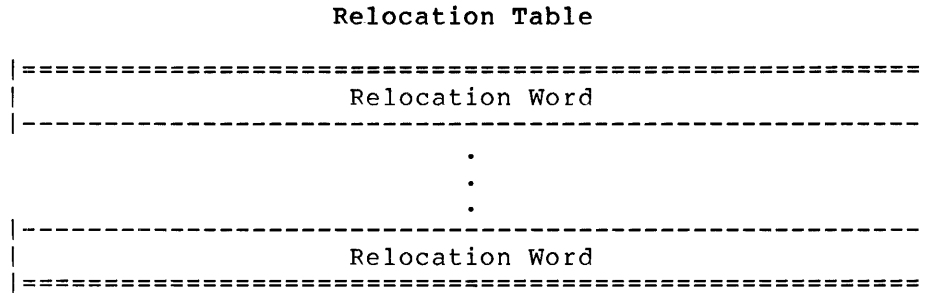
#### INTTAB

0 (Reserved)		Address of Entry Point
0 (Reserved)		Forward Pointer
.		
.		
.		
0 (Reserved)		Address of Entry Point
0 (Reserved)		Forward Pointer

## OVERLAYS

### The Relocation Table

The following diagram shows the contents of the Relocation Table:



The Relocation Table contains one bit for each halfword of the link. If the bit is on, the halfword is relocatable; if it is off, the halfword is not relocatable.

The first word contains the relocation bits for the first 22 (octal) words of the link; the second word contains the relocation bits for the next 22 (octal) words; and so forth for all words in the link.

This table exists only when relocatable overlays are requested with the /OVERLAY:RELOCATABLE switch.

## OVERLAYS

### The Other Relocation Tables

The following diagram shows the contents of the Other Relocation Tables:

#### Other Relocation Tables

Number of Words Following for This Link	
Link Number	Planned Load Address
Relocation Halfword	Ptr to Words of Code
.	
.	
.	
Relocation Halfword	Ptr to Words of Code
.	
.	
.	
Number of Words Following for This Link	
Link Number	Planned Load Address
Relocation Halfword	Ptr to Words of Code
.	
.	
.	
Relocation Halfword	Ptr to Words of Code

This table exists only when relocatable overlays have been requested with the OVERLAY/RELOCATABLE switch. The Other Relocation Tables are used to hold internal LINK references.



## CHAPTER 6

### PSECTS

PSECTS (Program SECTIONS) are programmer- or system-defined regions of code and data that LINK relocates in memory. PSECTS are used to structure a program's memory space, or to load a program that uses extended addressing.

#### 6.1 LOADING PROGRAMS WITH PSECTS

When loading programs with PSECTS, you must specify the origin of the PSECT. LINK then uses this PSECT origin to store the data in the PSECTS.

To specify a PSECT origin, include the origin in the source program or use the LINK /SET switch. See the appropriate language manual for including the origin in the source program and Chapter 3 for the /SET switch.

Defining an upper bound is also important when loading PSECTS. The LINK /LIMIT switch defines an upper bound for a PSECT. If the PSECT loads to this bound, LINK returns a warning and an error message. Despite these messages, LINK continues to process input files and to load code. The warning is:

```
%LNKPEL PSECT [psect] exceeded limit of [address]
```

Although LINK does continue to process input files and load code, the program is incomplete and should not be used. LINK does issue the following fatal error:

```
?LNKCFS Chained fixups have been suppressed
```

Chained fixups are a method that LINK uses to resolve symbol references.

Using /LIMIT to define an upper bound prevents unintended PSECT overlaps. PSECT overlaps can cause LINK to loop and produce other unpredictable behavior.

For example, the LRGPRO and BIGPRO modules each contain two PSECTS, BIG and GRAND. LRGPRO is loaded and /COUNTERS is used to check PSECT origins and current values. PSECT origins are found by looking under the initial value column and PSECT current values are found by looking under the current value column of the /COUNTERS output. The upper bound is found by looking under the limit value column.

```
.R LINK (RET)
*/SET:BIG:1000 (RET)
*/SET:GRAND:5400 (RET)
*LRGPRO (RET)
```

PSECTs

```
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.    initial value  current value  limit value
  .LOW.                0              140            1000000
  BIG                  1000            5100           1000000
  GRAND                5400            10500          1000000]
```

\*

/COUNTERS shows that the current value for PSECT BIG and the initial value for PSECT GRAND are close together in memory. The current value for BIG is 5100 and the PSECT origin for GRAND is 5400. The /LIMIT switch can now be used to restrict PSECT BIG's current value to PSECT GRAND's initial value using the following:

```
*/LIMIT:BIG:GRAND (RET)
```

/LIMIT prevents an unintended overlap because it causes LINK to issue a warning if the current value for BIG exceeds GRAND's origin. The warning is:

```
%LNKPEL PSECT [psect] exceeded limit of [address]
```

The warning message indicates that the PSECTs overlapped, and that PSECTs BIG and GRAND need to be farther apart in memory. The /COUNTERS switch shows a new current value greater than 5400. Notice that the limit set with the /LIMIT switch is shown in the limit value column.

```
*BIGPRO (RET)
```

```
%LNKPEL PSECT BIG exceeded limit of 5400
detected in module .MAIN from file DSK:BIGPRO.REL[12,3456]
```

```
*/COUNTERS (RET)
```

```
[LNKRLC Reloc. ctr.    initial value  current value  limit value
  .LOW.                0              140            1000000
  BIG                  1000            6300           5400
  GRAND                5400            10500          1000000]
```

/GO continues loading the program, and LINK issues a warning and fatal error message. The warning is:

```
%LNKPOV Psects [psect] and [psect] overlap from address [address] to
[address]
```

The fatal error message is:

```
?LNKCFS chained fixups have been suppressed
```

For example,

```
*/GO (RET)
```

```
%LNKPOV Psects BIG and GRAND overlap from address 5400 to 6300
?LNKCFS chained fixups have been suppressed
```

EXIT

Now, LINK is re-run and the PSECTs are moved farther apart in memory. In this example, GRAND's origin is reset from 5400 to 7000.

```
.R LINK (RET)
```

```
*/SET:BIG:1000 (RET)
```

```
*/SET:GRAND:7000 (RET)
```

```
*LRGPRO (RET)
```



## PSECTS

```
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.   initial value   current value   limit value
  .LOW.                0                140            1000000
  BIG                  1000             5100           1000000
  GRAND               7000             10500          1000000]
```

\*

```
*/LIMIT:BIG:GRAND (RET)
*BIGPRO (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.   initial value   current value   limit value
  .LOW.                0                140            1000000
  BIG                  1000             6300           5400
  GRAND               7000             10500          1000000]
```

```
*/GO (RET)
```

EXIT

.

### 6.2 PSECT ATTRIBUTES

PSECT attributes specify how LINK stores a PSECT in memory, and the page access of the PSECT.

The CONCATENATED or OVERLAID attribute specifies how LINK stores PSECTS.

#### 6.2.1 CONCATENATED and OVERLAID

LINK uses the CONCATENATED or OVERLAID attributes when loading PSECTS into memory. These attributes are specified when the PSECT is defined in the source program, and are contained in REL Blocks 24 and 1050. See Appendix A for information on these blocks. If the attribute is not specified, LINK uses CONCATENATED.

The following example illustrates how PSECTS are stored in memory. In this example, modules MAINKO and MAINKC contain three PSECTS, ALPHA, BETA, and GAMMA. There is an additional module named SUBMD1. The ALPHA and BETA PSECTS have the CONCATENATE attribute. The GAMMA PSECT, which is a data PSECT declared in each module, has the OVERLAID attribute defined in MAINKO and the CONCATENATE attribute defined in MAINKC.

First, LINK is run and the origin is set for PSECTS ALPHA, BETA, and GAMMA.

```
.R LINK (RET)
*/SET:ALPHA:3000/SET:BETA:5000/SET:GAMMA:7000 (RET)
```

Next, MAINKO is loaded with GAMMA defined as OVERLAID, and /COUNTERS is used to display the initial, current, and limit values.

```
*MAINKO ;OVERLAID GAMMA (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.   initial value   current value   limit value
  .LOW.                0                140            1000000
  ALPHA               3000             3017           1000000
  BETA                 5000             5011           1000000
  GAMMA               7000             7025           1000000]
```

PSECTS

Now, SUBMD1 is loaded, /COUNTERS is used, and /GO is used to load the modules and exit LINK.

Notice that the current values for ALPHA and BETA have increased, and that the current value for GAMMA remains the same.

```
*SUBMD1 (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.    initial value    current value    limit value
  .LOW.                0                140              1000000
  ALPHA               3000             3033             1000000
  BETA                5000             5041             1000000
  GAMMA               7000             7025             1000000]
*/GO (RET)
```

EXIT

.

In the following example, LINK is run and the origin is set for PSECTS ALPHA, BETA, and GAMMA.

```
.R LINK (RET)
*/SET:ALPHA:3000/SET:BETA:5000/SET:GAMMA:7000 (RET)
```

Now, MAINKC is loaded with GAMMA defined as CONCATENATE, and /COUNTERS is used.

```
*MAINKC                ;CONCATENATED GAMMA (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.    initial value    current value    limit value
  .LOW.                0                140              1000000
  ALPHA               3000             3017             1000000
  BETA                5000             5011             1000000
  GAMMA               7000             7025             1000000]
*
```

Next, SUBMD1 is loaded, /COUNTERS is used, and /GO is used to load the modules and exit LINK.

Notice that all current values have increased.

```
*SUBMD1 (RET)
*/COUNTERS (RET)
[LNKRLC Reloc. ctr.    initial value    current value    limit value
  .LOW.                0                140              1000000
  ALPHA               3000             3033             1000000
  BETA                5000             5041             1000000
  GAMMA               7000             7035             1000000]
*/GO (RET)
```

EXIT

.

## PSECTS

### 6.2.2 RWRITE

The RWRITE attribute sets the page access for PSECTS to read/write. This attribute can be set in the source program or omitted, as RWRITE is the default.

For example, the following MACRO statement defines read/write access for the ALL PSECT.

```
.PSECT ALL/RWRITE,1000
```

1000 is the PSECT's origin.

#### NOTE

The RONLY attribute does exist for compatibility with MACRO-20 and LINK-20, but is not recommended for use on TOPS-10. You should not use RONLY on TOPS-10 because TOPS-10 does not support read-only pages in the low-segment and LINK considers PSECTS to be part of the low segment. LINK does build an .EXE file from .REL files that have the RONLY attribute set, but these files cannot be run on TOPS-10.



## APPENDIX A

### REL BLOCKS

The object modules that LINK loads are output from the language translators. These object modules are formatted into REL (RELocatable) Blocks, each of which contains information for LINK.

This appendix describes each type of REL Block and gives its format. Terms used throughout this discussion are defined as follows:

**Header Word:** a fullword giving the REL Block Type in its left half and a short count or long count in its right half.

**Short Count:** a halfword giving the length of the REL Block, excluding relocation words (which appear before each group of 18 decimal, or 22 octal words), and excluding the header word.

**Long Count:** a halfword giving the length of the REL Block, including all words in the block except the header word itself.

**Relocation Word:** a fullword containing the relocation bits for up to 18 decimal or 22 octal words following words. Each relocation bit is either 1, indicating a relocatable halfword, or 0, indicating a nonrelocatable halfword.

The first two relocation bits give the relocatability of the left and right halves, respectively, of the next following word; the next two bits give the relocatability of the two halves of the second following word; and so forth for all bits in the word, except any unused bits, which will be zero.

If a REL Block has relocation words, the first one follows the header word. If more than 18 (decimal) data words follow this relocation word, the next word (after the 18 words) is another relocation word. Thus, a REL Block that has relocation words will have one for each 18 words of data that it contains. If the REL Block does not contain an integral multiple of 18 words, the last relocation word will have unused bits.

#### NOTE

A block with a zero short count does not include a relocation word.

**Data Word:** Any word other than a header word or a relocation word.

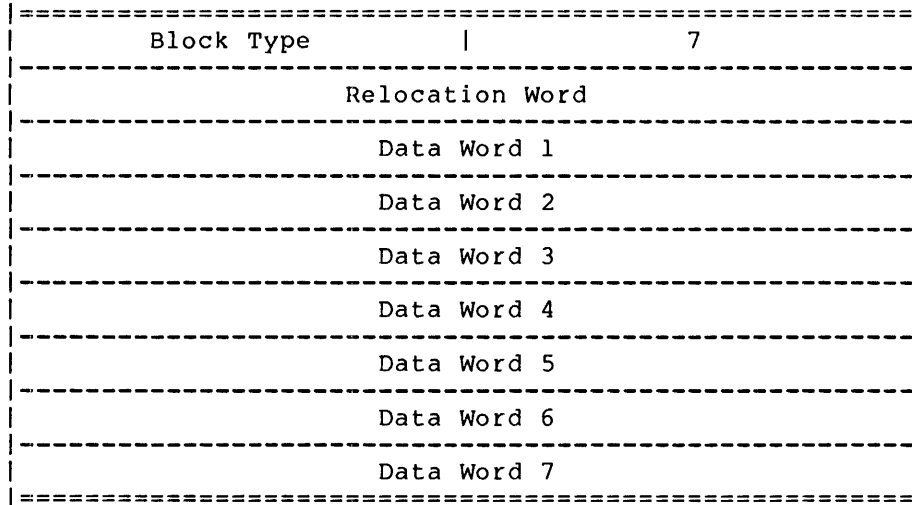
**MBZ:** Must Be Zero.

#### NOTE

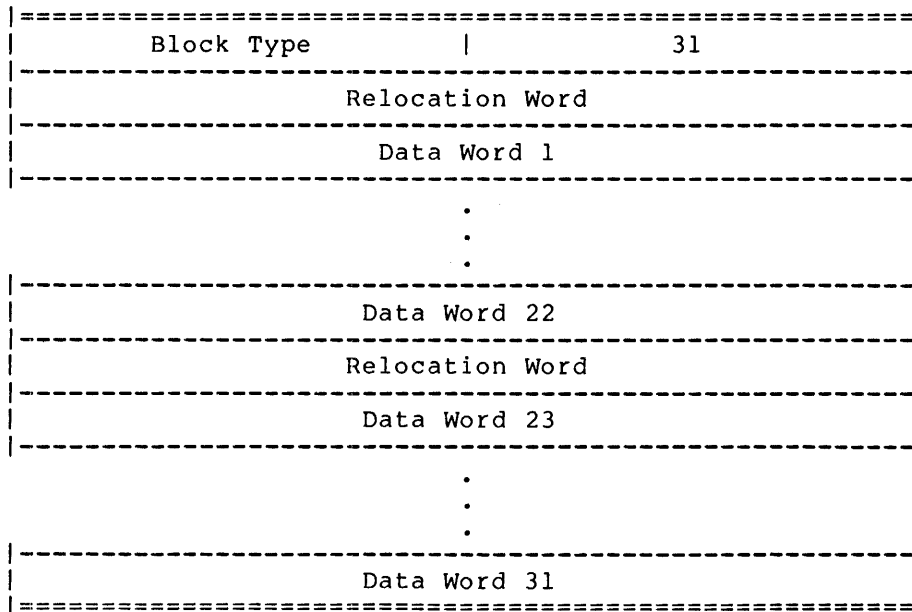
All numbers in this appendix are octal unless specifically noted as decimal.

## REL BLOCKS

The diagram below shows a REL Block having a short count of 7, and a relocation word.



The diagram below shows a REL Block having a short count of 31 and two relocation words.



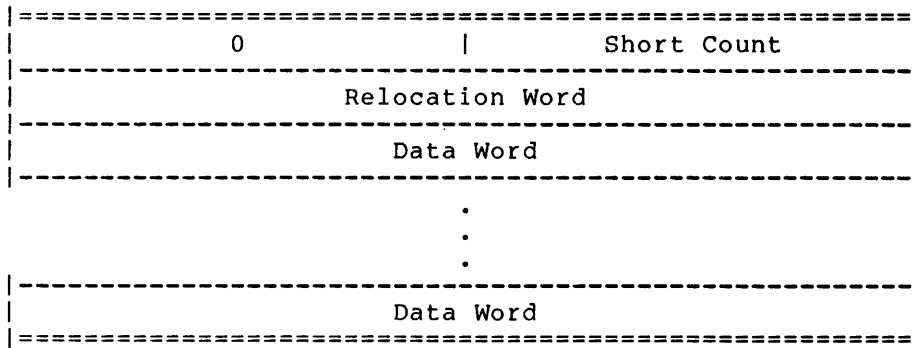
## REL BLOCKS

REL Block Types must be numbered in the range 0 to 777777. The following list shows which numbers are reserved for DIGITAL, and which for customers:

Type Numbers	Use
0 - 37	Reserved for DIGITAL
40 - 77	Reserved for customers
100 - 401	Reserved for DIGITAL
402 - 577	Reserved for customers
600 - 677	Reserved for customer files
700 - 777	Reserved for DIGITAL files
1000 - 1777	Reserved for DIGITAL
2000 - 3777	Reserved for customers
4000 - 777777	Reserved for ASCII text

REL BLOCKS

Block Type 0 (Ignored)



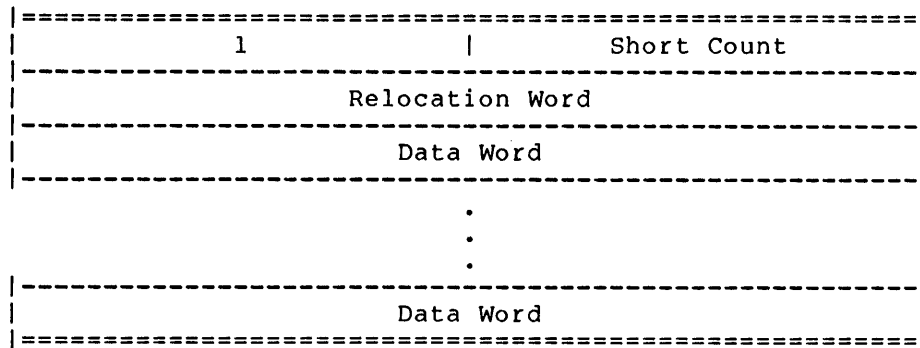
Block Type 0 is ignored by LINK.

If the short count is 0, then no relocation word follows, and the block consists of only one word. This is how LINK bypasses zero words in a REL file.



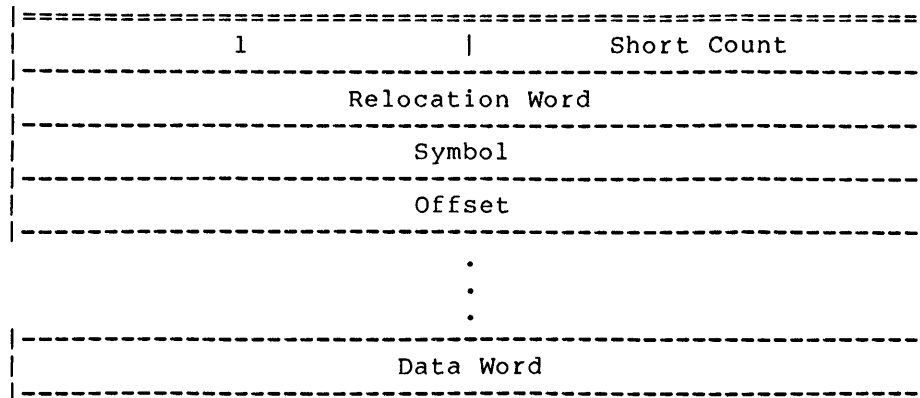
## REL BLOCKS

### Block Type 1 (Code)



Block Type 1 contains data and code. The first data word gives the address at which the data is to be loaded. This address can be relocatable or absolute, depending on the value of bit 1 of the relocation word. The remaining data words are loaded beginning at that address.

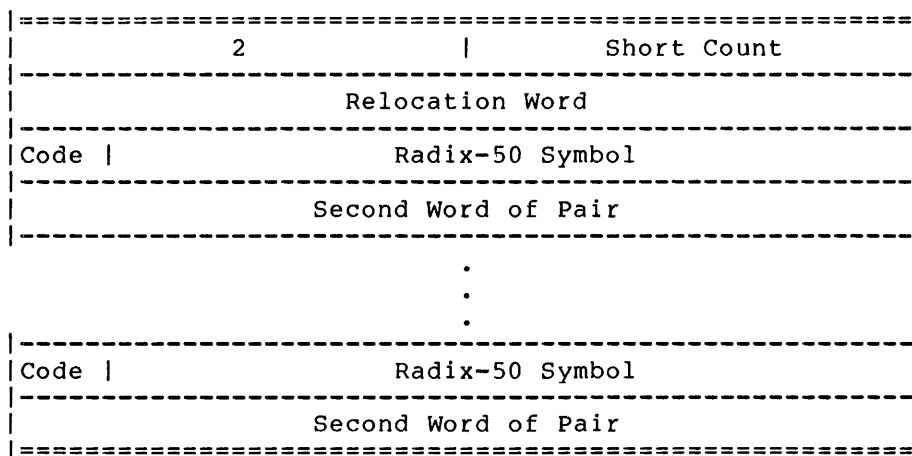
If the start address is given in symbolic, the following format of Block Type 1 is used:



In this alternate format, the first four bits of the first data word (Symbol) are 1100 (binary), and the word is assumed to be a Radix-50 symbol of type 60. The load address is calculated by adding the value of the global symbol to the offset given in the following word. The third and following data words are loaded beginning at the resulting address. The global symbol must be defined when the Type 1 Block is found.

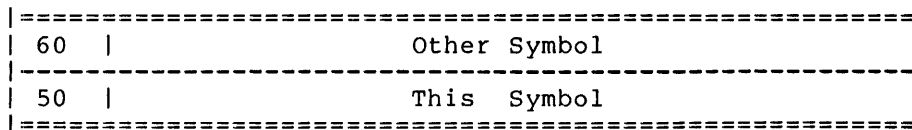
REL BLOCKS

Block Type 2 (Symbols)



The first word of each pair has a code in bits 0 to 3 and a Radix-50 symbol in bits 4 to 35 (decimal). The contents of the second word of a pair depends on the given code. The octal codes and their meanings are:

Code	Meaning
00	This code is illegal in a symbol block.
04	The given symbol is a global definition. Its value, contained in the second word of the pair, is available to other programs.
10	The given symbol is a local definition, and its value is contained in the second word of the pair. If the symbol is followed by one of the special pairs or by a Polish REL Block (as explained below, under code 24), the symbol is considered a partially defined local symbol. Otherwise, it is considered fully defined.
14	The given symbol is a block name (from a translator that uses block structure). The second word of the pair contains the block level. The symbol is considered local; if local symbols are loaded, the value of the block name is entered in the symbol table as its block level.
24	The given symbol is a global definition. However, it is only partially defined at this time, and LINK cannot yet use its value. If the symbol is defined in terms of another symbol, then the next entry in the REL file must be a word pair in a Block Type 2 as follows:



In this format, code 50 indicates that the right half of the word depends on the other symbol.

## REL BLOCKS

### Code

### Meaning

If the partially defined symbol is defined in terms of a Polish expression, then the next entry in the REL file must be Block Type 11 (Polish), whose store operator gives this symbol as the symbol to be fixed up. A fixup resolves the symbol. The store operator must be -4 or -6.

- 30 The given symbol is a global definition. However, it is only partially defined at this time, and LINK cannot yet use its value. If the symbol is defined in terms of another symbol, then the next entry in the REL file must be a word pair in a Block Type 2 as follows:

```

=====|
| 60 |           Other Symbol |
|-----|
| 70 |           This Symbol  |
|-----|
=====|
  
```

In this format, code 70 indicates that the left half of the word depends on the other symbol.

If the partially defined symbol is defined in terms of a Polish expression, then the next entry in the REL file must be Block Type 11 (POLISH), whose store operator gives this symbol as the symbol to be fixed up. The store operator must be -5.

- 34 The given symbol is a global definition. However, it is only partially defined at this time, and LINK cannot yet use its value. If the symbol is defined in terms of another symbol, then the next entry in the REL file must be a word pair in a Block Type 2 as follows:

```

=====|
| 60 |           Other Symbol |
|-----|
| 50 |           This Symbol  |
|-----|
| 60 |           Other Symbol |
|-----|
| 70 |           This Symbol  |
|-----|
=====|
  
```

This format indicates that both halves of the word depend on the other symbol.

- 44 The given symbol is a global definition exactly as in code 04, except that it is not output by DDT.
- 50 The given symbol is a local symbol exactly as in code 10, except that it is not output by DDT.
- 60 The given symbol is a global request. LINK's handling of the symbol depends on the value of the code in the first four bits of the second word of the pair. These codes and their meanings are:

- 00 The right half of the word gives the address of the first word in a chain of requests for the global memory address. In each request, the right half of the word gives the address of the next request. The chain ends when the address is 0.

## REL BLOCKS

Code	Meaning
40	The right half of the word contains an address. The right half of the value of the requested symbol is added to the right half of this word.
50	The rest of the word contains a Radix-50 symbol whose value depends on the requested global symbol. (If the given Radix-50 symbol is not the one defined in the previous word pair, then this word is ignored.) When the value of the requested symbol is resolved, it is added to the right half of the value of the Radix-50 symbol.
60	The right half of the word contains an address. The right half of the value of the requested symbol is added to the left half of this word.
70	The rest of the word contains a Radix-50 symbol whose value depends on the requested global symbol. (If the given Radix-50 symbol is not the one defined in the previous word pair, then this word is ignored.) When the value of the requested global symbol is resolved, it is added to the left half of the value of the Radix-50 symbol.
64	The given symbol is a global definition exactly as in code 24, except that it is not output by DDT.
70	The given symbol is partially defined, where the left half is deferred, as in code 30, except that it is not output by DDT.
74	The given symbol is partially defined, where the right half is deferred, as in code 34, except that it is not output by DDT.

Symbols are placed in the symbol table in the order that LINK finds them. However, DDT expects to find the symbols in a specific order.

For a non-block-structured program, that order is:

Program Name

Symbols for Program

For a block-structured program whose structure is:

```
Begin Block 1 (same as program name)
  Begin Block 2
  End Block 2
  Begin Block 3
    Begin Block 4
    End Block 4
  End Block 3
End Block 1
```

## REL BLOCKS

the order is:

Program Name (Block 1)  
Block Name 2  
Symbols for Block 2  
Block Name 4  
Symbols for Block 4  
Block Name 3  
Symbols for Block 3  
Block Name 1  
Symbols for Block 1

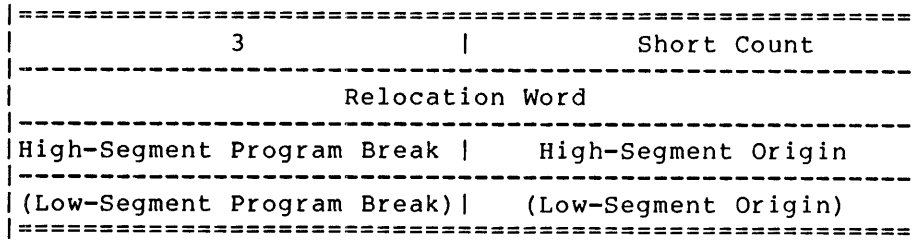
This ordering follows the rule that the name and symbols for each block must occur in the symbol table in the order of the block endings in the program.

### NOTES

1. Only one fixup by a Type 2, 10, 11, or 12 Block is allowed for a given field. (There can be separate fixups for the left and right halves of the same word.)
2. Fixups are not necessarily performed in the order LINK finds them.

## REL BLOCKS

### Block Type 3 (HISEG)



Block Type 3 tells LINK that code is to be loaded into the high segment.

Short Count is either 1 or 2.

If the left half of the first data word is 0, subsequent Type 1 blocks found are assumed to have been produced by the MACRO pseudo-op HISEG. This usage is not recommended. It means that the addresses in the blocks are relative to 0, but are to be placed in the program high segment. The right half of the first data word is the beginning of the high segment (usually 400000).

If the left half of the first data word is nonzero (the preferred usage), subsequent Type 1 blocks found are assumed to have been produced by the MACRO pseudo-op TWOSEG.

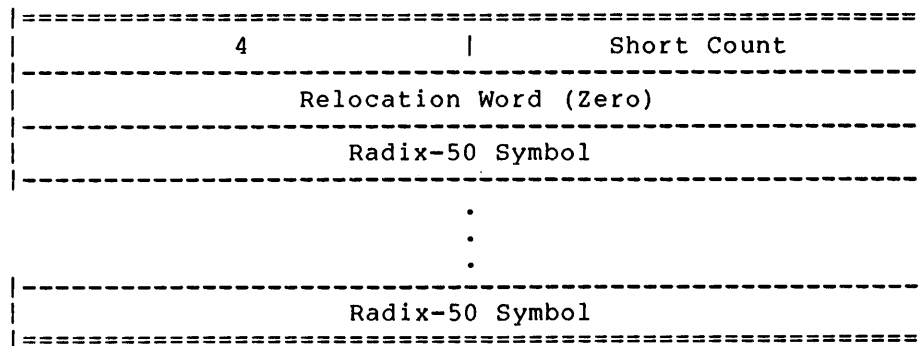
The right half is interpreted as the beginning of the high segment, and the left half is the high-segment break; the high-segment length is the difference of the left and right halves.

(One-pass translators that cannot calculate the high-segment break should set the left half equal to the right half.)

If the second word appears in the HISEG block, its left half shows the low-segment program break, and its right half shows the low-segment origin (usually 0).

## REL BLOCKS

### Block Type 4 (Entry)



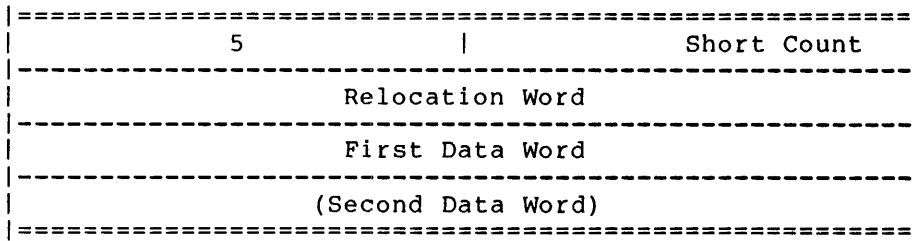
Block Type 4 lists the entry name symbols for a program module. If a Type 4 block appears in a module, it must be the first block in the module. A library file contains a Type 4 block for each of its modules.

When LINK is in library search mode, the symbols in the block are compared to the current list of global requests for the load. If one or more matches occur, the module is loaded and the name of the module is marked as an entry point in map files, etc. If no match occurs, the module is not loaded.

If LINK is not in library search mode, no comparison of requests and entry names is made, and the module is always loaded. Refer to Block Type 17 for more information about libraries. Refer to block type 14.

## REL BLOCKS

### Block Type 5 (End)



Block Type 5 ends a program module. A Block Type 6 must be encountered earlier in the module than the Type 5 block.

Short Count is 1 or 2.

If the module contains a two-segment program, the first data word is the high-segment break and the second data word is the low-segment break. If the module contains a one-segment program, the first data word is the program break and the second data word is the absolute break. If count is 1, then second word is assumed to be 0.

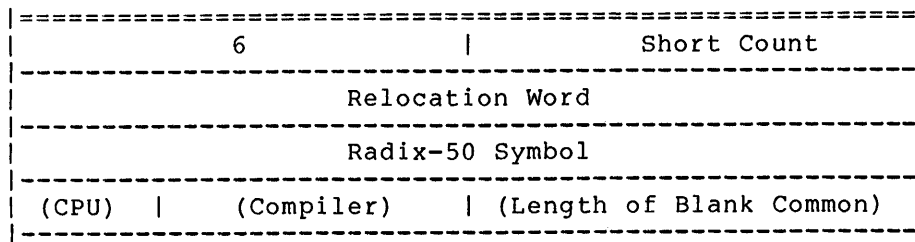
Each PRGEND pseudo-op in a MACRO program generates a Type 5 REL block. Therefore a REL file may contain more than one Type 5 block.

A library REL file has a Type 5 block at the end of each of its modules.



## REL BLOCKS

### Block Type 6 (Name)



Block Type 6 contains the program name, and must precede any Type 2 blocks. (A module should begin with a Type 6 block and end with a Type 5 block.)

Short Count is 1 or 2.

The first data word is the program name in Radix-50 format; this name cannot be blanks. The second data word is optional; if it appears, it contains CPU codes in bits 0 to 5, a compiler code in bits 6 to 17 (decimal), and the length of the program's blank COMMON in the right halfword.

The CPU codes specify processors for program execution as:

Bit 2	KS10
Bit 3	KL10
Bit 4	KI10
Bit 5	KA10

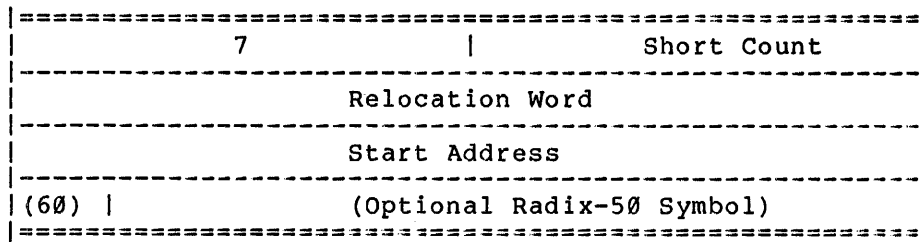
If all these bits are off, then any of the processors can be used for execution.

The compiler code specifies the compiler that produced the REL file. The defined codes are:

0	Unknown	7	SAIL	16	COBOL-74
1	Not used	10	FORTRAN	17	COBOL
2	COBOL-68	11	MACRO	20	BLISS-36
3	ALGOL	12	FAIL	21	BASIC
4	NELIAC	13	BCPL	22	SITGO
5	PL/I	14	MIDAS	23	(Reserved)
6	BLISS	15	SIMULA	24	PASCAL
				25	JOVIAL

## REL BLOCKS

### Block Type 7 (Start)



Block Type 7 contains the start address for program execution. LINK uses the start address in the last such block processed by the load, unless /START or /NOSTART switches specify otherwise.

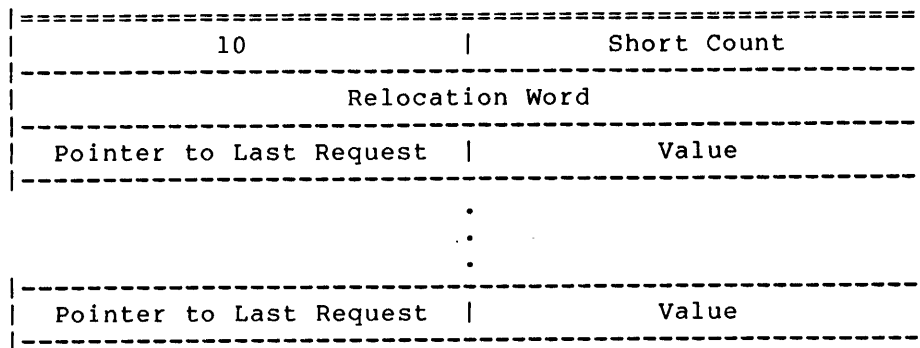
Short Count is 1 or 2.

If the second (optional) word is present, it must be a Radix-50 symbol with the code 60; LINK forms the start address by adding the value of the symbol to the value in the right half of the preceding word (Start Address).

LINK builds an entry vector if it is specified or non-zero.

## REL BLOCKS

### Block Type 10 (Internal Request)



Block Type 10 is generated by one-pass compilers to resolve requests caused by forward references to internal symbols. The MACRO assembler also generates Type 10 blocks to resolve requests for labels defined in literals; a separate chain is required for each PSECT in a PSECTed program.

Each data word contains one request for an internal symbol. The left half is the address of the last request for a given symbol. The right half is the value of the symbol. The right half of the last request contains the address of the next-to-last request, and so on, until a zero right half is found. (This is exactly analogous to Radix-50 code 60 with second-word code 00 in a Block Type 2.)

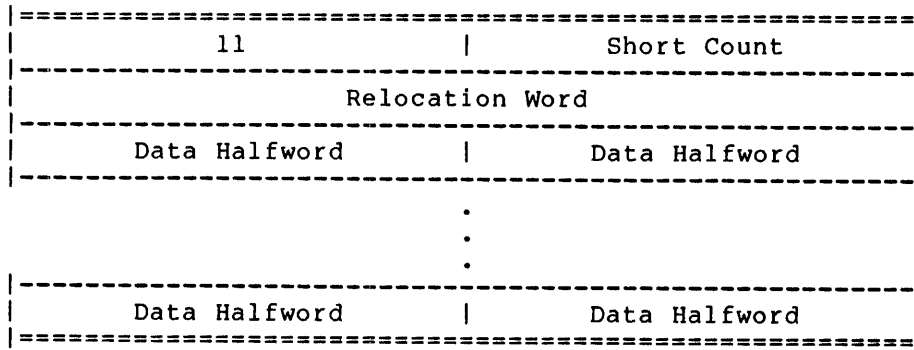
If a data word contains -1, then the following word contains a request for the left (rather than right) half of the specified word. In this case, the left half of the word being fixed up contains the address of the next-to-last left half request, and so on, until a zero left half is found. (This is a left half chain analogous to the right half chain described above.)

#### NOTES

1. Only one fixup by a Type 2, 10, 11, or 12 is allowed for a given field. (There can be separate fixups for the left and right halves of the same word.)
2. Fixups are not necessarily performed in the order LINK finds them.

## REL BLOCKS

### Block Type 11 (Polish)



Block Type 11 defines Polish fixups for operations on relocatable values or external symbols. Only one store operator code can appear in a Block Type 11; this store operator code can be either a symbol fixup code or a chained fixup code. The store operator code appears at the end of the block.

#### NOTES

1. Only one fixup by a Type 2, 10, 11, or 12 Block is allowed for a given field. (There can be separate fixups for the left and right halves of the same word.)
2. Fixups are not necessarily performed in the order LINK finds them.

The data words of a Type 11 block form one Polish string of halfwords. Each halfword contains one of the following:

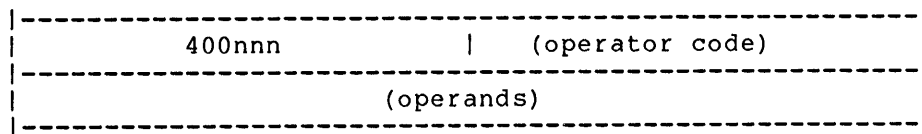
1. A symbol fixup store operator code.  
A symbol fixup defines the value to be stored in the value field of the symbol table for the given symbol. A symbol fixup store operator code is followed by two or four data halfwords.
2. A chained fixup store operator code.  
A chained fixup takes a relocatable address whose corrected virtual address is the location for storing or chaining. A chained fixup store operator code is followed by one data halfword.
3. A data type code. Data type code 0 is followed by a data halfword; a data type code 1 or 2 is followed by two data halfwords.
4. An arithmetic or logical operator code.

## REL BLOCKS

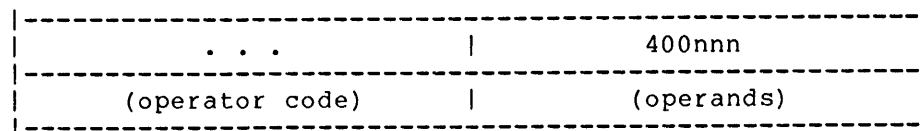
5. A PSECT index code. This code defines a PSECT index to be used for calculating the relocated addresses that appear in this block. PSECT indexes are needed only for PSECTed programs.

A global PSECT index is associated with a Block Type 11. This index appears as the first halfword after the relocation word, and it defines the PSECT for the store address or store symbol. Any addresses for a different PSECT must be preceded by a different PSECT index.

Thus, a relocatable data halfword in a different PSECT must appear in one of the following formats:



OR



where the different PSECT index is nnn+1.

Any relocatable address that does not have an explicit preceding PSECT index code preceding its data type code is assumed to be in the same PSECT as the store address for the block. The current PSECT may be set by a previous REL Block type.

6. A halfword of data (preceded by a data type 0 halfword) or two halfwords of data (preceded by a data type 1 or 2 halfword).

A sequence of halfwords containing a data type code 0 and a data halfword can begin in either half of a word.

The codes and their meanings are:

### Symbol Fixup Store Operator Codes:

- 7 Fullword replacement. No chaining is done.
- 6 Fullword symbol fixup. The following one or two words contain the Radix-50 symbol(s) (with their 4-bit codes). The first is the symbol to be fixed up, and the second is the block name for a block-structured program (0 or nonexistent for other programs).

## REL BLOCKS

- 5 Left half symbol fixup. The following one or two words contain the Radix-50 symbols. The first is the symbol to be fixed up, and the second is the block name for a block-structured program (0 or nonexistent for other programs).
- 4 Right half symbol fixup. The following one or two words contain the Radix-50 symbols. The first is the symbol to be fixed up, and the second is the block name for a block-structured program (0 or nonexistent for other programs).

### Chained Fixup Store Operator Codes:

- 3 Fullword chained fixup. The halfword following points to the first element in the chain. The entire word pointed to is replaced, and the old right half points to the next fullword.
- 2 Left half chained fixup. The halfword following points to the first element in the chain.
- 1 Right half chained fixup. The halfword following points to the first element in the chain.

### Data Type Codes:

- 0 The next halfword is an operand.
- 1 The next two halfwords form a fullword operand.
- 2 The next two halfwords form a Radix-50 symbol that is a global request. The operand is the value of the symbol.

### Arithmetic and Logical Operator Codes:

#### NOTE

Operands are read in the order that they are encountered.

- 3 Add.
- 4 Subtract.
- 5 Multiply.
- 6 Divide.
- 7 Logical AND.
- 10 Logical OR.
- 11 Logical shift. (A positive second operand causes a shift to the left. A negative operand causes a shift to the right.)
- 12 Logical XOR.
- 13 Logical NOT (one's complement).

## REL BLOCKS

- 14 Arithmetic negation (two's complement).
- 15 Count leading zeros (like JFFO instruction). Refer to the MACRO Assembler Reference Manual for information about the ^L operand, which this code implements.
- 16 Remainder.
- 17 Magnitude.
- 20 Maximum.
- 21 Minimum.
- 22 Comparison. Returns 0 if the two operands are different; -1 if they are equal.
- 23 Used to resolve the links in a chain. Refer to Block Type 12.
- 24 Symbol definition test. Returns 0 if the operand (a Radix-50 symbol) is unknown; 1 if it is known but undefined; -1 if it is known and defined.
- 25 Skip N words of Polish.
- 26 Skip N words of Polish on some condition.
- 27 Return contents of location N.

### PSECT Index Codes:

400nnn PSECT index nnn, where nnn is a 3-digit octal integer.

For an example of a Type 11 block, the MACRO statements

```

        EXTERN B
A:      EXP <A*B+A>

```

Generate (assuming that A has a relocatable value of zero):

=====	
11	6
-----	
00 01 00 00 10 10	0
-----	
3 (Add)	5 (Multiply)
-----	
0 (Halfword Operand Next)	0 (Relocatable)
-----	
2 (Fullword Radix-50 Next)	1st Half of Radix-50 B
-----	
2nd Half of Radix-50 B	0 (Halfword Operand Next)
-----	
0 (Relocatable)	-3 (Rh Chained Fixup Next)
-----	
0 (Chain Starts at 0')	. . .
=====	

The first word contains the block type (11) and the short count (6). The second word is the relocation word; it shows that the following halfwords are to be relocated: right half of second following word, left half of fifth following word, left half of sixth following word.

## REL BLOCKS

The next word shows that the two operations to be performed are addition and multiplication; because this is in Polish prefix format, the multiplication is to be performed on the first two operands first, then addition is performed on the product and the third operand.

The next two halfwords define the first operand. The first halfword is a data type code 0, showing that the operand is a single halfword; the next halfword is the operand (relocatable 0).

The next three halfwords define the second operand. The first of these halfwords contains a data type code 2, showing that the operand is two halfwords containing a Radix-50 symbol with code 60. The next two halfwords give the symbol (B).

The next two halfwords define the third operand. The first of these halfwords contains a data type code 0, showing that the operand is a single halfword; the next halfword gives the value of the operand (relocatable 0).

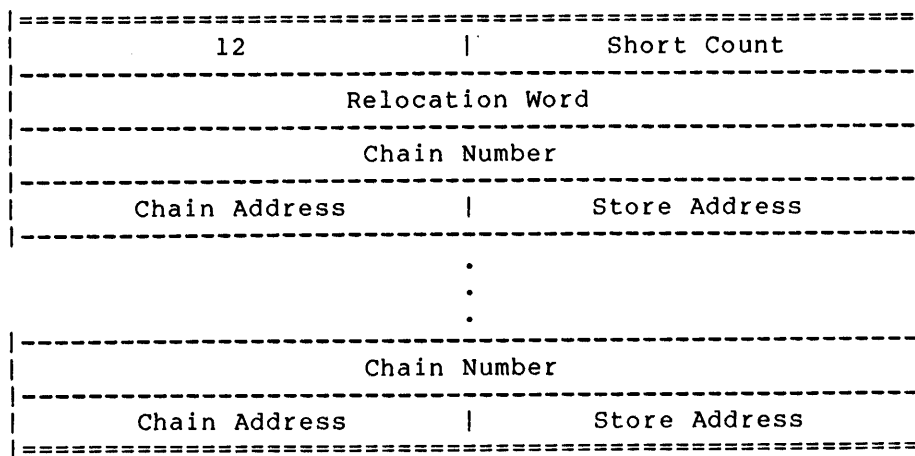
The next two halfwords give the store operator for the block. The first of these halfwords contains the chained fixup store operator code -3, showing that a fullword chained fixup is required; the next halfword contains the operand (relocatable 0), showing that the chain starts at relocatable zero.

The last halfword is irrelevant, and should be zero. If it is not, LINK issues the LNKJPB error message.



## REL BLOCKS

### Block Type 12 (Chain)



Block Type 12 chains together data structures from separately compiled modules. (The MACRO pseudo-ops `.LINK` and `.LNKEND` generate Type 12 blocks.) Block Type 12 allows linked lists that have entries in separately compiled modules to be constructed so that new entries can be added to one module without editing or recompiling any other module.

The data words in a Type 12 block are paired. The first word of each pair contains a chain number between 1 and 100 (octal). (The chain number is negative if the pair was generated by a `.LNKEND` pseudo-op.) The second word contains a store address in the right half, and a chain address in the left half. The store address points to the location where `LINK` will place the chain address of the last entry encountered for the current chain. The first entry in a chain has a zero in the word pointed to by the store address.

A MACRO statement of the form:

```
.LINK chain-number,store-address,chain-address
```

generates a word pair in a Type 12 block as shown above. A MACRO statement of the form:

```
.LINK chain-number,store-address
```

generates a word pair in a Type 12 block with a 0 for the chain address field in the REL block. A MACRO statement of the form:

```
.LNKEND chain-number,store-address
```

generates a word pair in a Type 12 Block with a 0 for the chain address and a negative chain number.

As `LINK` processes a load, it performs a separate chaining for each different chain number found; thus a word pair in a Type 12 block is related to all other word pairs having the same chain number (even in other loaded modules). Type 12 pairs having different chain numbers (even in the same module), are not related.

## REL BLOCKS

### NOTE

Chain numbers above 100 (octal) are reserved by Digital.

To show how the chains are formed, we will take some pairs from different programs having the same chain number (1 in the example). The following four programs contain .LINK or .LNKEND pseudo-ops for the chain numbered 1. After each program, the word pair generated in the Type 12 block appears.

### NOTES

1. When LINK stores an address that results from a Type 12 REL Block, only the right half of the receiving location is written. You can safely store another value in the left half; it will not be overwritten.
2. Only one fixup by a Type 2, 10, 11, or 12 Block is allowed for a given field. (There can be separate fixups for the left and right halves of the same word.)
3. Fixups are not necessarily performed in the order LINK finds them.

REL BLOCKS

Example

```
TITLE MOD0
.
.
TAG0:  BLOCK 1
.
.
.LNKEND 1, TAG0
.
.
END
```

-1	
0	Value of TAG0

```
TITLE MOD1
.
.
TAG1:  BLOCK 1
.
.
.LINK 1, TAG1
.
.
END
```

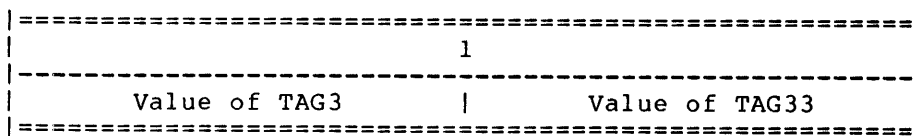
1	
0	Value of TAG1

```
TITLE MOD2
.
.
TAG2:  BLOCK 1
.
.
.LINK 1, TAG2
.
.
END
```

1	
0	Value of TAG2

## REL BLOCKS

```
TITLE MOD3
      .
      .
TAG3:  BLOCK 1
      .
      .
TAG33: BLOCK 1
      .
      .
      .LINK 1, TAG33, TAG3
      .
      .
      END
```



Suppose we load MOD0 first. The .LNKEND statement for MOD0 generates a negative chain number. LINK sees the negative chain number (-1) and recognizes this as the result of a .LNKEND statement for chain number 1. LINK remembers the store address (value of TAG0) as the base of the chain.

Next we load MOD1. The .LINK statement for MOD1 does not use the third argument, so the chain address is 0. LINK sees that this is the first entry for chain number 1. Because it is the first entry, LINK places a 0 in the store address (value of TAG1). LINK then remembers the value of TAG1 for use in the next chain entry. (If the chain address is 0, as it is in MOD1, LINK remembers the store address; if the chain address is nonzero, LINK remembers the chain address.)

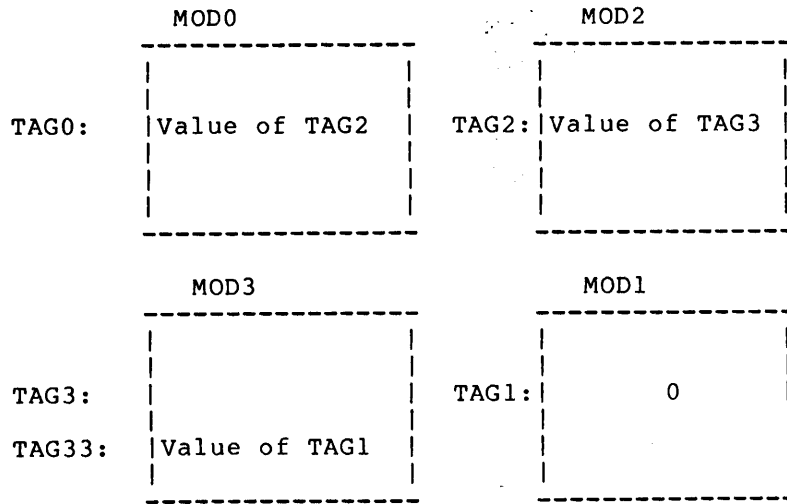
Next we load MOD3. The .LINK statement in MOD3 uses a third argument (TAG3), therefore, the value of TAG3 is used as the chain address. LINK places its remembered address (value of TAG1) in the store address (value of TAG33). Because the chain address (value of TAG3) is nonzero, LINK remembers it for the next entry.

Finally we load MOD2. Like MOD1, the .LINK statement for MOD2 does not take a third argument, and thus the chain address is 0. LINK places the remembered address (value of TAG3) in the store address (value of TAG2). Because the chain address is 0, LINK remembers the store address (value of TAG2).

At the end of loading, LINK places the last remembered address (value of TAG2) at the address (value of TAG0) given by the .LNKEND statement in MOD0.

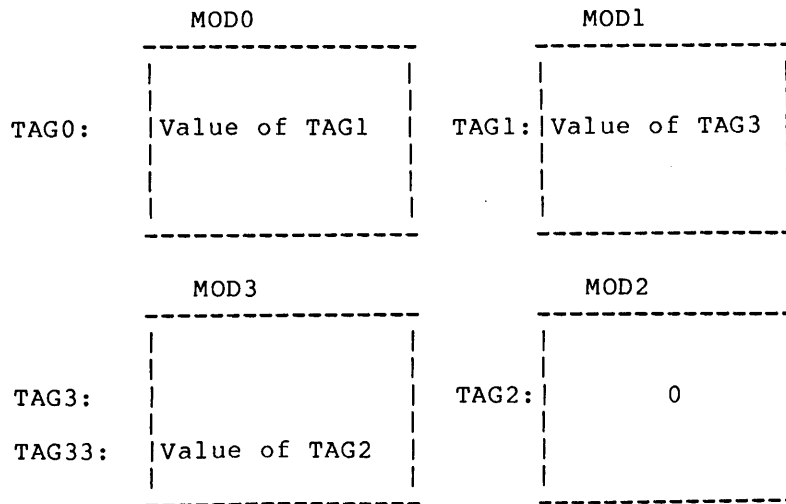
## REL BLOCKS

The results of the chaining can be seen in the following diagram of the loaded core image:



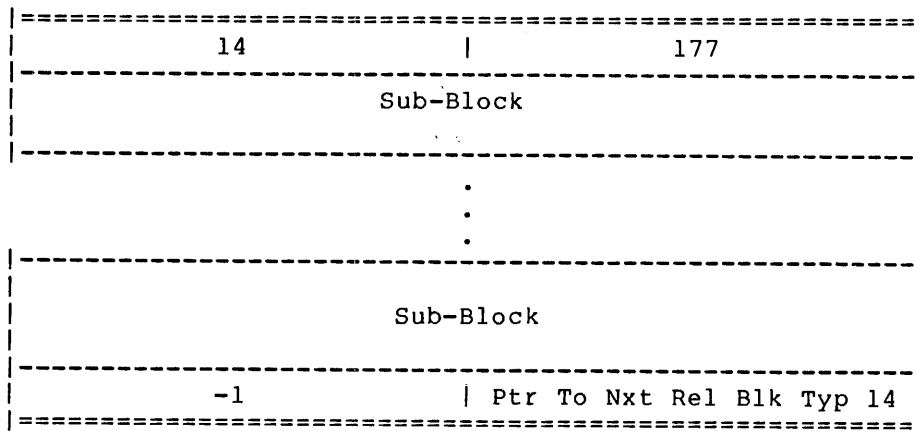
Note that the order of loading for modules with .LINK statements is critical. (A module containing a .LNKEND statement can be loaded any time; its treatment is not affected by the order of loading.)

For example, if we load the four programs in the order MOD2, MOD3, MOD0, MOD1, we get a different resulting core image:

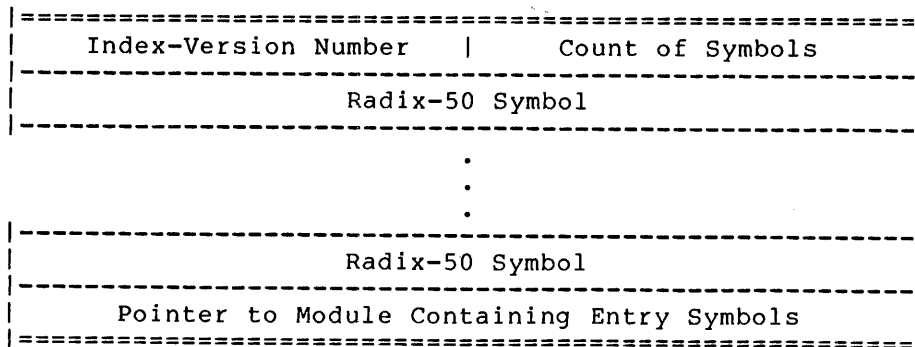


## REL BLOCKS

### Block Type 14 (Index)



Each sub-block is of the form:



Block Type 14 contains a list of all entry points in a library produced by MAKLIB. The block contains 177 (octal) data words (with no relocation words); if the index requires more entries, additional Type 14 blocks are used. If 177 data words are not needed, zero words pad the block to a length of 177. -1 indicates the end of the sub-block information.

The Type 14 block consists of a header word, a number of sub-blocks, and a trailer word containing the disk block address of the next Type 14 block, if any. Each disk block is 128 words.

Each sub-block is like a Type 4 block, with three differences:

1. The sub-block has no relocation words.
2. The last word of the sub-block points to the module that contains the entry points listed in the sub-block. The right half of the pointer has the disk block number of the module within the file; the left half has the number of words (in that block) that precede the module. If there is no next block, then the word after the last sub-block is -1.
3. The index-version number is used so that old blocks can still be loaded, even if the format changes in the future.

## REL BLOCKS

### Block Type 15 (ALGOL)

15	Short Count
Relocation Word	
Load Address	Length
Chain Address	Offset
.	
.	
Chain Address	Offset

Block Type 15 is used to build the special ALGOL OWN block.

The first data word contains the length of the module's OWN block in the right half, and the desired load address for the current OWN block in the left half. Each following word contains an offset for the start of the OWN block in the right half, and the address of a standard righthalf chain of requests for that word of the OWN block in the left half.

When LINK sees a REL Block Type 15, it allocates a block of the requested size at the requested address. The length of the block is then placed in the left half of the first word, and the address of the last OWN block seen is placed in the right half. If this is the first OWN block seen, 0 is stored in the right half of the first word.

The remaining data words are then processed by adding the address of the first word of the OWN block to each offset, and then storing the resulting value in all the locations chained together, starting with the chain address.

At the end of loading, LINK checks to see if the symbol %OWN is undefined. If it is undefined, then it is defined to be the address of the last OWN block seen. In addition, if LINK is creating an ALGOL symbol file, the file specification of the symbol file is stored in the first OWN block loaded. This file specification must be standard TOPS-10 format and can include (in order): device, file name, file extension, and project-programmer number.

## REL BLOCKS

### Block Type 16 (Request Load)

16	Short Count
Relocation Word (Zero)	
SIXBIT Filename	
Project-Programmer Number	
SIXBIT Device	
.	
.	
.	
SIXBIT Filename	
Project-Programmer Number	
SIXBIT Device	

Block Type 16 contains a list of files to be loaded. The data words are arranged in triplets; each triplet contains information for one file: file name, project-programmer number, and device. The file extension is .REL.

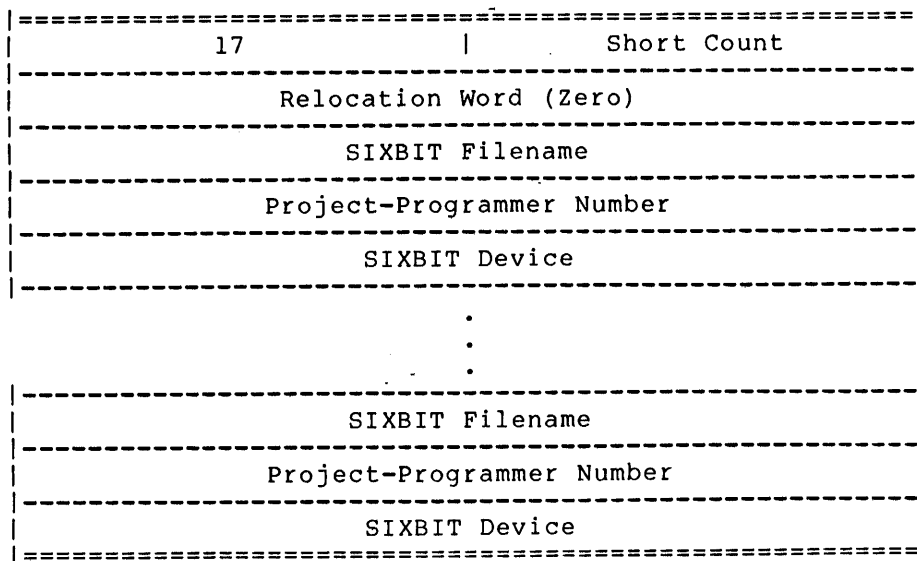
LINK saves the specifications for the files to be loaded, discarding duplicates. At the end of loading, LINK loads all specified files immediately before beginning library searches.

The MACRO pseudo-op .REQUIRE generates a Type 16 REL Block.



## REL BLOCKS

### Block Type 17 (Request Library)

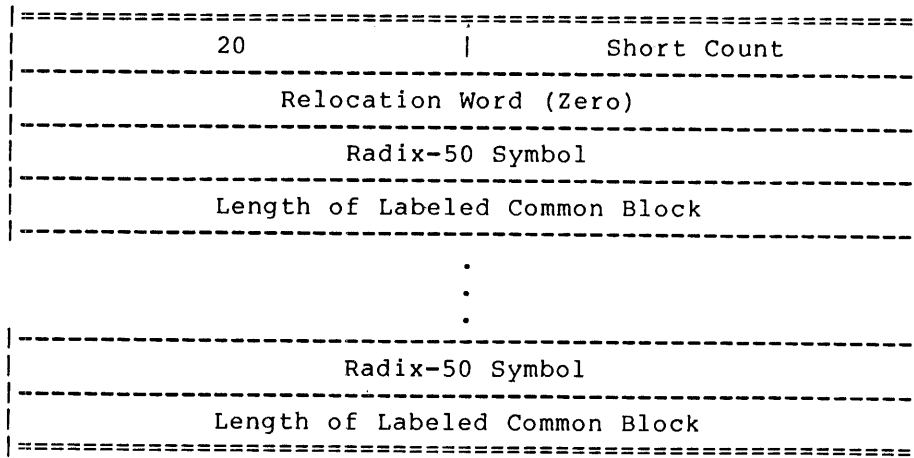


Block Type 17 is identical to Block Type 16 except that the specified files are loaded in library search mode. The specified files are searched after loading files given in Type 16 blocks, but before searching system or user libraries.

The MACRO pseudo-op `.REQUEST` generates a Type 17 REL Block.

## REL BLOCKS

### Block Type 20 (Common)



Block Type 20 allocates labeled COMMON areas. The label for unlabeled COMMON is ".COMM.". If a Block Type 20 appears in a REL file, it must appear before any other block that causes code to be loaded or storage to be allocated in the core image.

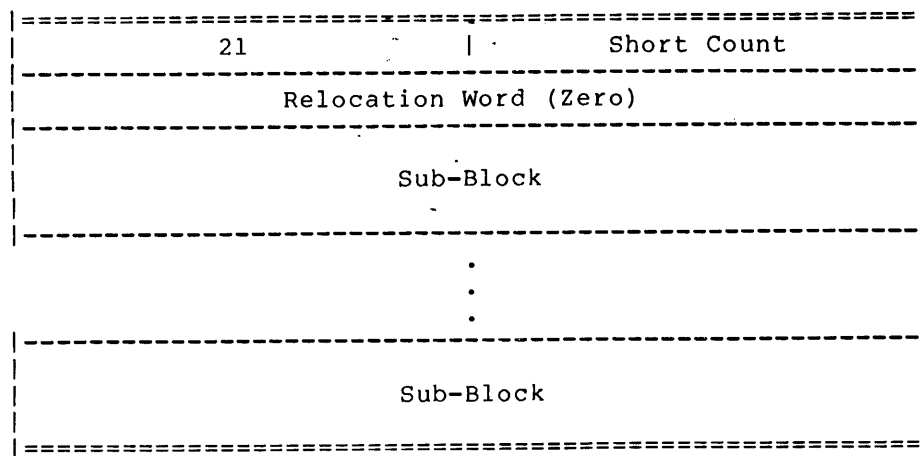
The data words are arranged in pairs. The first word of each pair contains a COMMON name in Radix-50 format (the four-bit code field must contain 60). The second contains the length of the area to be allocated.

For each COMMON entry found, LINK first determines whether the COMMON area is already allocated. If not, LINK allocates it. If the area has been allocated, the allocated area must be at least as large as the current requested allocation.

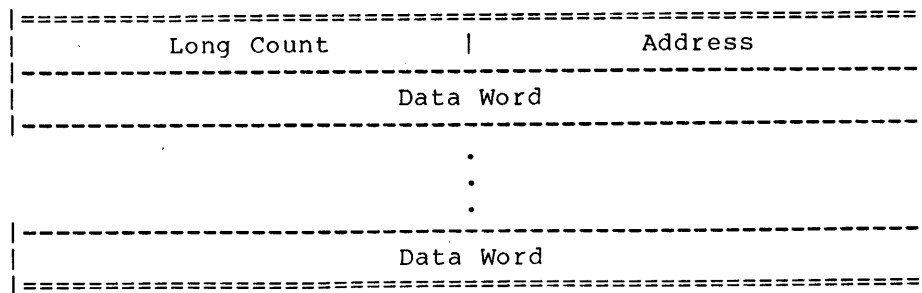
COMMON blocks can be referenced from other block types as standard globally defined symbols. However, a COMMON block must be initially allocated by a REL Block Type 20, by a REL Block Type 6 (for blank COMMON), or by the /COMMON switch to LINK. Any attempt to initially define a COMMON block with a standard global symbol definition causes the LNKSNC error when the redefining Block Type 20 is later seen.

## REL BLOCKS

### Block Type 21 (Sparse Data)



Each sub-block is of the form:



Block Type 21 contains data to be loaded sparsely in a large area. The first word of each sub-block contains the long count for the sub-block in the left half, and the address for loading the data words in the right half.

If the first four bits of the first data word of each sub-block are 1100 (binary) then the word is assumed to be a Radix-50 symbol of type 60; in this case the left half of the second word is the sub-block count, and the right half plus the value of the symbol is the load address.

## REL BLOCKS

### Block Type 22 (PSECT Origin)

22		Short Count
Relocation Word		
(SIXBIT PSECT Name) or (PSECT Index)		
PSECT Origin		

Block Type 22 contains the PSECT origin (base address).

Block Type 22 tells LINK to set the value of the relocation counter to the value of the counter associated with the given PSECT name. All following REL blocks are relocated with respect to this PSECT until the next Block Type 22 or 23 is found.

When data or code is being loaded into this PSECT, all relocatable addresses are relocated for the PSECT counter.

MACRO generates a Block Type 22 for each .PSECT and .ENDPS pseudo-op it processes. These Type 22 blocks are interleaved with the other blocks to indicate PSECT changes. A Type 22 block is also generated at the beginning of each symbol table to show to which PSECT the table refers.

REL BLOCKS

Block Type 23 (PSECT End Block)

23	Short Count
PSECT Index	
PSECT Break	

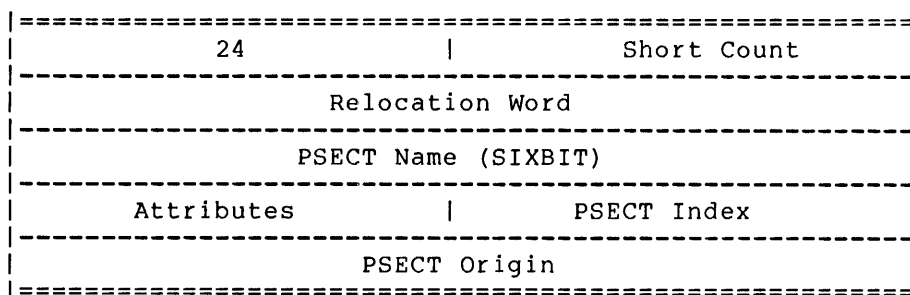
Block Type 23 contains information about a PSECT.

The PSECT index uniquely identifies the PSECT within the module being loaded. The Type 24 block assigns the index.

The PSECT break gives the length of the PSECT. This break is relative to the zero address of the current module, not to the PSECT origin.

## REL BLOCKS

### Block Type 24 (PSECT Header Block)



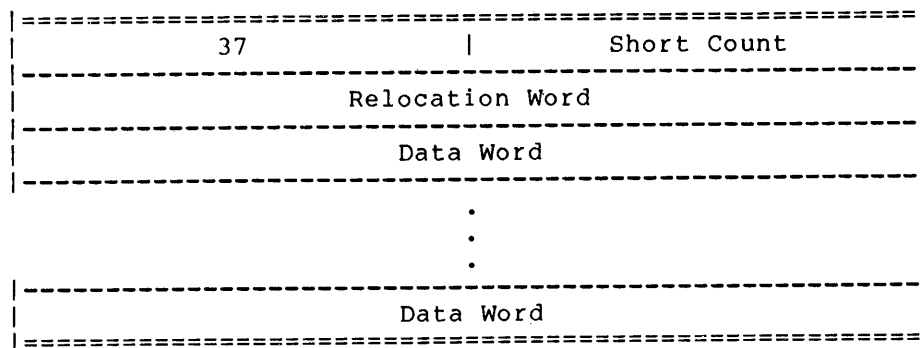
Block Type 24 contains information concerning a specified PSECT. The first word contains the block type number and the number of words associated with the block. The second word contains the relocation information. The third word contains the PSECT name in SIXBIT. The fourth word is the PSECT origin specified for this module.

Bit	Interpretation	MACRO .PSECT Keyword
13	PSECT is page-aligned.	PALIGNED
14	Concatenate parts of PSECTs seen in distinct modules.	CONCATENATE
15	Overlay parts of PSECTs seen in distinct modules.	OVERLAY
16	Read-only	RONLY
17	Read and write	RWRITE

LINK must find a Type 24 block for a PSECT before it finds the index for that PSECT. (MACRO generates a complete set of Type 24 blocks for all PSECTS in a module before generating Type 2 (Symbol Table) Blocks and Type 11 (POLISH) Blocks.)

## REL BLOCKS

### Block Type 37 (COBOL Symbols)



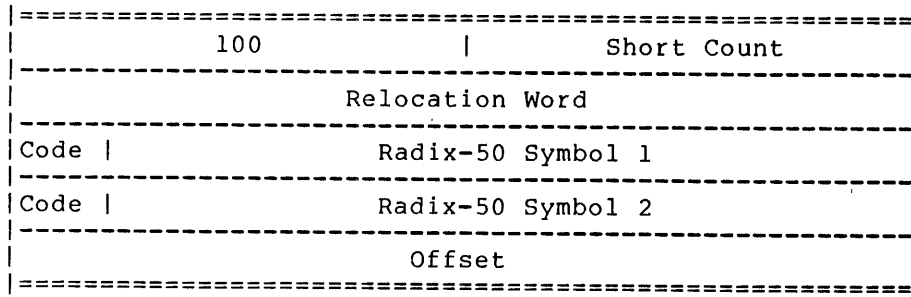
Block Type 37 contains a debugging symbol table for COBDDT, the COBOL debugging program. If local symbols are being loaded, the table is loaded.

If a REL file contains a Block Type 37, it must appear after all other blocks that cause code to be loaded or storage to be allocated in the core image.

This block is in the same format as the Type 1 REL Block.

REL BLOCKS

Block Type 100 (.ASSIGN)



Block Type 100 defines Symbol 1 (in the diagram above) as a new global symbol with the current value of Symbol 2, and then increases the value of Symbol 2 by the value of the given offset.

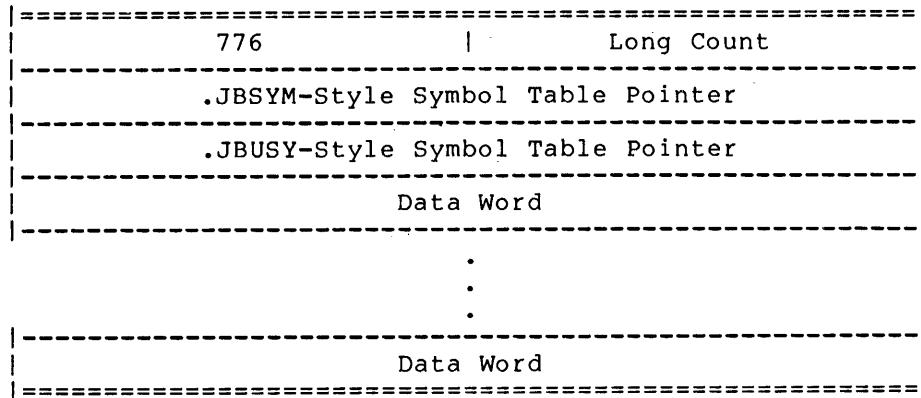
NOTE

Symbol 2 must be completely defined when the Block Type 100 is found.



REL BLOCKS

Block Type 776 (Symbol File)

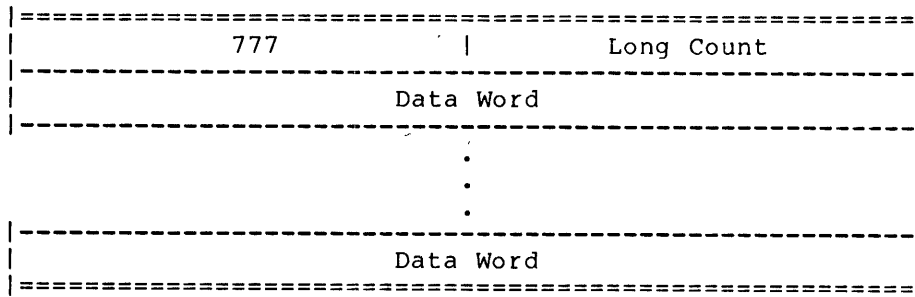


Block Type 776 must begin in the first word of the file, if it occurs at all. This block type shows that the file is a Radix-50 symbol file.

The data words form a Radix-50 symbol table for DDT in the same format as the table loaded for the switches /LOCALS/SYMSEG or the switch /DEBUG.

REL BLOCKS

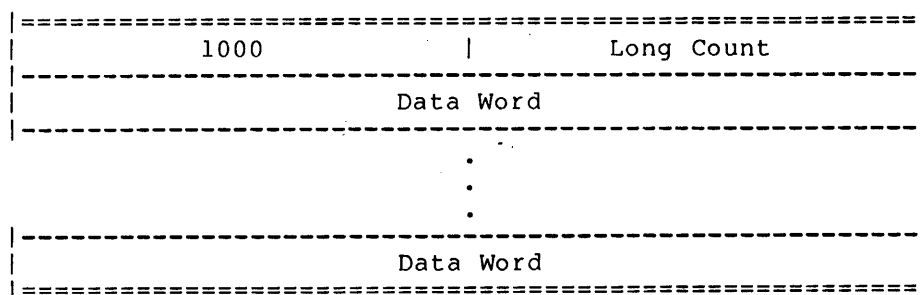
Block Type 777 (Universal File)



Block Type 777 is included in a universal (UNV) file that is produced by MACRO so that LINK will recognize when a UNV file is being loaded inadvertently. When a Block Type 777 is encountered, LINK produces a ?LNKUNS error.

REL BLOCKS

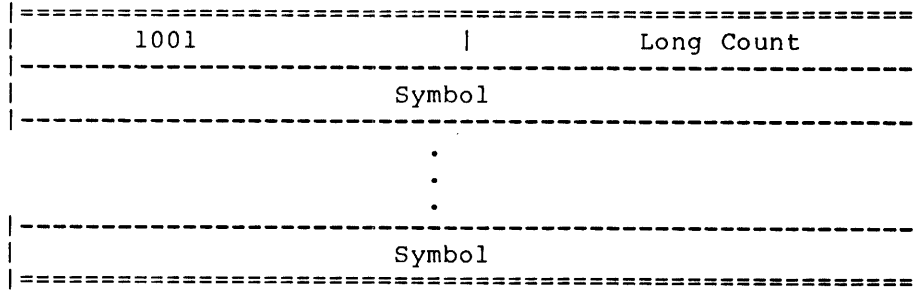
Block Type 1000 (Ignored)



Block Type 1000 is ignored by LINK.

REL BLOCKS

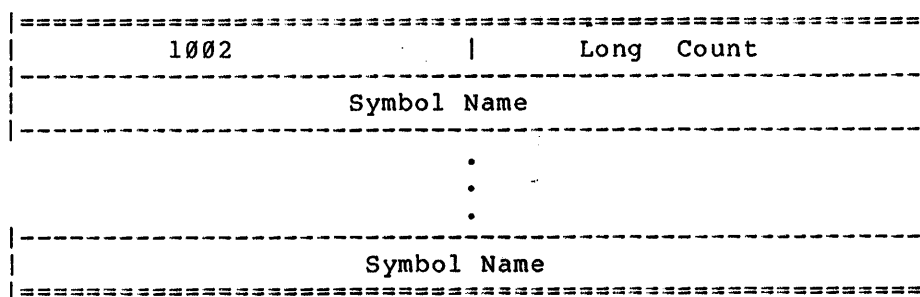
Block Type 1001 (Entry)



Block type 1001 is used to declare symbolic entry points. Each word contains one SIXBIT symbol. This block is similar in function to block type 4.

REL BLOCKS

Block Type 1002 (Long Entry)



Block Type 1002 is used to declare a symbolic entry point with a long name in SIXBIT. The count reflects the symbol length in words.

REL BLOCKS

Block Type 1003 (Long Title)

1003		Long Count
1		Count of Title words
Program Title		
Additional Program Title		
Additional Program Title		
:		
:		
:		
2		Count of ASCII Comment Words
More Comment Words		
More Comment Words		
:		
:		
:		
3		Count of Compiler Words
Compiler Code		CPU Bits
Compiler Name (in ASCII)		
Additional Compiler Name		
Additional Compiler Name		
:		
:		
:		
4		Ø
Compile Date and Time		
Compiler Version Number		

REL BLOCKS

Block Type 1003 (Cont.)

5		0
Device Name		
SIXBIT UFD		
6		0
TOPS-10 File Name		
File Extension		0
7		Number of SFDs
SIXBIT SFD 1		
SIXBIT SFD 2		
.		
.		
.		
10		Count of TOPS-10 File Spec Words (in ASCII)
TOPS-10 File Spec		
TOPS-10 File Spec		
.		
.		
.		
11		0
Source Version Number		
Date and Time		

Block Type 1003 is used to declare long title symbols in SIXBIT and to furnish other information about the source module. This block type contains the information that LINK prints in the map file.

Block Type 1003 consists of sub-blocks 1 through 11 (octal). The Title sub-block must be the first sub-block specified and cannot be omitted. You can omit other sub-blocks, but the sub-blocks must remain in numerical order.

The Program Title is a one word title from 1 to 72 SIXBIT characters long. You can specify a title of 0, and LINK defaults to .MAIN, but you may want to enter a more specific title.

For the compiler code and the CPU code, refer to the explanation of Block Type 6, where these codes are listed.

## REL BLOCKS

Sub-blocks 5 through 7 contain the device name, UFD, and SFD where the file resides.

In sub-block 10, the TOPS-10 file specification must be specified in the following format:

```
[dir]file.ext
```

This specification identifies the source file. LINK outputs this file specification to the map file in the order you enter it.

The Time and Date are in TOPS-10 format. The date is derived from a code that is given by the following formula:

$$\text{code} = 31[12(\text{year}-1964)+(\text{month}-1)]+(\text{day}-1)$$

You can obtain the current day, month, and year using the formulas:

```
day = mod(code,31)+1
month = mod(code/31,12)+1
year = (code/372)+1964
```

The Time is the time in milliseconds that has elapsed since midnight.

See the TOPS-10 Monitor Calls Manual for additional information on specifying date and time.



## REL BLOCKS

### Block Type 1004 (Byte Initialization)

1004		Long Count
Relocation Word		
Byte Count		
Byte Pointer		
Byte String		

·  
·

The above Block Type 1004 format is used to move a character string into static storage. This format uses old style relocation.

The byte count is the number of bytes in the string. The byte pointer is relocated and used to initialize a string in the user's program.

A second format for Block Type 1004 follows:

1004		Long Count
Relocation Word		
Global Symbol		
Byte Count		
Byte Pointer		
Byte String		

·  
·

In this format, the global symbol (in SIXBIT) is used to relocate the byte pointer. The symbol must be defined when this REL block is encountered.

## REL BLOCKS

### Block Types 1010 - 1037 (Code Blocks)

Block Types 1010 through 1037 are similar in function to Block Type 1. They contain code and data to be loaded. These blocks also contain relocation bytes that permit inclusion of PSECT indexes local to the module. For programs that use PSECTs with many inter-PSECT references, this permits a substantial decrease in the size of the REL files. The number of PSECTs that can be encoded in this manner is limited by the size of the relocation byte. A set of parallel code blocks differing only in the size of the relocation byte permits the compiler or assembler to select the most space efficient representation according to the number of PSECTs referenced in a given load module.

This set of blocks is divided by the type of relocation:

Right Relocation	Block Types 1010 - 1017
Left/Right Relocation	Block Types 1020 - 1027
Thirty-bit Relocation	Block Types 1030 - 1037

REL BLOCKS

Block Type 1042 (Request Load for SFDs)

1042		Long Count
Device		
SIXBIT Filename		
File Extension		Directory Count
Project-Programmer Number		
SFD1		
SFD2		
.		
.		
.		

Block Type 1042 contains a list of files to be loaded. It is similar to blocks of Type 16, but it supplies TOPS-10 sub-file directories for the files being requested. The first three data words (device, file name, and extension) are required. The right half of the third word (directory count) specifies the number of directory levels that are included. For example, the directory [27,5434,SFD1,SFD2] would have a directory count of 3.

LINK saves the specifications for the files to be loaded, discarding duplicates. LINK loads all specified files at the end of loading, and immediately before beginning library searches.

REL BLOCKS

Block Type 1043 (Request Library for SFDs)

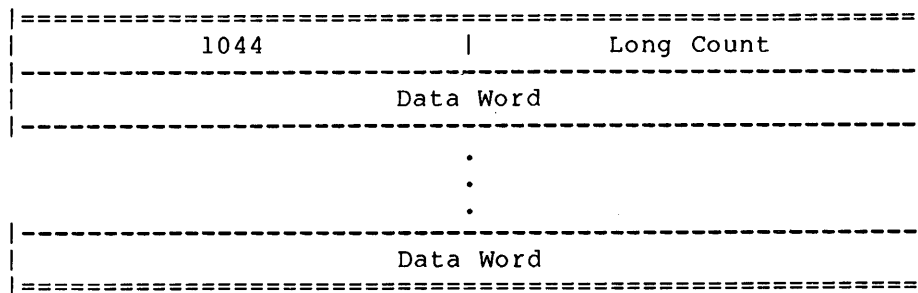
=====	
1043	Long Count
-----	
Device	
-----	
SIXBIT Filename	
-----	
File Extension	Directory Count
-----	
Project-Programmer Number	
-----	
SFD1	
-----	
SFD2	
-----	
.	
.	
.	

Block Type 1043 specifies the files to be searched as libraries. It is similar to Type 17 Blocks, except that it provides TOPS-10 sub-file directories. The first three data words (device, file name, and extension) are required. The right half of the third word (directory count) specifies the number of directory levels that are included. For example, the directory [27,5434,SFD1,SFD2] would have a directory count of 3.

The specified files are searched after requested files are loaded, but before user and system libraries are searched.

REL BLOCKS

Block Type 1044 (ALGOL Symbols)



Block Type 1044 contains a debugging symbol table for ALGDDT, the ALGOL debugging program.

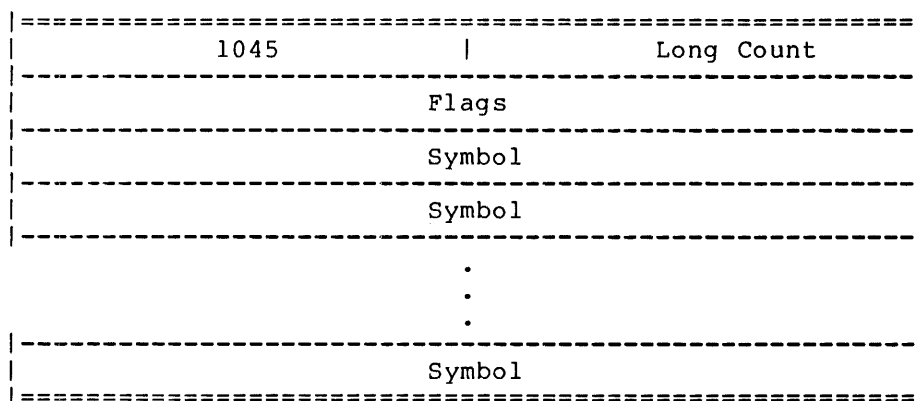
If an ALGOL main program has been loaded, or if you have used the /SYFILE:ALGOL switch, LINK writes the data words into a SYM file. In addition, if any Type 15 (ALGOL OWN) REL blocks have been seen, LINK stores the file specification of the file into the first OWN block loaded.

NOTE

If you have specified the /NOSYMBOLS switch, or if you have specified the /SYFILE switch with an argument other than ALGOL, then LINK ignores any Type 1044 blocks found.

## REL BLOCKS

### Block Type 1045 (Writable Links)

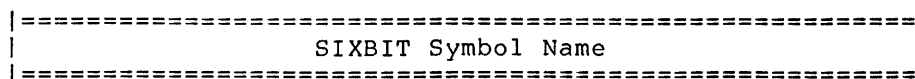


Block type 1045 declares as writable either the link containing the current module or the links containing the definitions of the specified symbols or both. This block type must follow any common block declarations (Types 20 or 6) in a module.

The flag word indicates which links are writable. If bit one is set then the link containing the current module and the links containing the definitions of the specified symbols are writable. If bit one of the flag word is not set then the link containing the current module is not writable, but the links containing the specified symbols are writable. All unused flag bits are reserved and should be zero.

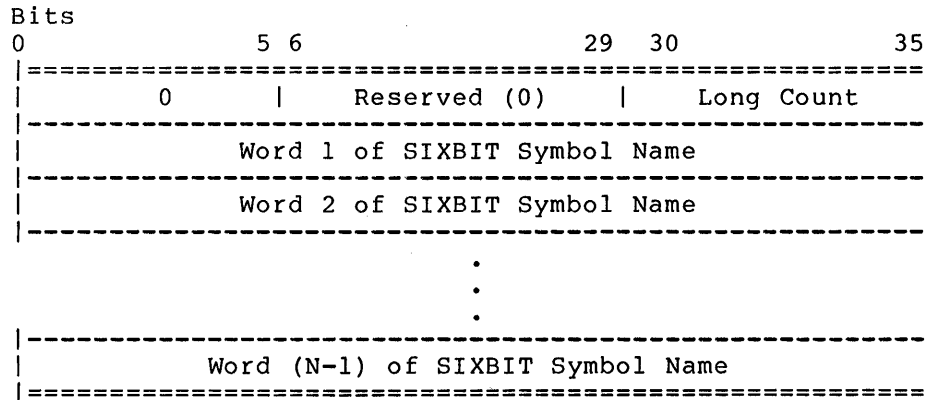
Any symbols specified in a block of Type 1045 must be defined in the path of links leading from the root link to the current link. A module cannot declare a parallel or inferior link to be writable.

If the symbol name contains six or fewer characters it is represented in a single word, left justified, with the following format:



## REL BLOCKS

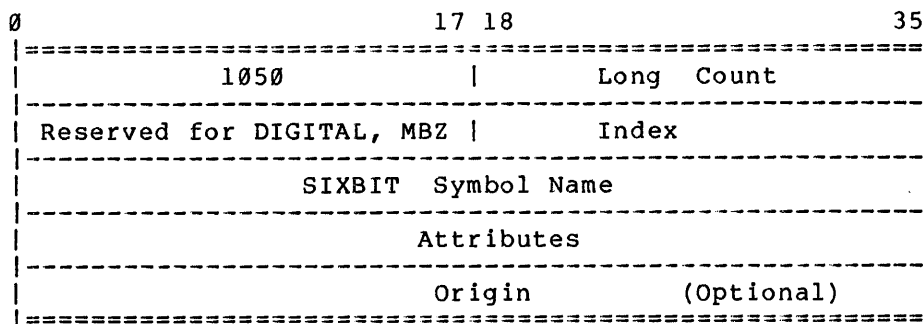
If the symbol name contains more than six characters it is represented in the following format:



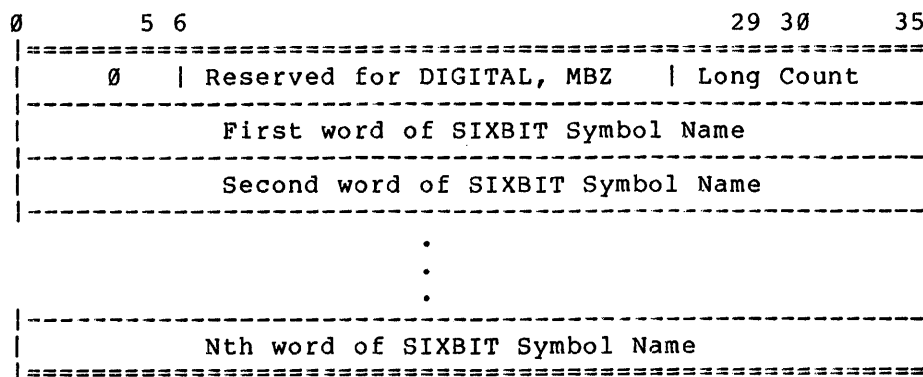
The first six bits of a long symbol are always 0. This distinguishes a long symbol name from a single word symbol name. N is the length of the symbol name including the header word. The remaining words contain the symbol name in SIXBIT, six characters to a word, left justified.

REL BLOCKS

Block Type 1050 (Long PSECT Name Block)



where SIXBIT Symbol Name may be either a word of up to six SIXBIT characters, or the following block.



Block Type 1050 creates a PSECT with the given name, if none currently exists. It also assigns a unique index number to the PSECT. This index is binding only in the current module. LINK clears PSECT indexes at the end of each module. PSECT indexes in any given module must be declared in consecutive order starting at index 1.

Block Type 1050 also assigns attributes to a PSECT and specifies the PSECT's origin address. The attributes that can be assigned are:

Bit	Description
11	PSECT is confined to one section. If this bit is set, LINK gives an error if the PSECT overflows. You can set Bit 11 or Bit 12, but not both. Bit 11 is the default. There is no equivalent MACRO .PSECT keyword.
12	PSECT is in a nonzero section. If this bit is set, LINK gives a warning if the PSECT is placed in Section 0. There is no equivalent MACRO .PSECT keyword.
13	PSECT is PAGE-ALIGNED. PALIGNED is the equivalent MACRO .PSECT keyword.



## REL BLOCKS

- | Bit | Description  |
|-----|--|
| 14  | <p>CONCATENATED parts of PSECTs seen in distinct modules.</p> <p>You can set Bit 14 (CONCATENATED) or Bit 15 (OVERLAID), but not both. The CONCATENATE and OVERLAID attributes are mutually exclusive. These attributes span modules; so if one module sets an attribute and a later module sets a mutually exclusive attribute, LINK issues the warning:</p> <pre>%LNKCOE Both CONCATENATE and OVERLAY attributes specified for psect [name].</pre> <p>If neither is set, CONCATENATED is the default, and a warning message is not returned if subsequent pieces of the PSECT are marked OVERLAID.</p> <p>CONCATENATED is the equivalent MACRO .PSECT keyword.</p> |
| 15  | <p>OVERLAID parts of PSECTs seen in distinct modules. OVERLAID is the equivalent MACRO .PSECT keyword.</p>   |
| 16  | <p>This PSECT must be READ-ONLY.</p> <p>You can set Bit 16 (READ-ONLY) or Bit 17 (WRITABLE), but not both. The READ-ONLY and WRITABLE attributes are mutually exclusive. These attributes span modules; so if one module sets an attribute and a later module sets a mutually exclusive attribute, LINK issues the warning:</p> <pre>%LNKRWA Both READ-ONLY and WRITABLE attributes specified for psect [name].</pre> <p>If neither is set, WRITABLE is the default, and a warning message is not returned if subsequent pieces of the PSECT are marked READ-ONLY.</p> <p>RONLY is the equivalent MACRO .PSECT keyword.</p>  |
| 17  | <p>This PSECT must be WRITABLE. RWRITE is the equivalent MACRO .PSECT keyword.</p>   |

All other bits in the Attributes word must be 0.

The origin specified in this block is absolute.

At least one Block Type 1050 (or the related Block Type 24) is required for each PSECT being loaded, and this block must be loaded prior to any other blocks that reference its PSECT (that is, use the unique index number).

REL BLOCKS

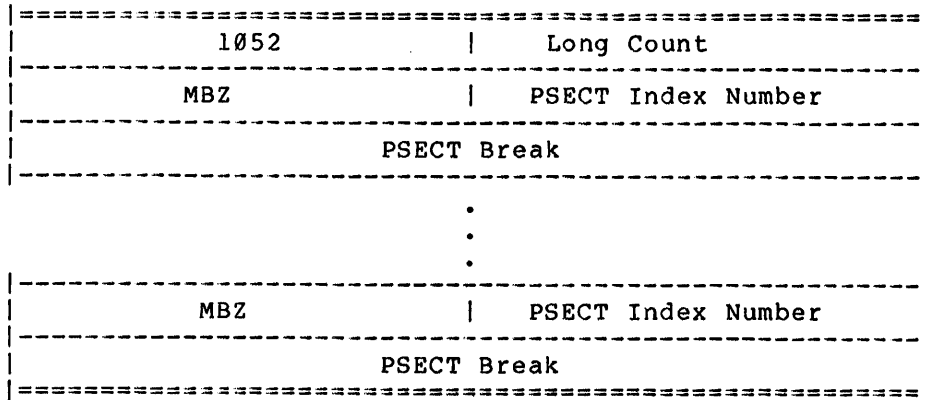
Block Type 1051 (Set Current PSECT)

1051	Long Count
Reserved for DIGITAL, MBZ	Index

Block Type 1051 resets the "current PSECT", against which LINK relocates subsequent REL blocks if no PSECT is explicitly specified.

REL BLOCKS

Block Type 1052 (PSECT End)



Block Type 1052 allocates additional space for a given PSECT. This space is located between the last address in the PSECT containing data and the address given by the PSECT break. A block of Type 1052 can contain more than one pair of PSECT indexes and breaks.

A module must contain a block of Type 24 (PSECT Name) or Type 1050 (Long PSECT Name) with the given PSECT index before a block of Type 1052 is generated. If a given PSECT has more than one Block Type 1052 in a single module, the block with the largest break address is used.

The break is interpreted as being relative to the PSECT's origin in the current module.

REL BLOCKS

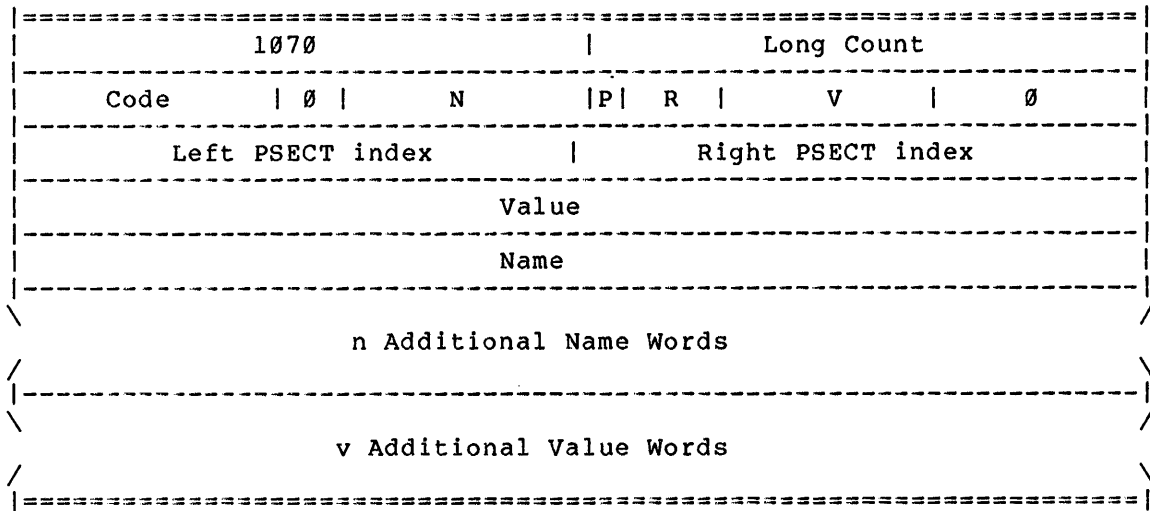
Block Type 1060 (Trace Block Data)

1060		Long Count
SIXBIT Edit Name		
Active Code		Last Changer
Creator Code		15-Bit Date Created
Installer Code		15-Bit Date Installed
Reserved		
Edit Count		PCO Group Count
Associated Edit Names and Codes		
Program Change Order Groups		

Block Type 1060 contains data used by the MAKLIB program. LINK ignores this block type.

REL BLOCKS

Block Type 1070 (Long Symbol Names)



This block defines a long symbol. A symbol defined with this block can:

- o be written to the DDT symbol table. Symbols longer than 6 characters are truncated when output to the DDT symbol table.
- o be written to the map file if requested.
- o have its value relocated as specified.
- o resolve global requests.

The Long Symbol Name Block is divided into two sections, the basic and the extension sections.

The basic section consists of four words: the flag word, an optional PSECT index word, the value word, and name word.

The Flag word contains information about the type of symbol, the length of the symbol name, and relocation. The optional word defines the PSECT index. The Value word contains the symbol's value. The Name word contains the symbol's name.

If the name or the value cannot fit in a single word, the block contains an extension section that consists of as many words as are necessary to accommodate the symbol name and the value. The length of the symbol name and value is stored in the Flag word and determines how many words are allocated for the long symbol name in the extension section. The maximum size of the symbol is 72 characters. In the case of a short symbol name, only the basic section is used.

The following pages provide detailed information on the block. For each word, the field, bits, and description is given.

## REL BLOCKS

### Header Word

Field	Bits	Description
Block Type	0-17	1070
Block Length	18-35	Number of words used in this block

### Flag Word

Field	Bits	Description
Code	0-8	A nine-bit code field:
		Bit 0 Must Be Zero
	000	Program name
	100	Local symbol definitions
	110	Suppressed to DDT
	120	MAP only
	200	Global symbols completely defined by one word
	202	Undefined
	203	Right fixup
	204	Left fixup
	205	Right and left fixups
	206	30-bit fixup
	207	Fullword fixup
	210-217	Suppress to DDT
	220-227	MAP only
	240-247	Global symbol request for chain fixup
	240	Ignored (no fixup)
	241	Undefined
	242	Undefined
	243	RH fixup
	244	LH fixup
	245	Undefined
	246	30-bit fixup
	247	Fullword fixup
	250-257	Global request for additive fixups (the value of x has the same meaning as in 0-7 above)
	260-267	Global request for additive symbol fixups (the value of x has the same meaning as in 0-7 above)
	300	Block names

#### NOTE

All symbols that require a fixup for their definition must have the fixup block immediately following the entry.

## REL BLOCKS

Field	Bits	Description												
Ø	9-1Ø	Must Be Zero												
N--Name length	11-17	If not zero, extended name field of length n words is used, so that the name occupies n+1 words.												
P--PSECT Flag	18	If Bit 18=Ø, relocate with respect to the current PSECT. No PSECT numbers are needed.  If Bit 18=1, relocate with respect to the PSECT specified in the next word.												
R--Relocation Type	19-21	3-bit relocation type field.  <table border="0" style="margin-left: 40px;"> <tr><td>Ø</td><td>Absolute</td></tr> <tr><td>1</td><td>Right half</td></tr> <tr><td>2</td><td>Left half</td></tr> <tr><td>3</td><td>Both halves</td></tr> <tr><td>4</td><td>3Ø-bit</td></tr> <tr><td>5</td><td>Fullword</td></tr> </table>	Ø	Absolute	1	Right half	2	Left half	3	Both halves	4	3Ø-bit	5	Fullword
Ø	Absolute													
1	Right half													
2	Left half													
3	Both halves													
4	3Ø-bit													
5	Fullword													
V--Value field	22-28	Number of additional value words if value is a long symbol.												
Ø	29-35	Not used												

### PSECT Indexes

PSECT Indexes                      Exists only if Bit 18 equals 1 in the Flag word. Contains Left and Right PSECT numbers. Bit Ø and Bit 18 of this word are zeros.

### Value

Value Word                              Contains the symbol value. This may be relocated as specified by the relocation type and the PSECT numbers provided. Contains a symbol for 26x codes.

### Name

Name Word                                Contains the symbol name in SIXBIT.

### N Additional Name Words

Additional name field                      Optional. It exists only if N > Ø. It contains the additional characters when a long symbol name is used.

### V Additional Value Words

Additional value field                      Optional. It exist only if the V field is greater than Ø. This field contains the additional characters when a long symbol name is being resolved. The first word contains the length of the extended field.

## REL BLOCKS

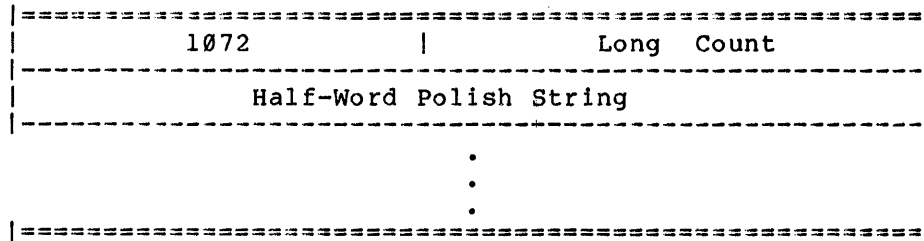
The following fixup rules apply to this block:

- Only one fixup by a Type 2, 10, 11, 12, 15, 1070, 1072, or 1120 Block is allowed for a given field. There can be separate fixups for the left and right halves of the same word.
- Fixups are not necessarily performed in the order LINK finds them.
- Chained halfword fixups cannot cross section boundaries; they wrap around to the beginning of the section. Also, they cannot fixup a location that resolves to word zero of a section unless it is the only address in the chain.
- Chained fixups must be in strict descending address order.
- A location must contain data before the location can be fixed up.



## REL BLOCKS

### Block Type 1072 (Long Polish Block)



Long Polish Blocks of Type 1072 define Polish fixups for operations on relocatable long external symbols. This Block Type is interpreted as a string of 18-bit operators and operands. The block is in Polish prefix format, with the store operator at the end of the block. Each halfword can contain one of the following:

- A halfword code in which the first 9 bits contain the data length (when applicable) and the second 9 bits contain the code telling LINK how to interpret the data that follows.
- A halfword data or a part of a larger data packet to be interpreted by LINK as indicated by the code that immediately precedes it.
- A PSECT index of the format 400000+n. The PSECT index field of a long Polish block causes LINK to relocate addresses against the PSECT number specified in the "n" of the PSECT index 400000+n.
- A Polish operator.

#### NOTE

Operations are performed in the order in which they are encountered.

#### CODE DEFINITIONS

##### Data Packet Codes

Category	Code	Description
Operand	xxxyyy	next "xxx+1" halfwords contain data of type "yyy"
	000000	halfword - absolute
	001000	fullword - absolute
	000001	halfword - relocatable
	001001	fullword - relocatable
	000010	fullword symbol name in Radix-50
	xxx010	xxx+1 halfwords of symbol name in SIXBIT

#### NOTE

You cannot store a symbol in a single halfword. You must place the symbol in the first halfword and fill the second halfword with zeroes.

## REL BLOCKS

### Polish Operator Codes

Category	Code	Description
Operator	000100	Add
	000101	Subtract
	000102	Multiply
	000103	Divide
	000104	Logical AND
	000105	Logical OR
	000106	Logical shift
	000107	Logical XOR
	000110	One's complement (not)
	000111	Two's complement (negative)
	000112	Count leading zeros
	000113	Remainder
	000114	Magnitude
	000115	Maximum
	000116	Minimum
	000117	Equal relation
	000120	Link
	000121	Defined
	000122-00177	Reserved

### Store Operator Codes

Store Operator	xxx=0 or 1	
	For xxx=0	Next two halfwords contain a Radix-50 symbol to be resolved.
	xxx777-xxx770	Chained fixup with relocatable addresses. Next xxx+1 halfwords contain the start address of the chain.
	000777	Right half chained fixup with relocatable address. Next halfword contains a relocatable address.
	000776	Left half chained fixup with relocatable address. Next halfword contains a relocatable address.
	000775	30-bit chained fixup with relocatable address. Next halfword contains a relocatable address.
	000774	Fullword chained fixup with relocatable address. Next halfword contains a relocatable address.
	001777	Right half chained fixup with relocatable address. Next fullword contains a relocatable address.
	001776	Left half chained fixup with relocatable address. Next fullword contains a relocatable address.

## REL BLOCKS

001775	30-bit chained fixup with relocatable address. Next fullword contains a relocatable address.
001774	Fullword chained fixup with relocatable address. Next halfword contains a relocatable address.
xxx767-xxx764	Chained fixups with absolute addresses.
000767-000764	Chained fixup with absolute address. Next halfword contains an absolute address.
001767-001764	Chained fixup with absolute fullword address. Next two halfwords contain absolute address.
xxx757-xxx754	Symbol fixup. For $1 \leq xxx \leq 377$ the next $xxx+1$ halfwords contain a SIXBIT symbol name to be resolved.
xxx757	Right half symbol fixup.
xxx756	Left half symbol fixup.
xxx755	30-bit symbol fixup.
xxx754	Fullword symbol fixup.
xxx747-xxx700	Not defined
PSECT index 4000000+n	PSECT index for PSECT N.

The following fixup rules apply to this block:

- Only one fixup by a Type 2, 10, 11, 12, 15, 1070, 1072, or 1120 Block is allowed for a given field. There can be separate fixups for the left and right halves of the same word.
- Fixups are not necessarily performed in the order LINK finds them.
- Chained halfword fixups cannot cross section boundaries; they wrap within a section. Also, they cannot fixup a location that resolves to word zero of a section unless it is the only address in the chain.
- Chained fixups must be in strict descending address order.
- A location must contain data before the location can be fixed up.

REL BLOCKS

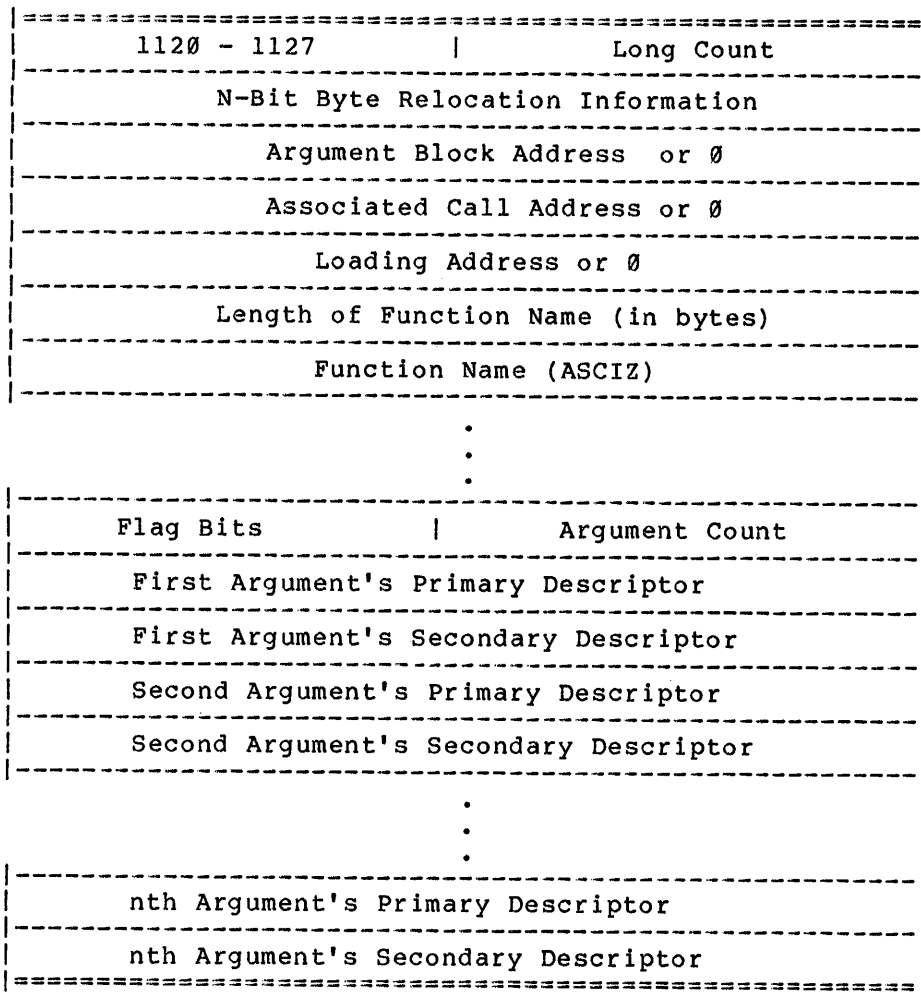
Block Type 1074 (Long Common Name)

1074		Long Count
PSECT Index		Symbol Length
Common Block Length		
Symbol (More Symbol)		

Block Type 1074 defines a long COMMON name.

## REL BLOCKS

Block types 1120-1127 (Argument Descriptor Blocks)



A block of this type is generated for the argument list to each subroutine call. The subroutine entry point also specifies one block with this format, though for the callee the argument block address is zero. If a descriptor block is associated with an argument list it must always follow the loading of the argument list.

The associated call address is used by LINK in diagnostic error messages and its value is determined by the compiler. The argument block address is nonzero if the descriptor block is associated with a call. In this case the argument block address points to the base of the argument block.

The argument block address, associated call address and the loading address are all relocatable.

The argument descriptors in these type blocks describe the properties of each formal (in the case of an entry point) or actual (in the case of a call). In either case the name of the associated routine is specified as a byte count, that can only be 6 characters long, followed by an ASCIZ string. Each primary description is optionally followed by a secondary descriptor.

## REL BLOCKS

There are five flag bits in the Descriptor Block:

Bit	Usage
0	If bit 0 is 1 then a difference between the actual number of arguments and the expected number of arguments is flagged as a warning at load time. If bit 0 is 0 no action is taken.
1	If bit 1 is 1 then the block is associated with a function call. If bit 1 is 0 then the block is associated with the function definition.
2	If bit 2 is 1 then the descriptor block is loaded into user memory at the loading address. This bit is ignored.
3	If bit 3 is 1 then the callee returns a value and the value's descriptor is the last descriptor specified.
4	If bit 4 is 1, and the caller expects a return value, which is not provided by the called function, or if the called function unexpectedly returns a value, then LINK will issue an error. The severity of the error is controlled by the coercion block.

The format for the argument descriptors is as follows:

Bit	Usage
0	(Reserved)
1	No update. In a caller block the argument is a literal, constant, or expression. In a callee block the argument won't be modified.
2-4	Passing mechanism 000 - pass by address 001 - pass by descriptor 010 - pass immediate value Others - reserved
5	Compile-time constant
6-11	Argument type code (see below)
12-17	(Reserved)
18	Implicit argument descriptor
19-26	(Reserved)
27-35	Number of secondary descriptors

## REL BLOCKS

The argument type codes are as follows:

Type-Code	Usage
0	
1	FORTRAN logical
2	Integer
3	(Reserved)
4	Real
5	(Reserved)
6	36-bit string
7	Alternate return (label)
10	Double real
11	Double integer
12	Double octal
13	G-floating real
14	Complex
15	COBOL format byte string descriptor, (for constant strings), or FORTRAN character
16	BASIC shared string descriptor
17	ASCIZ string
20	Seven-bit ASCII string

Secondary descriptors are used to convey information about the length of a data object passed as an argument and (in the case of the callee's argument descriptor block) whether or not a mismatched length is permissible. Secondary descriptors have the following format:

Bit Pos.	Usage
0-2	(For callee only) Defines the permissible relationships between formal and actual lengths. The values are:  000 - Any relationships are allowed 001 - Lengths must be equal 010 - Actual < formal 011 - Actual <= formal 100 - Actual > formal 101 - Actual >= formal 110 - Reserved 111 - Reserved
3-5	Length of argument (in words)

REL BLOCKS

Block Type 1130 (Coercion Block)

1130		Long Count
Field Code		Action
Formal Attribute		Actual Attribute
Field Code		Action
Formal Attribute		Actual Attribute
.		.
.		.
Field Code		Action
Formal Attribute		Actual Attribute

Block Type 1130 specifies which data type associations are permissible and what action LINK should take if an illegal type association is attempted. It may also specify actions to be taken by LINK to modify an actual parameter.

The Coercion Block must be placed before any instance of the caller/callee descriptor block in the REL file. If more than one coercion block is seen during a load, the last block seen is used for type checking.

If the descriptor block and command strings are not in the same section, no error message is given.

When a caller's argument descriptor block is compared to the descriptor block provided by the callee, LINK first checks bit 0 and the argument counts of the descriptor block. If bit 0 is set and the argument counts differ, a warning is given. However, if a byte descriptor is not word-aligned, no warning is given.

Next LINK compares the argument descriptors. The particular formal/actual pair is looked up in the internal table LINK builds using the information in the coercion block. The item field code designates which field of the argument descriptor is being checked. The field codes are defined as follows:

Field Code	Condition
0	Check update
1	Check passing mechanism
2	Check argument type code
3	Check if compile-time constant
4	Check number of arguments
5	Check for return value
6	Check length of argument



## REL BLOCKS

If the fields of the formal/actual pair do not match, LINK searches the internal table set up by the coercion block. If the table does not specify an action to take in the event of such a mismatch, LINK issues an informational message. If the formal/actual pair differs in more than one field then LINK takes the most severe action specified for the mismatches.

If an actual/formal pair differ and no coercion block has been seen, LINK ignores the difference. If the caller has specified a descriptor block but the subroutine has not, or if the subroutine has specified a descriptor and the caller has not, LINK does not flag the condition as an error and does not take any special action.

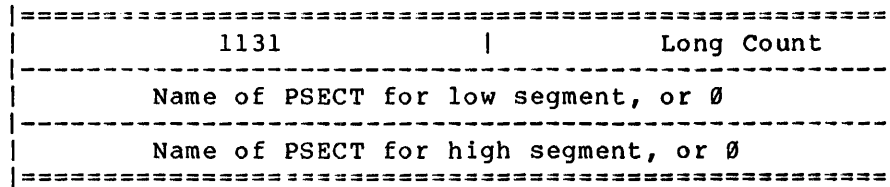
If LINK finds an entry in its internal table for a particular actual/formal mismatch, it uses the action code found in the entry to select one of the following five possible responses:

Code (18 Bits)	Action
0	Informational message
1	Warning
2	Error
3	Reserved for the specific conversion of static descriptor pointers (in the argument list) into addresses. The descriptor pointers are supplied by FORTRAN blocks of types 112x.
NOTE	
The actual conversion process involves the following actions:	
<ul style="list-style-type: none"><li>• If byte descriptor's P field is not word-aligned, issue a warning and continue.</li><li>• Pick up word address of start of string.</li><li>• Put the address of the string into the associated argument block in place of the address of the string descriptor.</li></ul>	
4	Suppress the message.
5-777776	Reserved
777777	Fatal error

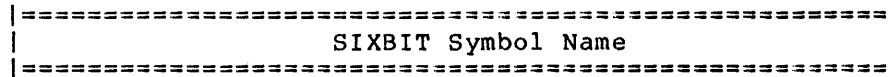
These messages can be displayed or suppressed. Refer to the descriptions of the /ERRORLEVEL and /LOGLEVEL switches.

REL BLOCKS

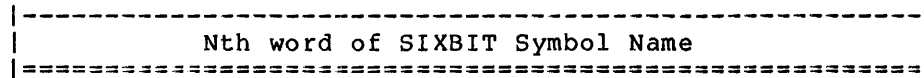
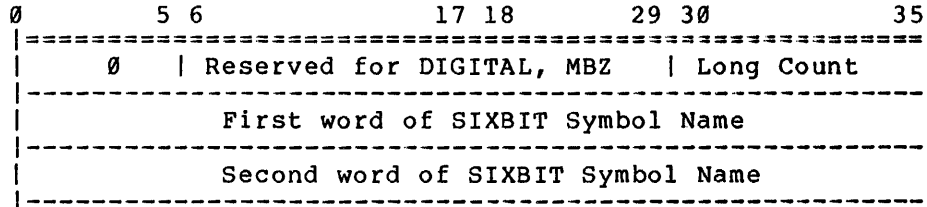
Block Type 1131 (TWOSEG Redirection Block)



where each PSECT name has the form:



or

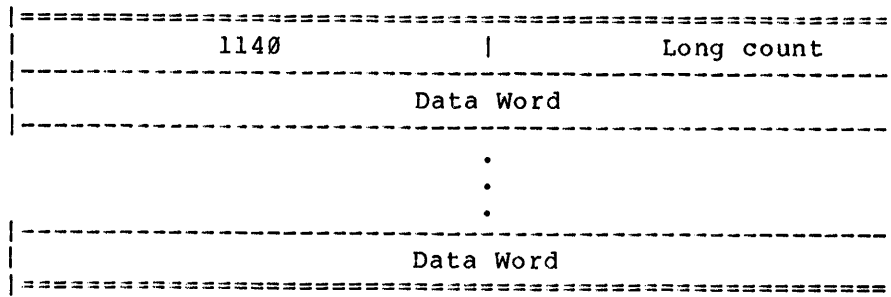


Block Type 1131 permits TWOSEG REL modules to be loaded into PSECTs by a compiler. You must redirect both the high and the low segment, you cannot redirect one or the other. Also, you cannot redirect both the high and low segment into the same PSECT.

This block does not affect the current module, but all subsequent modules to be loaded.

REL BLOCKS

Block Type 1140 (PL/1 debugger information)



Block type 1140 is ignored by LINK.

REL BLOCKS

Block Type 1160 (Extended Sparse Data Initialization Block)

1160						Long	Count
R	F	B	P	0	SYMLEN	PSECT	
Symbol (SYMLEN words)							
S	Origin Address						
Repetition Count if R=1							
Fill Count if F=1							
Fill Byte if F=1							
Byte Count if B=1							
Data Bytes							

Block Type 1160 supports the loading of data into different PSECTS and sections. This REL block allows separate program units to load data into different bytes in the same word of memory at different times during the loading process.

Block Type 1160 fields are described below.

Field Name	Position	Description
R	Bit 0	is a 1-bit field. If R is one, the Repetition Count word exists. If R is zero, the Repetition Count is assumed to be 1.
F	Bit 1	is a 1-bit field. If F is one, the Fill Count and Fill Byte words exist. If F is zero, no fill is used.
B	Bit 2	is a 1-bit field. If B is one, the Byte Count word exists. If B is zero, one Data Byte is assumed.
P	Bits 3-8	is a 6-bit field. This is the position within the word where the first byte is to be stored.
Unused	Bit 9	is an unused bit that must be zero.

## REL BLOCKS

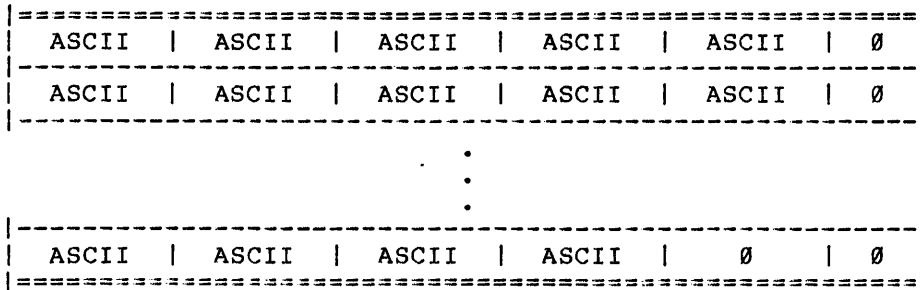
Field Name	Position	Description
SYMLEN	Bits 10-17	is an 8-bit field. SYMLEN is the length in words of the global symbol to be used to calculate the address to store the byte string. If SYMLEN is zero, there is no global symbol. The value of the symbol is added to the origin address. The symbol must be completely defined before this addition occurs.
PSECT	Bits 18-35	is an 18-bit field. PSECT is the PSECT to relocate the Origin Address against. The relocation is 30-bit. If PSECT is zero, the Origin Address is absolute.
Symbol	Bits 0-35	is a SIXBIT symbol name of the length specified in SYMLEN. The value of this symbol is added to the Origin Address. The symbol must be defined when the block is seen, or a fatal error occurs.
S	Bits 0-5	is a 6-bit field. S is the size of the data bytes.
Origin Address	Bits 6-35	is a 30-bit field. Origin Address is the address where LINK begins to store Data Bytes.
Repetition Count	Bits 0-35	is a 36-bit field. If flag bit R is one, Repetition Count exists and contains the number of times to repeat the data store. The Data Bytes are stored and the fill operation is performed as many times as specified in the Repetition Count.
Fill Count	Bits 0-35	is a 36-bit field. If flag bit F is one, Fill Count exists and specifies how many times to store the Fill Byte after storing Data Bytes.

## REL BLOCKS

Field Name	Position	Description
Fill Byte	Bits 0-35	is a 36-bit field. If flag bit F is one, Fill Byte exists and contains the right justified value to be used in the fill operation.
Byte Count	Bits 0-35	is a 36-bit field. If flag bit B is one, Byte Count exists and specifies the number of Data Bytes to be stored.
Data Bytes	Bits 0-35	are the data to be stored, of the length specified by the Byte Count, or 1 if flag bit B is not set. This data is stored left-justified, packed as many to a word as fit without overlapping a word boundary.

REL BLOCKS

Block Type Greater Than 3777 (ASCIZ)



When LINK reads a number larger than 3777 in the left half of a REL Block header word, the block is assumed to contain ASCIZ text. If the module containing the text is being loaded, LINK reads the ASCII characters as if they were a command string, input from the user's terminal.

LINK reads the string as five 7-bit ASCII characters per word; bit 35 of each word is ignored. The string and the block end when the first null ASCII character (000) is found in the fifth 7-bit byte of a word (bits 28-34).

After loading the current REL file, LINK processes text statements in the reverse order in which they are encountered -- from the end to the beginning of a module. For example, the first, second, and third statements from the beginning of a module are processed third, second, and first. As a result, search requests may be processed in the reverse order of entered /SEARCH switches. Keep this in mind when specifying the order the modules are to be searched.





APPENDIX B  
LINK MESSAGES

This appendix lists all of LINK's messages. (The messages from the overlay handler, which have the OVL prefix, are given in Chapter 5.)

**B.1 DESCRIPTION OF MESSAGES**

Section B.2 lists LINK's messages. For each message, the last three letters of the 6-letter code, the level, the severity, and its medium-length message are given in **boldface**. Then, in lightface type, comes the long message.

When a message is issued, the three letters are suffixed to the letters LNK, forming a 6-letter code of the form LNKxxx.

The level of a message determines whether it will be issued to the terminal, the log file, or both. You can use the /ERRORLEVEL and /LOGLEVEL switches to control message output. For some messages an asterisk (\*) is given for the level or severity. This means that the value is variable, and depends on the conditions that generated the message.

The severity of a message determines whether the load will be terminated when the message is issued. Table B-1 lists the severity codes used in LINK, along with their meanings. The /SEVERITY switch provides a means for lowering the severity that is considered fatal.

The severity also determines the first character on the message line printed to the terminal. This character can then be detected by the batch system. For all informational messages, the character is [. Warnings use %, and fatal errors use ?.

## LINK MESSAGES

Table B-1  
Severity Codes

Code	Meaning
1-7	Informational; messages of this severity generally indicate LINK's progress through the load.
8-15	Warning; LINK is able to recover by itself and continue the load.
16	Warning if timesharing, but fatal and stops the load if running under batch.
20	Fatal; LINK can only partially recover and continue the load. The loaded program may be incorrect. Undefined symbols cause this action.
24	This is for file access errors. Under batch, this is fatal and stops the load. Under timesharing, this is a warning, and LINK prompts for the correct file specification if possible.
31	Always fatal; LINK stops the load.

The /VERBOSITY switch determines whether the medium-length and long messages are issued. If you use /VERBOSITY:SHORT, only the 6-letter code, the level, and the severity are issued. If you use /VERBOSITY:MEDIUM, the medium-length message is also issued. If you use /VERBOSITY:LONG, the code, level, severity, medium-length message, and long message are issued.

Those portions of the medium-length messages enclosed in braces ( { and } ) are optional, and are only printed in appropriate circumstances.

Those portions of the medium-length messages enclosed in square brackets are filled in at runtime with values pertinent to the particular error. Table B-2 describes each of these bracketed quantities.

## LINK MESSAGES

Table B-2  
Special Message Segments

[area]	The name of one of LINK's internal memory management areas.
[date]	The date when LINK is running.
[decimal]	A decimal number, such as a node number.
[device]	A device name.
[file]	A file specification.
[label]	An internal label in LINK.
[memory]	A memory size, such as 17P.
[name]	The name of the loaded program or a node in an overlaid program.
[octal]	An octal number, such as a symbol value.
[reason]	The reason for a file access failure, one of the messages shown in Section B.3.
[switch]	The name of a switch associated with the error.
[symbol]	The name of a symbol, such as a subroutine or common block name.
[type]	The type or attributes associated with a symbol.

The octal status codes that appear in TOPS-10 messages are described more thoroughly in the TOPS-10 Monitor Calls Manual.

Whenever possible, LINK attempts to indicate the module and file associated with an error. This information represents the module currently being processed by LINK, and may not always be the actual module containing the error. For instance, if LINK detects a multiply-defined symbol, either value may be the incorrect one. In this case, LINK reports only the second (and subsequent) redefinition and the module containing it.

## LINK MESSAGES

### B.2 LIST OF MESSAGES

Code	Lev	Sev	Message
ABT	31	31	<p>Load aborted due to %LNKTMA errors, max. /ARSize: needed was [decimal]</p> <p>You loaded programs containing more ambiguous subroutine requests than can fit in the tables of one or more overlay links. You received a LNKARL message for each ambiguous request, and a LNKTMA message for each link with too many requests. You can solve this problem by using the /ARSize switch just before each /LINK switch to expand the tables separately.</p>
AIC	31	31	<p>Attempt to increase size of {blank common} {common [symbol]} from [decimal] to [decimal] {Detected in module [symbol] from file [file]}</p> <p>FORTTRAN common areas cannot be expanded once defined. Either load the module with the largest definition first, or use the /COMMON: switch to reserve the needed space.</p>
AMM	‡	‡	<p>Argument mismatch in argument [decimal] in call to routine [symbol] called from module [symbol] at location [octal]</p> <p>The caller supplied argument does not match the argument expected by the callee.</p>
AMP	8	8	<p>ALGOL main program not loaded</p> <p>You loaded ALGOL procedures, but no main program. The missing start address and undefined symbols will cause termination of execution.</p>
ANM	31	31	<p>Address not in memory</p> <p>LINK expected a particular user address to be in memory, but it is not there. This is an internal LINK error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ARL	8	8	<p>Ambiguous request in link [decimal] {name [name]} for [symbol] defined in links [decimal], [decimal], ...</p> <p>More than one successor link can satisfy a call from a predecessor link. The predecessor link requested an entry point that is contained in two or more of its successors. You should revise your overlay structure to remove the ambiguity.</p>

‡ The level and severity of this message is determined by compiler-generated coercion block. See Block Type 1130 in Appendix A.

## LINK MESSAGES

Code	Lev	Sev	Message
			<p>If you execute the current load, one of the following will occur when the ambiguous call is executed:</p> <ul style="list-style-type: none"> <li>• If only one module satisfying the request is in memory, that module will be called.</li> <li>• If two or more modules satisfying the request are in memory, the one with the most links in memory will be called.</li> <li>• If no modules satisfying the request are in memory, the one with the most links in memory will be called.</li> </ul> <p>If a module cannot be selected by the methods 2 or 3 above, an arbitrarily selected module will be called.</p>
AZW	31	31	<p><b>Allocating zero words</b></p> <p>LINK's memory manager was called with a request for 0 words. This is an internal LINK error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
CBO	31	31	<p><b>Cannot build overlays outside section 0.</b></p> <p>You tried to build an overlay structure for a program that is either too large to fit in Section 0 or is loaded outside Section 0 by default. Check the commands input to LINK and the language compiler.</p>
CCD	31	31	<p><b>CPU conflict</b>  {Detected in module [symbol] from file [file]}</p> <p>You have loaded modules compiled with conflicting CPU specifications, such as loading a MACRO program compiled with the statement .DIRECTIVE KL10 and another compiled with .DIRECTIVE KIL0. Recompile the affected modules with compatible CPU specifications.</p>
CCE	8	8	<p><b>Character constant not word aligned in call to routine [routine] called from module [module] at location [address]</b></p> <p>LINK detected a character constant that did not begin on a word boundary. This error probably results from a language translator error. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>

LINK MESSAGES

Code	Lev	Sev	Message
CCS	31	31	<p><b>Cannot create section [octal]</b></p> <p>LINK is unable to create the specified section. This could be because your system does not have extended addressing hardware, or because there are insufficient resources to create a section.</p>
CFS	31	31	<p><b>Chained fixups have been suppressed</b></p> <p>The specified PSECT grew beyond the address specified in the /LIMIT switch. The program is probably incorrect. Use the /MAP or /COUNTER switch to check for accidental PSECT overlaps. Refer to Section 3.2.2 for more information about the /LIMIT switch.</p>
CLF	1	1	<p><b>Closing log file, continuing on file [file]</b></p> <p>You have changed the log file specification. The old log file is closed; further log entries are written in the new log file.</p>
CM6	31	31	<p><b>Cannot mix COBOL-68 and COBOL compiled code {Detected in module [symbol] from file [file]}</b></p> <p>You cannot use COBOL-68 and COBOL files in the same load. Compile all COBOL programs with the same compiler and reload.</p>
CM7	31	31	<p><b>Cannot mix COBOL-74 and COBOL compiled code {Detected in module [symbol] from file [file]}</b></p> <p>You cannot use COBOL-74 and COBOL files in the same load. Compile all COBOL programs with the same compiler and reload.</p>
CMC	31	31	<p><b>Cannot mix COBOL-68 and COBOL-74 compiled code {Detected in module [symbol] from file [file]}</b></p> <p>You cannot use COBOL-68 and COBOL-74 files in the same load. Compile all COBOL programs with the same compiler and reload.</p>
CMF	31	31	<p><b>COBOL module must be loaded first {Detected in module [symbol] from file [file]}</b></p> <p>You are loading a mixture of COBOL-compiled files and other files. Load one of the COBOL-compiled files first.</p>
CMP	31	28	<p><b>Common [symbol] declared in multiple PSECTS {Detected in module [symbol] from file [file]}</b></p> <p>You cannot load modules produced by FORTRAN with modules produced by G-floating FORTRAN. Compile all FORTRAN modules the same way, then reload.</p>

LINK MESSAGES

Code	Lev	Sev	Message
CMX	8	8	<p>Cannot mix G-floating FORTRAN compiled code with FORTRAN compiled code  {Detected in module [symbol] from [file]}</p> <p>You loaded a module which specified that the named common block must be loaded in a PSECT which is not compatible with the PSECT in which it was originally loaded. Compile the module with the common in the same PSECT as the original.</p>
CNW	31	31	<p>Code not yet written at [label]</p> <p>You attempted to use an unimplemented feature. This is an internal LINK error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
COE	8	8	<p>Both CONCATENATE and OVERLAY attributes specified for PSECT [name]  {Detected in module [symbol] from file [file]}</p> <p>One of the modules you loaded explicitly sets an attribute for the named PSECT which conflicts with the declaration of PSECT attributes in the current module. Check the compiler switches or assembly language directives that were used in the generation of this module.</p>
COF	*	*	<p>Cannot open file [file]</p> <p>LINK cannot open the specified file for input.</p>
CPU	31	31	<p>Module incompatible with specified CPU  {Detected in module [symbol] from file [file]}</p> <p>The module you are attempting to load does not contain a .DIRECTIVE for any of the CPUs you specified with the /CPU switch. Recompile the module with the proper .DIRECTIVE, or use a different /CPU switch.</p>
CRS	1	1	<p>Creating section [octal]</p> <p>LINK prints this informational message when a module is loaded into a new section. The message is printed only if you have specified /ERROR:Ø.</p>
CSF	1	1	<p>Creating saved file</p> <p>LINK is generating your executable (.EXE) file.</p>
DEB	31	1	<p>[name] execution</p> <p>LINK is beginning program execution at the named debugger.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
DLT	31	1	<p><b>Execution deleted</b></p> <p>Though you have asked for program execution, LINK cannot proceed due to earlier fatal compiler or LINK errors. Your program is left in memory or in an executable file.</p>
DNS	8	8	<p><b>Device not specified for switch [switch]</b></p> <p>You used a device switch (for example, /REWIND, /BACKSPACE), but LINK cannot associate a device with the switch. Neither LINK's default device nor any device you gave with the /DEFAULT switch can apply. Give the device with or before the switch (in the same command line).</p>
DRC	8	8	<p><b>Decreasing relocation counter [symbol] from [octal] to [octal] {Detected in module [symbol] from file [file]}</b></p> <p>You are using the /SET switch to reduce the value of an already defined relocation counter. Unless you know exactly where each module is loaded, code may be overwritten.</p>
DSC	31	31	<p><b>Data store to common [symbol] not in link number [decimal] {Detected in module [symbol] from file [file]}</b></p> <p>You loaded a FORTRAN-compiled module with DATA statement assignments to a common area. The common area is already defined in an ancestor link. Restructure the load so that the DATA statements are loaded in the same link as the common area to which they refer.</p>
DSL	31	*	<p><b>Data store to location [octal] not in link number [decimal] {Detected in module [symbol] from file [file]}</b></p> <p>You have a data store for an absolute location outside the specified link. Load the module into the root link.</p>
NOTE			
<p>If the location is less than 140, this message has level 8 and severity 8.</p>			
DUZ	31	31	<p><b>Decreasing undefined symbol count below zero</b></p> <p>LINK's undefined symbol count has become negative. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>



## LINK MESSAGES

Code	Lev	Sev	Message
EAS	31	31	<p><b>Error creating area AS overflow file [reason] [file]</b></p> <p>LINK could not make the ALGOL symbol table on disk. You could be over your disk quota, or the disk could be full or have errors.</p>
ECE	31	31	<p><b>Error creating EXE file [reason] [file]</b></p> <p>LINK could not write the saved file on disk. You could be over your disk quota, or the disk could be full or have errors.</p>
EGD	1	1	<p><b>Emergency GETSEG done</b></p> <p>LINK has expanded its low segment so large that one of its larger high segments will not now fit in memory. LINK will attempt to shrink its internal tables so that the GETSEG will succeed and loading continue.</p>
EHC	31	31	<p><b>Error creating area HC overflow file [reason] [file]</b></p> <p>LINK could not write your high-segment code on disk. You could be over your disk quota, or the disk could be full or have errors.</p>
EIF	31	31	<p><b>Error for input file Status [octal] for file [file]</b></p> <p>A read error has occurred on the input file. Use of the file is terminated and the file is released.</p>
ELC	31	31	<p><b>Error creating area LC overflow file [reason] [file]</b></p> <p>LINK could not write your low-segment code on the disk. You could be over your disk quota, or the disk could be full or have errors.</p>
ELF	1	1	<p><b>End of log file</b></p> <p>LINK has finished writing your log file. The file is closed.</p>
ELN	1	1	<p><b>End of link number [decimal] {name [name]}</b></p> <p>The link is loaded.</p>
ELS	31	31	<p><b>Error creating area LS overflow file [reason] [file]</b></p> <p>LINK could not write your local symbol table on the disk. You could be over your disk quota, or the disk could be full or have errors.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
EMS	1	1	<p><b>End of MAP segment</b></p> <p>The map file is completed and closed.</p>
EOE	31	31	<p><b>.EXE file output error Status [octal] for file [file]</b></p> <p>LINK could not write the saved file on the disk.</p>
EOI	31	31	<p><b>Error on input Status [octal] for file [file]</b></p> <p>An error has been detected while reading the named file.</p>
EOO	31	31	<p><b>Error on output [file]</b></p> <p>An error has been detected while writing the named file.</p>
EOV	31	31	<p><b>Error creating overlay file [reason] [file]</b></p> <p>LINK could not write the overlay file on the disk.</p>
ESN	31	31	<p><b>Extended symbol not expected</b></p> <p>Long symbol names (more than six characters) are not implemented. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ETP	31	31	<p><b>Error creating area TP overflow file [reason] [file]</b></p> <p>LINK could not make the type checking area on disk. You are over your disk quota, the disk is full, or the disk has errors.</p>
EXP	1	1	<p><b>Expanding low segment to [addr]</b></p> <p>This informational message is printed when LINK expands the low segment memory allocation.</p>
EXS	1	1	<p><b>EXIT segment</b></p> <p>LINK is in the last stages of loading your program (for example, creating .EXE and symbol files, preparing for execution if requested).</p>
FCF	1	1	<p><b>Final code fixups</b></p> <p>LINK is reading one or both segment overflow files backwards to perform any needed code fixups. This may cause considerable disk overhead, but occurs only if your program is too big for memory.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
FEE	*	*	<p><b>ENTER error [reason] [file]</b></p> <p>See Section B.3 for the list of possible long messages.</p>
FIN	1	1	<p><b>LINK finished</b></p> <p>LINK is finished. Control is passed to the monitor, or to the loaded program for execution.</p>
FLE	*	*	<p><b>LOOKUP error [reason] [file]</b></p> <p>See Section B.3 for the list of possible long messages.</p>
FRE	*	*	<p><b>RENAME error [reason] [file]</b></p> <p>See Section B.3 for the list of possible long messages.</p>
FSN	31	31	<p><b>FUNCT. subroutine not loaded</b></p> <p>During final processing of your root link, LINK found that the FUNCT. subroutine was not loaded. This would cause an infinite recursion if your program were executed. The FUNCT. subroutine is requested by the overlay handler, and is usually loaded from a default system library. Either you prevented searching of system libraries, or you did not load a main program from an overlay-supporting compiler into the root link.</p>
FTH	15	15	<p><b>Fullword value [symbol] truncated to halfword</b></p> <p>This message is printed when a symbol that has a value greater than 777777 is used to resolve a halfword reference. This warning message helps you to be sure that global addresses are used properly throughout the modules in a load.</p>
GSE	*	*	<p><b>GETSEG error [reason] [file]</b></p> <p>See Section B.3 for the list of possible long messages.</p>
HCL	31	31	<p><b>High segment code not allowed in an overlay link {Detected in module [symbol] from file [file]}</b></p> <p>You have attempted to load high segment code into an overlay link other than the root. Any high segment code in an overlaid program must be in the root.</p>

LINK MESSAGES

Code	Lev	Sev	Message
HSL	31	31	<p><b>Attempt to set high segment origin too low</b>  {Detected in module [symbol] from file [file]}</p> <p>You have set the high-segment counter to a page containing low-segment code. Reload, using the /SET:.HIGH.:n switch, or (for MACRO programs) reassemble after changing your TWOSEG pseudo-op.</p>
HTL	31	31	<p><b>Symbol hash table too large</b></p> <p>Your symbol hash table is larger than the maximum LINK can generate (about 50P). This table size is an assembly parameter. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
IAS	31	31	<p><b>Error inputting area as Status [octal] for file [file]</b></p> <p>An error occurred while reading in the ALGOL symbol table.</p>
ICB	8	8	<p><b>Invalid chain REL block (type 12) link number [octal]</b>  {Detected in module [symbol] from file [file]}</p> <p>REL block type 12 (Chain), generated by the MACRO pseudo-op .LINK and .LNKEND, must contain a number from 1 to 100 (octal) in its first word. The link word is ignored.</p>
IDM	31	31	<p><b>Illegal data mode for device [device]</b></p> <p>You specified an illegal combination of device and data mode (for example, terminal and dump mode). Specify a legal device.</p>
IHC	31	31	<p><b>Error inputting area HC Status [octal] for file [file]</b></p> <p>An error occurred while reading in your high-segment code.</p>
ILC	31	31	<p><b>Error inputting area LC Status [octal] for file [file]</b></p> <p>An error occurred while reading in your low-segment code.</p>
ILS	31	31	<p><b>Error inputting area LS Status [octal] for file [file]</b></p> <p>An error occurred while reading in your local symbol table.</p>
IMA	8	8	<p><b>Incremental maps not yet available</b></p> <p>The INCREMENTAL keyword for the /MAP switch is not implemented. The switch is ignored.</p>

LINK MESSAGES

Code	Lev	Sev	Message
IMI	31	31	<p>Insufficient memory to initialize LINK</p> <p>LINK needs more memory than is available.</p>
IMM	*	1	<p>[Decimal] included modules missing {from file [file]}</p> <p>You have requested with the /INCLUDE switch that the named modules (if any) be loaded. Specify files containing these modules.</p>
INS	31	31	<p>I/O data block not set up</p> <p>LINK attempted a monitor call (for example, LOOKUP, ENTER) for a channel that is not set up. This is an internal LINK error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
IOV	31	31	<p>Input error for overlay file Status [octal] for file [file]</p> <p>An error occurred when reading the overlay file.</p>
IPO	31	31	<p>Invalid Polish operator [octal] {Detected in module [symbol] from file [file]}</p> <p>You are attempting to load a file containing an invalid REL Block Type 11 (Polish). This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
IPX	31	31	<p>Invalid PSECT index {for PSECT [symbol]} {Detected in module [symbol] from file [file]}</p> <p>A REL block contains a reference to a nonexistent PSECT. This error is probably caused by a fault in the language translator used for the program. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
IRB	31	31	<p>Illegal REL block type [octal] {Detected in module [symbol] from file [file]}</p> <p>The file is not in the proper binary format. It may have been generated by a translator that LINK does not recognize, or it may be an ASCII or .EXE file.</p>
IRC	31	31	<p>Illegal relocation counter {Detected in module [symbol] from file [file]}</p> <p>One of the new style 1000+ block types has an illegal relocation counter. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
IRR	8	8	<p><b>Illegal request/require block</b>  {Detected in module [symbol] from file [file]}</p> <p>One of the REL block types 1042 or 1043 is in the wrong format. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ISM	31	31	<p><b>Incomplete symbol in store operator in Polish block (type 11)</b>  {Detected in module [symbol] from file [file]}</p> <p>The specified module contains an incorrectly formatted Polish Fixup Block (Type 11). The store operator specifies a symbol fixup, but the block ends before the symbol is fully specified. This error is probably caused by a fault in the language translator used for the program. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ISN	31	31	<p><b>Illegal symbol name [symbol]</b>  {Detected in module [symbol] from file [file]}</p> <p>The LINK symbol table routine was called with the blank symbol. This error can be caused by a fault in the language translator used for the program. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ISP	31	31	<p><b>Incorrect symbol pointer</b></p> <p>There is an error in the global symbol table. This is an internal LINK error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ISS	8	8	<p><b>Insufficient space for symbol table after PSECT [symbol] -- table truncated</b></p> <p>There is insufficient address space for the symbol table between the named PSECT and the next higher one or the end of the address space. Restructure your PSECT layout to allow sufficient room for the symbol table, or use /UPTO to allow more room.</p>
IST	31	31	<p><b>Inconsistency in switch table</b></p> <p>LINK has found errors in the switch table passed from the SCAN module. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
ITB	31	31	<p>Invalid text in ASCII block from file [file]</p> <p>LINK has failed to complete the processing of an ASCII text REL block from the named file. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
ITP	31	31	<p>Error inputting area TP {Status [octal]} for file [file]</p> <p>An error occurred while reading in the type checking area. The status is represented by the right half of the I/O status word. Refer to the <u>TOPS-10 Monitor Calls Manual</u> for more information.</p>
IUU	*	31	<p>Illegal user UUU at PC [octal]</p> <p>LINK's user UUU (LUUU) handler has detected an illegal UUU. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
IVC	31	31	<p>Index validation check failed at address [octal]</p> <p>The range checking of LINK's internal tables and arrays failed. The address given is the point in a LINK segment at which failure occurred. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
JPB	8	8	<p>Junk at end of Polish block {Detected in module [symbol] from file [file]}</p> <p>The specified module contains an incorrectly formatted Polish Fixup Block (Type 11). Either the last unused halfword (if it exists) is nonzero, or there are extra halfwords following all valid data.</p>
LDS	1	1	<p>LOAD segment</p> <p>The LINK module LNKLOD is beginning its processing.</p>
LFB	1	1	<p>LINK log file begun on [date]</p> <p>LINK is creating your log file as a result of defining the logical name LOG:.</p>
LFC	1	1	<p>Log file continuation</p> <p>LINK is continuing your log file as a result of the /LOG switch.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
LFI	1	1	<p><b>Log file initialization</b></p> <p>LINK is beginning your log file as a result of the /LOG switch.</p>
LII	8	1	<p><b>Library index inconsistent, continuing</b></p> <p>A REL Block Type 14 (Index) for a MAKLIB generated library file is inconsistent. The library is searched, but the index is ignored.</p>
LIN	1	1	<p><b>LINK initialization</b></p> <p>LINK is beginning its processing by initializing its internal tables and variables.</p>
LMM	‡	‡	<p><b>Length mismatch for argument [decimal] in call to routine [symbol] called from module [symbol] at location [octal]</b></p> <p>The length of the argument passed by the caller does not match what the called routine expects it to be.</p>
LMN	6	1	<p><b>Loading module [symbol] from file [file] LINK is loading the named module.</b></p>
LNA	8	8	<p><b>Link name [name] already assigned to link number [decimal]</b></p> <p>You used this name for another link. Specify a different name for this link.</p>
LNL	8	8	<p><b>Link number [decimal] not loaded</b></p> <p>The link with this number has not yet been loaded. The /NODE switch is ignored. If you have used link numbers instead of link names with the /NODE switch, you may have confused the link numbers. To avoid this, use link names.</p>
LNM	31	31	<p><b>Link number [decimal] not in memory</b></p> <p>LINK cannot find the named link in memory. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>

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‡ The level and severity of this message is determined by a compiler-generated coercion block. See Block Type 1130 in Appendix A.



## LINK MESSAGES

Code	Lev	Sev	Message
LNN	8	8	<p><b>Link name [name] not assigned</b></p> <p>The name you gave with the /NODE switch is not the name of any loaded link. The switch is ignored.</p>
LNS	31	8	<p><b>Low segment data base not same size</b></p> <p>The length of LINK's low segment differs from the length stored in the current LINK high segment. This occurs if some but not all of LINK's .EXE files have been updated after rebuilding LINK from sources. Update all of LINK's .EXE files.</p>
LSM	8	8	<p><b>/LINK switch missing while loading link number [decimal] -- assumed</b></p> <p>Your use of the /NODE switch shows that you want to begin a new overlay link, but the current link is not yet completely loaded. LINK assumes a /LINK switch immediately preceding the /NODE switch, and loads the link (without a link name).</p>
LSS	31	1	<p><b>{No} Library search symbols (entry points) {[symbol] [octal]}</b></p> <p>The listed symbols and their values (if any) are those that are library search entry points.</p>
MDS	8	8	<p><b>Multiply-defined global symbol [symbol] {Detected in module [symbol] from file [file]} Defined value = [octal], this value = [octal]</b></p> <p>The named module contains a new definition of an already defined global symbol. The old definition is used. Make the definitions consistent and reload.</p>
MEF	31	31	<p><b>Memory expansion failed</b></p> <p>LINK cannot expand memory further. All permitted overflows to disk have been tried, but your program is still too large for available memory. A probable cause is a large global symbol table, which cannot be overflowed to disk. It may be necessary to restructure your program, or use overlays, to alleviate this problem.</p>
MMF	31	31	<p><b>Memory manager failure</b></p> <p>The internal memory manager in LINK has failed a consistency check. This error should not occur. If it does, contact your software specialist or send a Software Performance Report (SPR) to DIGITAL.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
MOV	1	1	<p><b>Moving low segment to expand area [area]</b></p> <p>LINK is rearranging its low segment to make more room for the specified area. Area is one of the following:</p> <p>AS ALGOL symbol table            BG bound global symbols            DY dynamic free memory            FX fixup area            GS global symbol table            HC your high-segment code            LC your low-segment code            LS local symbol tables            RT relocation tables</p>
MPS	1	1	<p><b>MAP segment</b></p> <p>The LINK module LNKMAP is writing a map file.</p>
MPT	31	31	<p><b>Mixed PSECT and TWOSEG code in same module {Detected in module [symbol] from file [file]}</b></p> <p>This module contains both PSECT code and TWOSEG code. LINK cannot load such a module. Change the source code to use PSECTS .HIGH. and .LOW. as the high and low segments, and remove the TWOSEG or HISEG pseudo-ops.</p>
MRN	1	1	<p><b>Multiple regions not yet implemented</b></p> <p>The REGION keyword for the /OVERLAY switch is not implemented. The argument is ignored.</p>
MSN	8	8	<p><b>Map sorting not yet implemented</b></p> <p>Alphabetical or numerical sorting of the map file is not implemented. The symbols in the map file appear in the order they are found in the REL files.</p>
MSS	8	8	<p><b>/MAXCOR: set too small, expanding to [memory]</b></p> <p>The current value of MAXCOR is too small for LINK to operate. You can speed up future loads of this program by setting the /MAXCOR switch to this expanded size at the beginning of the load.</p>
MTB	8	8	<p><b>/MAXCOR: too big, [memory] used</b></p> <p>You are attempting to specify the /MAXCOR switch so large that the low segment cannot fit before the high segment. LINK will use only the core indicated.</p>

LINK MESSAGES

Code	Lev	Sev	Message
MTS	8	8	<p><b>/MAXCOR: too small, at least [memory] is required</b></p> <p>LINK needs more space than you gave with the /MAXCOR switch. Give a new /MAXCOR switch with at least the required size.</p>
NAP	31	31	<p><b>No store address in polish block (Type 11 or 1072)</b>  <b>{Detected on module [symbol] from R6 [file]}</b></p> <p>The specified module contains an incorrectly formatted polish fixup block (Type 11 or 1072). The store operator specifies a memory fixup, but the block ends before the address is specified. This error is probably caused by a fault in the language translator used for the program. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
NBR	31	31	<p><b>Attempt to position to node before the root</b></p> <p>The argument you gave for the /NODE switch would indicate a link before the root link. (For example, from a position after the third link in a path, you cannot give /NODE:-4.)</p>
NEB	8	8	<p><b>No end block seen</b>  <b>{Detected in module [symbol] from file [file]}</b></p> <p>No REL Block Type 5 (End) was found in the named module. This will happen if LINK finds two Type 6 blocks (Name) without an intervening end, or if an end-of-file is found before the end block is seen. LINK simulates the missing end block. However, fatal messages usually follow this, because this condition usually indicates a bad REL file.</p>
NED	31	24	<p><b>Non-existent device [device]</b></p> <p>You gave a device that does not exist on this system. Correct your input files and reload.</p>
NHN	31	31	<p><b>No high segment in non-zero section</b></p> <p>You have attempted to load high segment code into a program with a non-zero section. High segments must reside in Section 0.</p>
NPS	8	8	<p><b>Non-existent PSECT [symbol] specified for symbol table</b></p> <p>You have specified the name of a PSECT after which LINK should append the symbol table, but no PSECT with that name was loaded. Load the named PSECT or specify an existing PSECT for the symbols.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
NSA	31	1	<p><b>No start address</b></p> <p>Your program does not have a starting address. This can happen if you neglect to load a main program. Program execution, if requested, will be suppressed unless you specified debugger execution.</p>
NSM	31	31	<p><b>/NODE switch missing after /LINK switch</b></p> <p>You used the /LINK switch, which indicates that you want to begin a new overlay link, but you have not specified a /NODE switch to tell LINK where to put the new overlay link.</p>
NSO	31	31	<p><b>No store operator in Polish block (type 11) {Detected in module [symbol] from file [file]}</b></p> <p>The specified module contains an incorrectly formatted Polish Fixup Block (Type 11). Either the block does not have a store operator, or LINK was not able to detect it due to the block's invalid format. This error is probably caused by a fault in the language translator used for the program. This error is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
NVR	‡	‡	<p><b>No value returned by routine [symbol] called from module [symbol] at location [octal]</b></p> <p>The called routine does not return a value, however the caller expected a returned value.</p>
OAS	31	31	<p><b>Error outputting area as Status [octal] for file [file]</b></p> <p>An error occurred while writing out the ALGOL symbol table.</p>
OEL	8	8	<p><b>Output error on log file, file closed, load continuing {Status [octal] for file [file]}</b></p> <p>An error has occurred on the output file. The output file is closed at the end of the last data successfully output.</p>
OEM	8	8	<p><b>Output error on map file, file closed, load continuing Status [octal] for file [file]</b></p> <p>An error has occurred on the output file. The output file is closed at the end of the last data successfully output.</p>

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‡ The level and severity of this message is determined by a compiler-generated coercion block. See Block Type 1130 in Appendix A.

## LINK MESSAGES

Code	Lev	Sev	Message
OES	8	8	<p><b>Output error on symbol file, file closed, load continuing Status [octal] for file [file]</b></p> <p>An error has occurred on the output file. The output file is closed at the end of the last data successfully output.</p>
OEX	8	8	<p><b>Output error on XPN file, file closed, load continuing Status [octal] for file [file]</b></p> <p>An error has occurred on the output file. The output file is closed at the end of the last data successfully output.</p>
OPD	31	31	<p><b>OPEN failure for device [device]</b></p> <p>An OPEN or INIT monitor call for the specified device failed. The device may be under another user's control.</p>
OFN	31	31	<p><b>Old FORTRAN (F40) module not available {Detected in module [symbol] from file [file]}</b></p> <p>The standard released version of LINK does not support F40 code.</p>
OFS	31	31	<p><b>Overlay file must be created on a file structure</b></p> <p>Specify a disk device for the overlay file.</p>
OHC	31	31	<p><b>Error outputting area HC Status [octal] for file [file]</b></p> <p>An error occurred while writing out your high-segment code.</p>
OHN	31	31	<p><b>Overlay handler not loaded</b></p> <p>Internal symbols in the overlay handler could not be referenced. If you are using your own overlay handler, this is a user error; if not, it is an internal error and is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
OLC	31	31	<p><b>Error outputting area LC Status [octal] for file [file]</b></p> <p>An error occurred while writing out your low-segment code.</p>
OLS	31	31	<p><b>Error outputting area LS Status [octal] for file [file]</b></p> <p>An error occurred while writing out your local symbol table.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
OMB	31	31	<p><b>/OVERLAY switch must be first</b></p> <p>The /OVERLAY switch must appear before you can use any of the following switches: /ARSize, /LINK, /NODE, /PLOT, /SPACE. (It is sufficient that the /OVERLAY switch appear on the same line as the first of these switches you use.)</p>
ONS	8	1	<p><b>Overlays not supported in this version of LINK</b></p> <p>LINK handles overlays with its LNKOV1 and LNKOV2 modules. Your installation has substituted dummy versions of these. You should request that your installation rebuild LINK with the real LNKOV1 and LNKOV2 modules.</p>
OOV	31	31	<p><b>Output error for overlay file Status [octal] for file [file]</b></p> <p>An error has occurred while writing the overlay file.</p>
OS2	1	1	<p><b>Overlay segment phase 2</b></p> <p>LINK's module LNKOV2 is writing your overlay file.</p>
OSL	8	8	<p><b>Overlaid program symbols must be in low segment</b></p> <p>You have specified /SYMSEG:HIGH or /SYMSEG:PSECT when loading an overlay structure. Specify /SYMSEG:LOW or /SYMSEG:DEFAULT.</p>
OTP	31	31	<p><b>Error outputting area TP {Status [octal] for file [file]}</b></p> <p>An error occurred while writing out the typechecking area. The status is represented by the right half of the I/O status word (refer to the <u>TOPS-10 Monitor Calls Manual</u>).</p>
PAS	1	1	<p><b>Area AS overflowing to disk</b></p> <p>The load is too large to fit into the allowed memory and the ALGOL symbol table is being moved to disk.</p>
PBI	8	8	<p><b>Program break [octal] invalid {Detected in module [symbol] from file [file]}</b></p> <p>The highest address allocated in the named module is greater than 512P. This is usually caused by dimensioning large arrays. Modify your programs or load list to reduce the size of the load.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
PCL	8	8	<p><b>Program too complex to load, saving as file [file]</b></p> <p>Your program is too complex to load into memory for one of the following reasons:</p> <ul style="list-style-type: none"><li>• There are page gaps between PSECTs (except below the high segment).</li><li>• There are PSECTs above the origin of the high segment.</li><li>• Your program will not fit in memory along with LINK's final placement code.</li><li>• One or more PSECTs has the read-only attribute.</li></ul> <p>LINK has saved your program as an .EXE file on disk and cleared your user memory. You can use a GET or RUN command to load the .EXE file.</p>
PCX	8	1	<p><b>Program too complex to load and execute, will run from file [file]</b></p> <p>Your program is too complex to load into memory for one of the following reasons:</p> <ul style="list-style-type: none"><li>• There are page gaps between PSECTs (except below the high segment).</li><li>• There are PSECTs above the origin of the high segment.</li><li>• Your program will not fit in memory along with LINK's final placement code.</li><li>• One or more PSECTs has the read-only attribute.</li></ul> <p>LINK will save your program as an .EXE file on disk and automatically run it, but the .EXE file will not be deleted.</p>
PEF	31	8	<p><b>Premature end of file from file [file]</b></p> <p>LINK found an end-of-file inside a REL block (that is, the word count for the block extended beyond the end-of-file). This error may be caused by a fault in the language translator used for the program.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
PEL	15	15	<p><b>PSECT [symbol] exceeded limit of [octal]</b></p> <p>The specified PSECT grew beyond the address specified in the /LIMIT switch. The program is probably incorrect. Use the /MAP or /COUNTER switch to check for accidental PSECT overlaps. Refer to Section 3.2.2 for more information about the /LIMIT switch.</p>
PHC	1	1	<p><b>Area HC overflowing to disk</b></p> <p>The load is too large to fit into the allowed memory and your high-segment code is being moved to disk.</p>
PLC	1	1	<p><b>Area LC overflowing to disk</b></p> <p>The load is too large to fit into the allowed memory and your low-segment code is being moved to disk.</p>
PLS	1	1	<p><b>Area LS overflowing to disk</b></p> <p>The load is too large to fit into the allowed memory and your local</p>
PMA	‡	‡	<p><b>Possible modification of argument [decimal] in call to routine [symbol] called from module [symbol] at location [octal]</b></p> <p>The caller has specified that the argument should not be modified. The called routine contains code which may modify this argument. In some cases this message will occur although the argument is not actually modified by the routine.</p>
POT	1	1	<p><b>Plotting overlay tree</b></p> <p>LINK is creating your overlay tree file.</p>
POV	8	8	<p><b>PSECTs [symbol] and [symbol] overlap from address [octal] to [octal]</b></p> <p>The named PSECTs overlap each other in the indicated range of addresses. If you do not expect this message, restructure your PSECT origins with the /SET switch.</p>
PTP	1	1	<p><b>Area TP overflowing to disk</b></p> <p>The load is too large to fit into the allowed memory and your argument typechecking tables are being moved to disk.</p>

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‡ The level and severity of this message is determined by a compiler-generated coercion block. See Block Type 1130 in Appendix A.



## LINK MESSAGES

Code	Lev	Sev	Message
PTL	31	31	<p><b>Program too long</b>  {Detected in module [symbol] from file [file]}</p> <p>Your program extends beyond location 777777, which is the highest location that LINK is capable of loading. You may be able to make your program fit by moving PSECT origins, lowering the high-segment origin, loading into a single segment, reducing the size of arrays in your program, or using the overlay facility.</p>
PUF	31	31	<p><b>PAGE. UWO failed, error code was [code]</b></p> <p>Refer to the description of the PAGE. UWO in the TOPS-10 Monitor Calls Manual.</p>
RBS	31	31	<p><b>REL block type [octal] too short</b>  {Detected in module [symbol] from file [file]}</p> <p>The REL block is inconsistent. This may be caused by incorrect output from a translator (for example, missing argument for an end block). Recompile the module and reload.</p>
RED	1	1	<p><b>Reducing low segment to [memory]</b></p> <p>LINK is reclaiming memory by deleting its internal tables.</p>
RER	*	1	<p><b>{No} Request external references (inter-link entry points)</b>  {[symbol] [octal]}</p> <p>The listed symbols and their values (if any) represent subroutine entry points in the current link.</p>
RGS	1	1	<p><b>Rehashing global symbol table from [decimal] to [decimal]</b></p> <p>LINK is expanding the global symbol table either to a prime number larger than your /HASHSIZE switch requested, or by about 50 percent. You can speed up future loads of this program by setting /HASHSIZE this large at the beginning of the load.</p>
RLC	31	1	<p><b>Reloc ctr. initial value current value</b>  {[symbol] [octal] [octal]}</p> <p>The listed symbols and values represent the current placement of PSECTs in your address space.</p>

LINK MESSAGES

Code	Lev	Sev	Message
RME	31	*	<p>REMAP error{, high segment origin may be incorrect}</p> <p>The REMAP UUO to place your program's high segment has failed. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
RUM	31	31	<p>Returning unavailable memory</p> <p>LINK attempted to return memory to the memory manager, but the specified memory was not previously allocated. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
RWA	8	8	<p>Both READ-ONLY and WRITABLE attributes specified for PSECT [name]</p> <p>One of the modules you loaded explicitly sets an attribute for the named PSECT which conflicts with the declaration of PSECT attributes in the current module. Check the compiler switches or assembly language directives that were used in the generation of this module.</p>
SFU	8	8	<p>Symbol table fouled up</p> <p>There are errors in the local symbol table. Loading continues, but any maps you request will not contain control section lengths. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
SIF	31	31	<p>Symbol insert failure, non-zero hole found</p> <p>LINK's hashing algorithms failed; they are trying to write a new symbol over an old one. You may be able to load your files in a different order. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
SMP	8	8	<p>SIMULA main program not loaded</p> <p>You loaded some SIMULA procedures or classes, but no main program. Missing start address and undefined symbols will terminate execution.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
SNC	31	31	<p><b>Symbol [symbol] already defined, but not as common</b>  <b>{Detected in module [symbol] from file [file]}</b></p> <p>You defined a FORTRAN common area with the same name as a non-common symbol. You must indicate which definition you want. If you want the common definition, load the common area first.</p>
SNL	1	1	<p><b>Scanning new command line</b></p> <p>LINK is ready to process the next command line.</p>
SNP	8	8	<p><b>Subroutine [symbol] in link number [decimal] not on path for call from link number [decimal]</b>  <b>{name [name]}</b></p> <p>The named subroutine is in a different path than the calling link. Redefine your overlay structure so that the subroutine is in the correct path.</p>
SNS	31	31	<p><b>SITGO not supported</b>  <b>{Detected in module [symbol] from file [file]}</b></p> <p>LINK does not support the REL file format produced by the SITGO compiler. Load your program by using SITGO.</p>
SOE	31	31	<p><b>Saved file output error Status [octal] for file [file]</b></p> <p>An error occurred in outputting the .EXE file.</p>
SRB	8	8	<p><b>Attempt to set relocation counter [symbol] below initial value of [octal]</b>  <b>{Detected in module [symbol] from file [file]}</b></p> <p>You cannot use the /SET switch to set the named relocation counter below its initial value. The attempt is ignored.</p>
SRP	31	31	<p><b>/SET: switch required for PSECT [symbol]</b>  <b>{Detected in module [symbol] from file [file]}</b></p> <p>Relocatable PSECTS are not implemented; you must specify an explicit absolute origin with the /SET switch for the named PSECT.</p>
SSN	8	8	<p><b>Symbol table sorting not yet implemented</b></p> <p>Alphabetical or numerical sorting of the symbol table is not implemented. The symbols appear in the order they are found.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
SST	1	1	<p><b>Sorting symbol table</b></p> <p>LINK is rearranging the symbol table, and if required, is converting the symbols from the new to old format as indicated on the /SYMSEG, /SYFILE, or /DEBUG switch.</p>
STC	1	1	<p><b>Symbol table completed</b></p> <p>The symbol table has been sorted and moved according to the /SYMSEG, /SYFILE, or /DEBUG switch.</p>
STL	31	31	<p><b>Symbol too long</b></p> <p>A symbol specified in a REL block is larger than the maximum allowed by LINK.</p>
T13	31	31	<p><b>LVAR REL block (type 13) not implemented</b>  {Detected in module [symbol] from file [file]}</p> <p>REL Block Type 13 (LVAR) is obsolete. Use the MACRO pseudo-op TWOSEG.</p>
TDS	8	8	<p><b>Too late to delete initial symbols</b></p> <p>LINK has already loaded the initial symbol table. To prevent this loading, place the /NOINITIAL switch before the first file specification.</p>
TMA	31	8	<p><b>Too many ambiguous requests in link [decimal]</b>  {name [name]}, use /ARSIZE:[decimal]  {Detected in module [symbol] from file [file]}</p> <p>You have more ambiguous subroutine requests (indicated by LNKARL messages) than will fit in the table for this link. Continue loading. Your load will abort at the end with a LNKABT message; if you have loaded all modules, the message will give the size of the needed /ARSIZE switch for a reload.</p>
TML	31	31	<p><b>Too many links use /MAXNODE</b></p> <p>You specified more overlay links than were allowed by the current value for the /MAXNODE switch. Reload the program with a larger /MAXNODE value.</p>

## LINK MESSAGES

Code	Lev	Sev	Message
TMM	‡	‡	<p><b>Type mismatch seen for argument [decimal] in call to routine [symbol] called from module [symbol] at location [octal]</b></p> <p>The data type of the argument passed by the caller does not match what the called routine expects.</p>
TTF	8	8	<p><b>Too many titles found</b></p> <p>In producing the index for a map file, LINK found more program names than there are programs. The symbol table is in error. This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
UAR	8	8	<p><b>Undefined assign for [symbol] {Detected in module [symbol] from file [file]}</b></p> <p>The named symbol was referenced in a REL Block Type 100 (ASSIGN), but the symbol is undefined. This is generated with the MACRO pseudo-op .ASSIGN. The assignment is ignored. You should load a module that defines the symbol.</p>
UCB	8	8	<p><b>Unknown common [symbol] referenced</b></p> <p>A reference was made to a common block that does not exist.</p>
UGS	*	1	<p><b>{No} Undefined global symbols {[symbol] [octal]}</b></p> <p>The listed symbols and their values (if any) represent symbols not yet defined by any module. Each value is the first address in a chain of references for the associated symbol.</p> <p>If this message resulted automatically at the end of loading, this is a user error. In this case, the load will continue, leaving references to these symbols unresolved.</p>
UNS	31	31	<p><b>Universal file REL block (type 777) not supported from file [file]</b></p> <p>Extraction of symbols from a MACRO universal file is not implemented.</p>

---

‡ The level and severity of this message is determined by a compiler-generated coercion block. See Block Type 1130 in Appendix A.

## LINK MESSAGES

Code	Lev	Sev	Message
URC	31	1	<p><b>Unknown Radix-50 symbol code [octal] [symbol]</b>  <b>{Detected in module [symbol] from file [file]}</b></p> <p>In a REL Block Type 2 (Symbols), the first 4 bits of each word pair contain the Radix-50 symbol code. LINK found one or more invalid codes in the block. This error can be caused by a fault in the language translator used for the program.</p>
URV	‡	‡	<p><b>Unexpected return value in call to routine [symbol] called from module [symbol] at location [octal]</b></p> <p>The called routine returns a value which was not expected by the caller.</p>
USA	8	8	<p><b>Undefined start address [symbol]</b></p> <p>You gave an undefined global symbol as the start address. Load a module that defines the symbol.</p>
USB	8	8	<p><b>Undefined symbol in byte array (type 1004) block</b></p> <p>LINK has detected an undefined global symbol in a Type 1004 REL block. This global symbol is used to relocate a byte pointer and must be defined before the 1004 block that uses it is seen. This error is probably the result of an error in the language translator used to generate the REL file.</p>
USC	31	8	<p><b>Undefined subroutine [symbol] called from link number [decimal] {name [name]}</b></p> <p>The named link contains a call for a subroutine you have not loaded. If the subroutine is required for execution, you must reload, including the required module in the link.</p>
USD	8	8	<p><b>Undefined symbol [symbol] in spare data (type 1160)</b></p> <p>LINK has detected an undefined global symbol in a Type 1160 REL block. This global symbol is used to relocate a byte pointer and must be defined before the 1160 block that uses it is seen. This error is probably the result of an error in the language translator used to generate the REL file.</p>

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‡ The level and severity of this message is determined by a compiler-generated coercion block. See Block Type 1130 in Appendix A.

LINK MESSAGES

Code	Lev	Sev	Message
USI	8	16	<p><b>Undefined symbol [symbol] illegal in switch [switch]</b></p> <p>You have specified an undefined symbol to a switch that can only take a defined symbol or a number. Specify the correct switch value.</p>
UUA	8	8	<p><b>Undefined /UPTO: address [symbol]</b></p> <p>You gave the named symbol as an argument to the /UPTO switch, but the symbol was never defined. Load a module that defines the symbol, or change your argument to the /UPTO switch.</p>
VAL	31	1	<p><b>Symbol [symbol] [octal] [type]</b></p> <p>LINK has printed the specified symbol, its value and its attributes as requested.</p>
WNA	*	*	<p><b>Wrong number of arguments in call to routine [symbol] called from module [symbol] at location [octal]</b></p> <p>The number of arguments in the routine call is not the number of arguments expected by the called routine.</p>
XCT	31	1	<p><b>[Name] execution</b></p> <p>LINK is beginning execution of your program.</p>
ZSV	8	8	<p><b>Zero switch value illegal</b></p> <p>You omitted required arguments for a switch (for example, /REQUIRE with no symbols). Respecify the switch.</p>

## LINK MESSAGES

### B.3 INDEXED MESSAGES

The following pages list the indexed messages issued with the LINK messages LNKFEE, LNKFLE, LNKFRE, and LNKGSE. The medium-length messages in the table are substituted for the [reason] field in the main messages.

The level and severity of these messages depend on the particular file access error. Thus, each indexed message includes the level and severity actually assigned to the main message.

Idx	Lev	Sev	Message
Ø	31	31	(Ø) Illegal file name One of the following conditions occurred: <ul style="list-style-type: none"><li>o The specified file name was illegal.</li><li>o When updating a file, the specified file name did not match the file being updated.</li><li>o The RENAME monitor call following a LOOKUP monitor call failed.</li></ul>
Ø	31	24	(Ø) file was not found The named file was not found. Specify an existing file.
1	31	24	(1) no directory for project-programmer number The named directory does not exist on the named file structure, or the project-programmer number given was incorrect.
2	31	24	(2) protection failure You do not have sufficient access privileges to use the named file.
2	31	31	(2) directory full The directory on the DECTape has no room for the file. Delete some files from the DECTape or specify another device.
3	31	24	(3) file was being modified Another job is currently modifying the named file. Try accessing the file later.
4	31	24	(4) rename file name already exists The named file already exists, or a different file was specified on the ENTER monitor call following a LOOKUP monitor call.



## LINK MESSAGES

Idx	Lev	Sev	Message
5	31	31	<p>(5) illegal sequence of UUOs</p> <p>LINK has specified an illegal sequence of monitor calls. (For example, a RENAME without a preceding LOOKUP or ENTER monitor call, or a LOOKUP after an ENTER.) This is an internal error. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
6	31	31	<p>(6) bad UFD or bad RIB</p> <p>One of the following conditions occurred:</p> <ul style="list-style-type: none"><li>o A transmission, device or data error occurred while attempting to read the directory or the RIB of the named file.</li><li>o A hardware-detected device or data error was detected while reading the named directory's RIB or data block.</li><li>o A software-detected data inconsistency error was detected while reading the named directory's or file's RIB.</li></ul>
7	31	31	<p>(7) not a saved file</p> <p>The named file is not a saved file. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
10	31	31	<p>(10) not enough memory</p> <p>The system cannot supply enough memory to use as buffers or to read in a program. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
11	31	31	<p>(11) device not available</p> <p>The named device is currently not available. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>
12	31	31	<p>(12) no such device</p> <p>The named device does not exist. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.</p>

LINK MESSAGES

Idx	Lev	Sev	Message
13	31	31	(13) not two reloc reg. capability  The machine does not have a two-register relocation capability. This message can never occur and is included only for completeness of the LOOKUP, ENTER and RENAME error codes. If it does occur, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
14	31	24	(14) no room or quota exceeded  You have exceeded the quota of the named directory, or the entire capacity of the file structure. Delete some files, or specify a directory or structure with sufficient space.
15	31	24	(15) write lock error  The named device is write-locked. Specify a write-enabled device or ask the operator to write-enable the named device.
16	31	31	(16) not enough monitor table space  There is not enough internal monitor table space for the named file. Try the load at a later time.
17	1	1	(17) partial allocation only  Because of the named directory's quota or the available space on the file structure, the total number of blocks requested could not be allocated. A partial allocation was given.
20	31	31	(20) block not free on allocation  The block required by LINK is not available for allocation. This message can never occur and is included only for completeness of the LOOKUP, ENTER and RENAME error codes. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
21	31	31	(21) can't supersede (enter) an existing directory  You have attempted to supersede the named directory.
22	31	31	(22) can't delete (rename) a non-empty directory  You have attempted to delete a directory that is not empty. This message can never occur and is included only for completeness of the LOOKUP, ENTER and RENAME error codes. If it does occur, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.

## LINK MESSAGES

Idx	Lev	Sev	Message
23	31	24	(23) SFD not found  One of the sub-file directories in the named path was not found.
24	31	24	(24) search list empty  A LOOKUP or ENTER monitor call was performed on generic device DSK: and the search list is empty.
25	31	24	(25) SFD nested too deeply  You have attempted to access a sub-file directory nested deeper than the maximum level allowed.
26	31	24	(26) no-create on for specified SFD path  No file structure in your job's search list has both the no-create bit and the write-lock bit equal to zero, and has the named directory.
27	31	24	(27) segment not on swap space  A GETSEG monitor call was issued from a locked low segment to a high segment which is not a dormant, active or idle segment. This message can never occur and is included only for completeness of the LOOKUP, ENTER and RENAME error codes. If it does occur, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
30	31	24	(30) can't update file  A LOOKUP and ENTER monitor call was given to update a file, but the file cannot be updated for some reason. (For example, another user is superseding it or the file was deleted between the time of the LOOKUP and the ENTER).
31	31	24	(31) low segment overlaps high segment  The end of the low segment is above the beginning of the high segment.
32	31	31	(32) RUN not allowed when not logged in  An attempt has been made to run a program from a not-logged-in job. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.

LINK MESSAGES

Idx	Lev	Sev	Message
33	31	31	(33) file still has outstanding ENQ/DEQ locks  The ENQ/DEQ facility has been used for simultaneous updating of the named file, but some ENQ/DEQ requests are still outstanding and the file cannot be closed. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
34	31	31	(34) bad .EXE file directory format  The named file has a bad .EXE format directory. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
35	31	31	(35).EXE format files must have .EXE extension  An attempt has been made to run an .EXE-format file with a non-.EXE extension. .EXE format files (those with an internal directory) must have the extension .EXE. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
36	31	31	(36) .EXE file directory is too big  An attempt has been made to run an .EXE-format file with a directory that is too large for the monitor to handle. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.
37	31	31	(37) network capability exceeded for TSK:  The monitor's ability to accept another network connection has been exceeded.
40	31	31	(40) task is not available  The named task is not available. Specify an existing task name.
41	31	31	(41) undefined network node for TSK:  You have specified a network node that does not exist. Wait for the node to come up or specify an existing network node.
42	31	31	([octal]) Unknown cause  This message indicates that a LOOKUP, ENTER or RENAME monitor call error occurred with an error code larger in number than the errors LINK knows about. This message is not expected to occur. If it does, please notify your Software Specialist or send a Software Performance Report (SPR) to DIGITAL.

## APPENDIX C

### JOB DATA AREA LOCATIONS SET BY LINK

LINK sets a number of locations between 40 and 140 (octal) in the user's program. These locations are known as the Job Data Area (commonly abbreviated to JOBDAT). They are used by the TOPS-10 monitor. In addition, two segment programs will have a Vestigial Job Data Area of eight words following the high segment origin.

#### Job Data Area

Address	Mnemonic	Use
41	.JB41	HALT if not specified otherwise. Executes by LUUOs.
42	.JBERR	Right: Number of errors during loading.
74	.JBDDT	Left: Highest location occupied by DDT. Right: Start address of DDT if loaded.
115	.JBHRL	Left: High segment length. Right: Highest address in high segment.
116	.JBSYM	Left: Negative length of symbol table. Right: Address of table.  For extended symbol tables, this word contains a positive value which is the pointer to the symbol table vector (see Section 4.3).
117	.JBUSY	Left: Negative length of undefined symbol table. Right: Address of undefined symbol table.  If 0, there is no undefined symbol table, or the pointer to an extended symbol table is stored in .JBSYM.
120	.JBSA	Left: First free location in low segment. Right: Start address of program.
121	.JBFF	Right: First free location in low segment.
124	.JBREN	Right: Reenter address of program.

## JOB DATA AREA LOCATIONS SET BY LINK

Address	Mnemonic	Use
131	.JBOVL	Address of header block for the root link in an overlaid program.
133	.JBCOR	Left: Highest location of low segment loaded with data.
137	.JBVER	Version number: see description of /VERSION switch in Section 3.2.2.

### Vestigial Job Data Area

Offset	Mnemonic	Use
0	.JBHSA	Copy of .JBSA.
1	.JBH41	Copy of .JB41.
2	.JBHCR	Copy of .JBCOR.
3	.JBHRH	LH: left half of .JBHRL. RH: right half of .JBREN.
4	.JBHVR	Copy of .JBVER.
5	.JBHNM	Program Name.
6	.JBHSM	High segment symbol table, if any.
7	.JBHGA	High segment origin page in bits 9-17.

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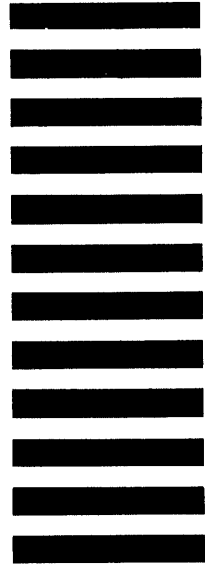


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### **TOPS-10 LINK Reference Manual AD-0988D-T1**


**January 1985**

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